

Article

A Power-Gated 8-Transistor Physically Unclonable Function Accelerates Evaluation Speeds

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Abstract: The proposed 8-Transistor (8T) Physically Unclonable Function (PUF), in conjunction with the power gating technique, can significantly accelerate a single evaluation cycle more than 100,000 times faster than a 6-Transistor (6T) Static Random-Access Memory (SRAM) PUF. The 8T PUF is built to swiftly eliminate data remanence and maximise physical mismatch. Moreover, a two-phase power gating module is devised to provide controllable power on/off cycles for the chosen PUF clusters in order to facilitate fast statistical measurements and curb the in-rush current. The architecture and hardware implementation of the power-gated PUF are developed to accommodate fast multiple evaluations of PUF Responses. The fast speed enables a new data processing method, which coordinates Dark-bit masking and Multiple Temporal Majority Voting (TMV) in different Process, Voltage and Temperature (PVT) corners or during field usage, hence greatly reducing the Bit Error Rate (BER) and the hardware penalty for error correction. The designs are based on the UMC 65 nm technology and aim to tape out an Application-Specific Integrated Circuit (ASIC) chip. Post-layout Monte Carlo (MC) simulations are performed with Cadence, and the extracted PUF Responses are processed with Matlab to evaluate the 8T PUF performance and statistical metrics for subsequent inclusion in PUF Responses, which comprise the novelty of this approach.

Keywords: physically unclonable function; PUF; power gating; SRAM; dark bit; metastability; data remanence; data retention; reset



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1. Introduction

1.1. Background

The Internet of Things (IoT) market is growing fast in annual revenue. However, this globally interconnected world poses a much more severe challenge to security. Due to the extremely resource-constrained nature of many IoT devices, e.g., wearables, lightweight security schemes and enhanced energy efficiency are consequently in great demand. To build a strong physical cyber-defence capability from the very start of the Integrated Circuit (IC) design [1], PUFs are promising hardware security primitives because they are easy to evaluate and physically hard to duplicate [2]. PUFs are functions that map Challenges to Responses through physically unclonable devices [2,3], and each PUF device is unique. Their uniqueness comes from the uncontrollable physical parameter mismatch generated in semiconductor fabrication. The proposed 8T PUF is derived from conventional 6T-SRAM. SRAM is an indispensable part of mainstream embedded designs because of its symmetric structure and mass entropy. The SRAM PUF was first proposed by Guajardo et al. [4] and Holcomb et al. [5], who discovered the intrinsically random start-up values of SRAM cells. Accordingly, these repeatable start-up values are the raw data for creating the Response of the PUF, and the address used to read them is the Challenge, as shown in Figure 2.

However, in a small number of SRAM cells, random logical states inevitably appear after every power-up due to negligible physical mismatch. They cause unreliable PUF readings, which cannot be tolerated for authentication. Moreover, some stable cells become

unstable with the change in environmental conditions, such as temperature or supply voltage, because the physical mismatch in cells is affected, e.g., the threshold voltages of transistors.

One PUF application is extracting and regenerating *Secret Keys*. The raw data read from the PUF cells are the source for creating *Secret Keys*. The keys are normally extracted during manufacturing in a stable nominal temperature and voltage condition and regenerated in the field with various environmental conditions. Since no error can be tolerated for *Secret Key* application, there are some requisite techniques for identifying unstable PUF cells during manufacturing, such as Multiple Evaluation [6], TMV [6–9], Dark-bit masking [7–10], etc. Then, the stable cells can be used to derive the *Secret Key* for authentication. Whilst the keys are regenerated in the field, SRAM PUF is sensitive to environmental changes and ambient noise, which both cause bit errors. However, time is too limited to execute any aforementioned techniques to reduce errors. Thus, bit error correction techniques are vital, such as BCH codes [11], Hamming codes [12], etc. For error correction and its required PUF entropy, the hardware overheads increase exponentially with the growth in error numbers [13]. This is unsuitable for lightweight IoT applications.

1.2. Main Contributions

Two fundamental circuits are built in this work: one is the custom 8T PUF in Figure 1a; the other is the two-phase power gating cell in Figure 1b. A single-phase power gating cell is also implemented for comparison, as shown in Figure 1c. Based on these circuits, a power-gated 8-Transistor (8T) PUF architecture presented in Figure 2 is developed to alleviate some of the aforementioned issues. This power gating structure can also be utilised with the other bistable PUFs to improve PUF performance. The main contributions are listed below:

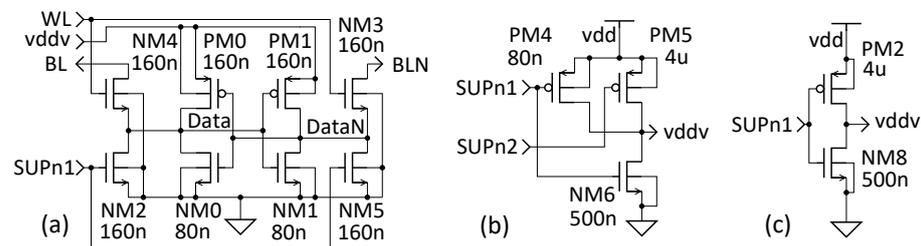


Figure 1. Schematics of (a) 8T PUF cell, (b) two-phase power gating cell and (c) single-phase power gating cell.

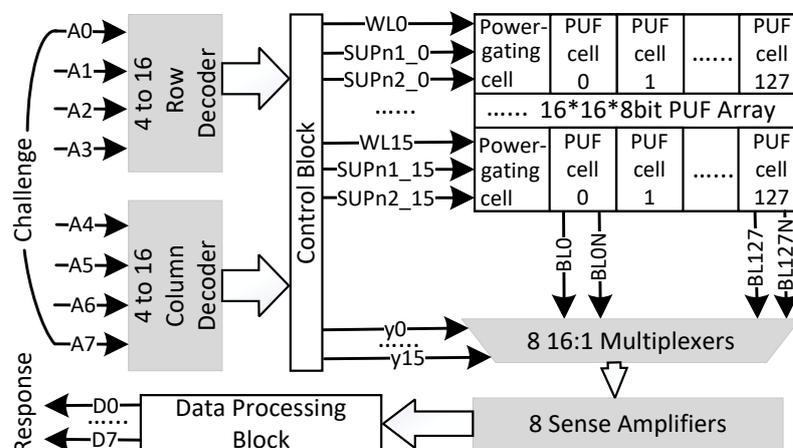


Figure 2. A two kilo-bit power-gated PUF architecture.

- **The custom 8T PUF facilitates fast statistical measurements and improves security:** The 8T PUF maximises physical mismatch and eliminates data retention swiftly for

high-speed evaluations and countering security attacks [14–18]. It does not require a special process for high-density SRAM manufacturing and can be fabricated in the same process as microcontrollers (MCUs). Fast statistical measurements can be performed on this platform to extract raw PUF bits. These raw bits are then processed and marked onto a bitmap to identify the PUF cell instability. These unstable bits, which are discarded in the references, can be used for *True Random Number* generation or as part of the *PUF Response* [19–22].

- **Two-phase power gating improves PUF performance and security whilst saving energy and delaying ageing:** This newly developed power-switching process includes three stages:
 1. Reset stage: The reset stage drops the virtual power supply (vddv), quickly drains the remaining current and eliminates retained data.
 2. Phase I power-up stage: Phase I slowly powers up the chosen PUFs to prolong the metastability-resolving process in the hope of reducing EMI and crosstalk amongst PUF cells [23].
 3. Phase II power-up stage: Finally, phase II speeds up the voltage ramp-up process.

In addition, different combinations of power gating parameters can curb the in-rush current, thus shielding it from side channel attacks, e.g., Differential Power Analysis (DPA) [14,15]. Moreover, the 8T PUF cells are partitioned into rows, and only the chosen rows will be powered up during the reading process. Besides saving energy proportionally, PUF cells without a power supply cannot be read out and are protected from security attacks. Furthermore, power gating can delay the transistor ageing effect.
- **A new data processing method marks out most errors:** The high-speed measurements not only reduce the time needed for the enrolment phase in manufacturing but also enable multiple evaluations in the key regeneration phase. Since the positions of the unstable PUF cells drift away in different voltage or temperature corners, and extreme corners cause more unstable PUF cells, TMV plus Dark-bit masking in nominal conditions [7–9] is insufficient. We propose a new data processing method here. First, during manufacturing, Multiple TMV under nominal conditions and extreme corners are used to mark out unstable readings and flipped readings as Dark bits onto a bitmap. Then, during field usage, the Dark bits are discarded first and followed by the regeneration of Secret Keys using fast TMV. Finally, with a significantly reduced number of error bits remaining, an error correction technique with a lower hardware penalty will be applied. When the BER drop close to 0%, the hardware penalty for error correction can be reduced significantly. The stability levels of PUFs are also recorded on the bitmap. These unstable readings can be used further for True Random Number generation or as part of the *PUF Response*.

1.3. State of the Art and Related PUF Works

Before starting a new silicon-based PUF design, the implementation method needs to be considered first. The silicon-based PUFs can be implemented in an ASIC or a Field-Programmable Gate Array (FPGA). Since FPGAs have ready-to-use resources and can be purchased off-the-shelf, many early PUF works are based on FPGAs [2–4]. However, ASIC-based PUFs can offer far more energy and have a cost-efficient design, which is crucial for lightweight IoT devices [8,9,24–34]. However, FPGA-based PUFs [35–41] are indispensable parts for applications which are not area or energy sensitive but time sensitive instead. In addition, FPGA is a feasible platform for assessing PUF design methods and performance. There are several novel FPGA-based bistable PUF works [36–41], which we intend to implement in our test chip with the power gating method to compare the metastable behaviour and PUF performance.

There are several PUF applications [42–45] that previously applied the power gating method, but their purposes and implementations are different to this approach. Maes et al. [42] implemented power-gated PUF blocks for further investigation. Xu et al. [43] utilised the random duration of multiple power gating to replace voltage control so as

to induce failure patterns to determine the Data Retention Voltage (DRV). In comparison, our design diminishes data retention swiftly for high-speed evaluation. Although with a different purpose, this research and its predecessor [46] have clarified that “a strong DRV fingerprint is correlated with power-up tendency”. This substantiates that the stability or instability degrees of PUF cells come from their innate physical mismatch. However, it can be seen that 40 μs is not enough to eliminate retention data at 25 °C. In 2020, a 2D power gating scheme to relieve an enhancement–enhancement (EE) SRAM PUF from short-circuit currents and also to protect PUF data from attacks was presented by Liu et al. [44]. However, there was no consideration of the major power gating parameters, such as the *SLEEP* transistor design [47], power distribution network [48], etc. This scheme also has a half-selected cell problem, which requires additional peripheral circuits to lower the extra energy consumption.

1.4. Paper Structure

The remainder of this paper is organised as follows. Section 2 describes the power-gated PUF architecture, the 8T PUF design, the two-phase power gating implementation and a new data processing method combining Multiple TMV and Dark-bit masking. Section 3 firstly analyses the power gating parameters, then compares the power-gated 8T PUF behaviour with the 6T PUF, and lastly discusses the importance of a thorough reset for PUF applications. Section 4 measures the processed data for a PUF property comparison. Finally, Section 5 concludes the main contributions and outlines the ongoing work and future plan.

2. Power-Gated PUF Architecture and Design Methods

An example architecture of a power-gated PUF array is illustrated in Figure 2. This architecture consists of some general SRAM function blocks in grey, a 2 kilo-bit power-gated PUF array and two functional blocks, i.e., a Control block and a Data Processing block. The two function blocks will be implemented and evaluated in due course.

- **Challenge–Response pair (CRP):** The 8-bit address inputs are the PUF *Challenges*, and the 8-bit corresponding data outputs are the PUF *Responses*. Together, they form a Challenge–Response pair (CRP).
- **PUF array:** In the PUF array, there are 128 PUFs in a row gated by a power gating cell. Since the main purpose of power gating is facilitating fast evaluations by switching the power supply, the general term *SLEEP* for normal power gating is replaced by *SUPPLY* in this work.
- **Control block:** The switching activity is controlled by the Control block. Apart from passing the decoded higher 4-bit address to choose an 8-bit word from 16:1 multiplexers, it generates *SUPn1*, *SUPn2* and *WL* (Word Line) signals from the decoded lower 4-bit address. Once a PUF row is chosen by the lower 4-bit address, the *SUPn1* signal is discharged to ‘0’ to switch on the power supply. After *Data* are settled down, the *SUPn2* will be discharged to ‘0’, and *WL* will be asserted to ‘1’ in sequence for a reading process. Then, in the Reset stage, both *SUPn1* and *SUPn2* signals are asserted to ‘1’. There will be no power supply to the PUF cells, and the data will be discharged to ‘0’. The protocol of the two-phase power gating method is shown in Figure 3.

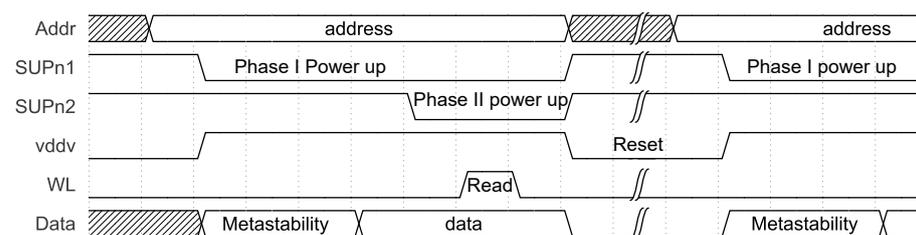


Figure 3. Protocol of two-phase power gating.

- **Data Processing block:** The Data Processing block is there to evaluate the raw read-out bits and marking out different stability levels of each PUF cell. Using TMV under nominal or various voltage and temperature conditions is to sift the unstable or flipped bits to achieve a lower Bit Error Rate (BER).

The input vectors and the output vectors of multiple PUF arrays can be concatenated to match the PUF entropy requirement. The parameters of this architecture, e.g., the type of PUF cell, the number of rows or columns, word width, etc., can be altered for various purposes or implemented in different fabrication techniques. It is worth noting that with larger PUF array dimensions, the sizes of *SUPPLY* transistors need to be evaluated. The driven ability of *SUPn1* and *SUPn2* signals needs to be improved.

The design is implemented with UMC 65 nm technology. The layout of one test circuit is partly illustrated in Figure 4. It includes rows of two-phase power-gated 8T PUFs on the left-hand side and rows of single-phase power-gated 6T-SRAM PUFs on the right-hand side for comparison. Each row has 128 PUF cells.

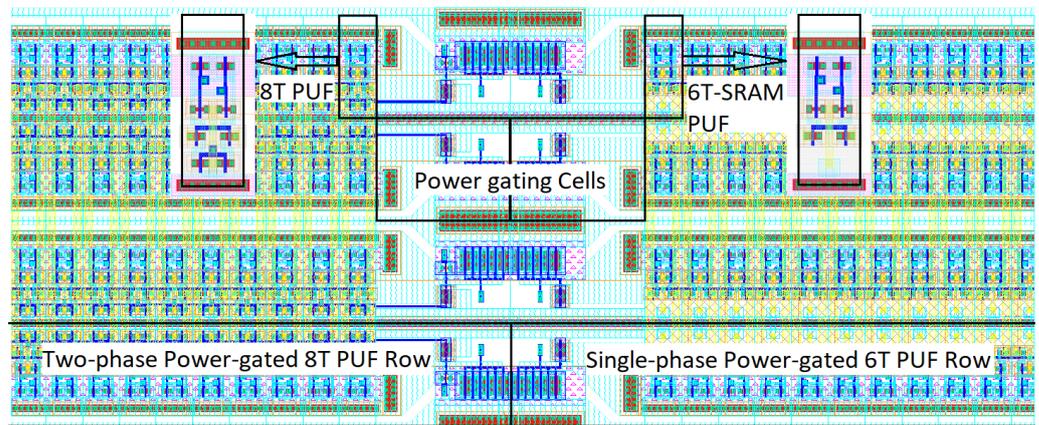


Figure 4. Part of the power-gated PUF test circuit layout.

2.1. 8-Transistor PUF Design

A distinctive feature of the 8T PUF is preparing a clean status for the next power-on cycle with a high reset speed. This clean status means attackers cannot read out any retention data. Unlike some 8T SRAM designs [49,50] concentrating on improving the write margin and dealing with the opposite requirements for read stability and write ability, this 8T PUF abandons the write function and keeps only the SRAM read ability. The physical parameter mismatch, which affects SRAM stability, is actually a vital feature of PUFs. Since mismatch amongst transistors increases with the scaling down the physical size, the 8T PUF uses the smallest transistor to maximise the mismatch. An 8T PUF can be fabricated by the same process as MCUs, so a special process for high-density SRAM is not needed.

The 8T PUF cell shown in Figure 1a stems from conventional 6T-SRAM [51]. Four transistors, namely *PM0*, *PM1*, *NM0* and *NM1*, make up two cross-coupled inverters for storing data. The unique PUF readings predominantly come from the threshold voltage mismatch among these transistors. After powering up, the mutual input and output nodes, i.e., *Data* and *DataN*, both reach a metastable state [52], at which time the outputs of two cross-coupled inverters are lingering between logical '0' and '1', then cross it and settle in one of these stable states. *NM4* and *NM3* are access transistors. They connect the internal nodes *Data* and *DataN* to bit lines *BL* and *BLN*. In addition, two NMOS transistors *NM2* and *NM5* are designed to rapidly discharge the *Data* and *DataN* nodes during the reset stage with both *SUPn1* and *SUPn2* signals asserting. As a result, the reset time can be shortened by 5 orders of magnitude, i.e., from 120 microseconds to 1 nanosecond, hence facilitating high-rate PUF evaluation. Furthermore, no data remanence can be exploited by attackers [53,54]. The PUF layouts were implemented symmetrically. During the evaluations, it is shown that a tiny bias of layout influences the proportion of the ones and

zeros of the PUF readings. With two additional transistors, the area cost of the 8T PUF is 13.2% more than its 6T-SRAM PUF counterpart, which was implemented for comparison as shown in Figure 4.

2.2. Two-Phase Power Gating Method and Design

Originally, the power gating technique was introduced to reduce leakage, which is independent of the transistor switching activity and thus lessens power dissipation. This design eliminates the leakage current from unchosen PUF rows, so the energy consumption is roughly proportional to the chosen PUF percentage. In the two-phase power gating method, the power gating cells switch the chosen rows of PUF cells on and off in three stages: a reset stage, phase I power-up and phase II power-up. First and foremost, this enables fast statistical evaluations, which improve PUF performance. Secondly, this method facilitates different combinations of power gating parameters in hardware implementation, so that the in-rush current can be curbed by adjusting these parameters. The flattened currents provide a method against side channel attacks. Moreover, PUF rows without power supply cannot be read out and are protected from security attacks.

Conventional power gating is conducted via a *SLEEP* transistor to enable a power or ground connection [47]. Figure 1b shows the single-phase power gating cell is operated via a PMOS transistor *PM2* as a *SUPPLY* transistor to switch power on/off and an NMOS transistor *NM8* for fast resetting. As illustrated in Figure 1c, the two-phase power gating cell employs two PMOS transistors as *SUPPLY* transistors for two-stage power switching and an NMOS transistor *NM6* for resetting. In the test circuits, each power gating cell controls the power supply of a row of 128 PUF cells. The reset transistor *NM6* quickly drains the remaining current and drops *v_{ddv}* to 0 V. Since a smaller transistor curbs the drain current and the *v_{ddv}* output, the two-phase power gating exploits this to generate a gentle voltage incline with a smaller *SUPPLY* transistor *PM4* in phase I and create a steep slope of *v_{ddv}* using a larger *SUPPLY* transistor *PM5* in phase II. Consequently, the metastability resolving time in phase I is lengthened in the hope of minimum mutual disturbance amongst PUF cells. Meanwhile, if all PUF cells start metastability in a very short period of time, the in-rush current will be significant due to all the cross-coupled transistors in 8T PUF cells being in saturation mode. The prolonged phase I is able to flatten the current peak as well. Afterwards, phase II takes control and increases *v_{ddv}* swiftly. The sizes of two *SUPPLY* transistors and the time duration of the two phases need to be evaluated to obtain the best trade-off. This will be discussed in the evaluation part in Section 3.1.

2.3. Multiple TMV and Dark-Bit Masking

In some previous methods combining TMV with Dark-bit masking [7–9], TMV cannot eliminate the unstable bits in different environmental conditions, and the Dark-bit mismatch increases bit errors more than 10-fold during field usage. This is because the physical mismatch, e.g., the threshold voltages, vary under different voltages or temperatures. Hence, some stable cells become unstable and vice versa. However, there are some cells that even change their bias tendency and flip their readings. PUF cells with flipped readings cannot be distinguished by TMV under nominal condition.

We introduce a Multiple TMV method with additional TMV thanks to the high-speed evaluation enabled by the proposed design and architecture. During manufacturing, the golden reading is extracted using TMV in the nominal condition first. Then, by comparing the TMV results at all worst corners with the golden reading, any flipped bits and unstable bits are marked as Dark bits. Afterwards, while the *Secret Key* regenerates in the field, the marked-out Dark bits will be eliminated from the raw data first. Then TMV will be executed next, and the unstable bits at this moment are the error bits. Then, these error bits will be corrected with the error correction technique with a lower hardware penalty. The simulation results show that the BER can drop to 0%, which is presented in Section 4.

3. Evaluations and Results Analyses

Due to the sensitivity of bistable PUFs, e.g., the 8T PUF and the SRAM PUF, many aspects of parameter mismatch count for the PUF stability. Although the layouts were implemented symmetrically to avoid human-made bias, the physical parameters still exhibit variations. For this reason, post-layout simulation is a much more accurate way to evaluate the sizes of SUPPLY transistors and measure PUF performance, such as uniqueness, robustness and randomness. For various evaluation purposes, different simulation methods were executed with corresponding test circuits.

3.1. Power Gating Parameter Evaluation

A DC sweep and post-layout transient simulations were performed to evaluate the size of SUPPLY transistors for a cluster of 128 PUF cells and the corresponding behaviour of the PUF cells.

For two-phase power gating, the phase one SUPPLY transistor PM4 must be small enough to prolong the metastability process. Meanwhile, the phase two SUPPLY transistor PM5 should be large enough to supply power to 128 PUF cells quickly. As presented in Figure 5, there is a roughly linear relationship between the SUPPLY transistor width and the approximate value of the vddv plateau while metastability resolves. There is also a nearly quadratic dependence between the SUPPLY transistor width and the time duration for vddv to reach 1.2 V or the metastability resolving time duration. Hence, by varying the size of the SUPPLY transistor, the metastability resolving time can be manipulated.

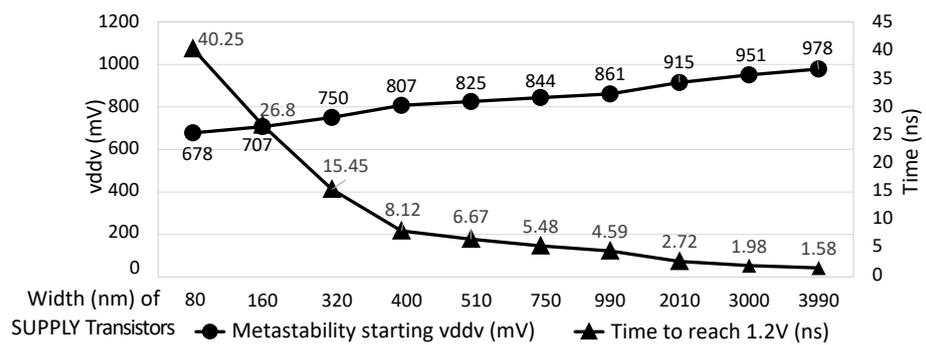


Figure 5. SUPPLY transistor sizes vs. metastability starting vddv and time to reach 1.2 V.

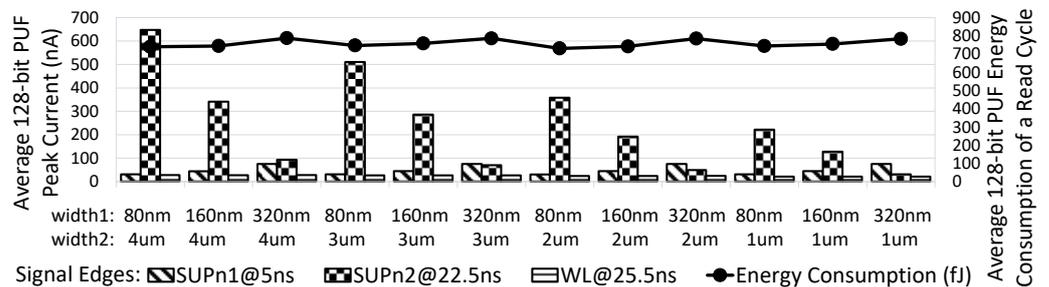


Figure 6. SUPPLY transistor sizes vs. currents and energy consumption.

Phase I SUPPLY transistor widths of 80 nm to 320 nm and phase II SUPPLY transistor widths of 1 μm to 4 μm were picked to evaluate various combinations of two-phase power gating. Figure 6 illustrates the comparison of the current peaks of different signal edges in different combinations of SUPPLY transistors whilst power gating a 128-bit PUF. It can be seen that a higher power-up current corresponds to a larger SUPPLY transistor. The phase II current also relates to the vddv value at the start point of phase II. The energy consumption of a PUF cell for each reading cycle exhibits minimal variation, i.e., from 5.7 fJ to 6.2 fJ. These adjustable power-up currents can limit the in-rush current and be exploited further as a candidate against side channel attacks. In our design, the SUPPLY transistor combination

of 80 nm for PM4 and 4 μm for PM5 is implemented to emphasise the prolonged PUF metastability-resolving process and the rapid reset process.

The power consumption of a power-gated PUF array is proportional to the chosen PUF percentage, in addition to the power consumption of the power-gating cells. The post-layout simulation results show that the power consumption is about 36.5% when only a quarter of PUF rows are enabled. For example, a row of 128 8T PUFs consumes 954 fJ in an array of four PUF rows, which consume 2.613 pJ under the same conditions. The dynamic power of one two-phase power gating cell is about 812.5 aJ.

3.2. Power-Gated PUF Behaviour

To examine the PUF behaviour of 8T PUFs, post-layout Monte Carlo simulations with a Gaussian distribution of transistor threshold voltages were carried out under nominal conditions, i.e., a supply voltage of 1.2 V, an ambient temperature of 27 °C and a *Typical–Typical (TT)* process corner. The test circuit includes a row of 128 two-phase power-gated 8T PUFs and a row of 128 single-phase power-gated 6T-SRAM PUFs on the right-hand side for comparison.

The 128-run Monte Carlo simulation results illustrate the different power-up and reset behaviours of PUF cells, as shown in Figure 7. In Figure 7a, with a 4 μm width SUPPLY transistor, the vddv of single-phase power gating reaches 1.2 V in roughly 1 ns, whilst quickly resolving the metastability of inside node pairs, i.e., Data and DataN. In contrast, Figure 7b illustrates that vddv reaches around 0.7 V in 5 ns and lingers for about 4 ns, then gradually climbs up towards 1.2 V at phase I power-up with the SUPPLY transistor width of 80 nm. Following this, with a 4 μm width SUPPLY transistor, phase II starts from 22.5 ns and swiftly increases vddv to 1.2 V in around 1 ns with the help of the 80 nm SUPPLY transistor. It can be seen that the low start-up voltages in phase I lengthen the metastability resolving time of 8T PUFs. Simultaneously, Data and DataN start wrestling while vddv is ramping up slowly, then escape out of metastability and tend to their distinct logical status in various resolving times due to their intrinsically varied physical parameters. These opposite tendencies of Data and DataN resemble the random PUF behaviour in real circuits.

After powering up, the WL signal asserts at 25 ns for data reading. Finally, in the reset stage, PUF cells are powered down. The remaining currents are drained away quickly from the 8T PUF cells within 1 ns. However, the traditional 6T-SRAM PUFs still have data retention, which not only affects the initial states of Data and DataN pairs in the following cycles but also can be targeted by attackers. It is worth noting that with the technology scaling down from 90 nm to 65 nm, the reset period for the same 6T-SRAM design lengthens from 5 μs to 120 μs, as listed in Table 1.

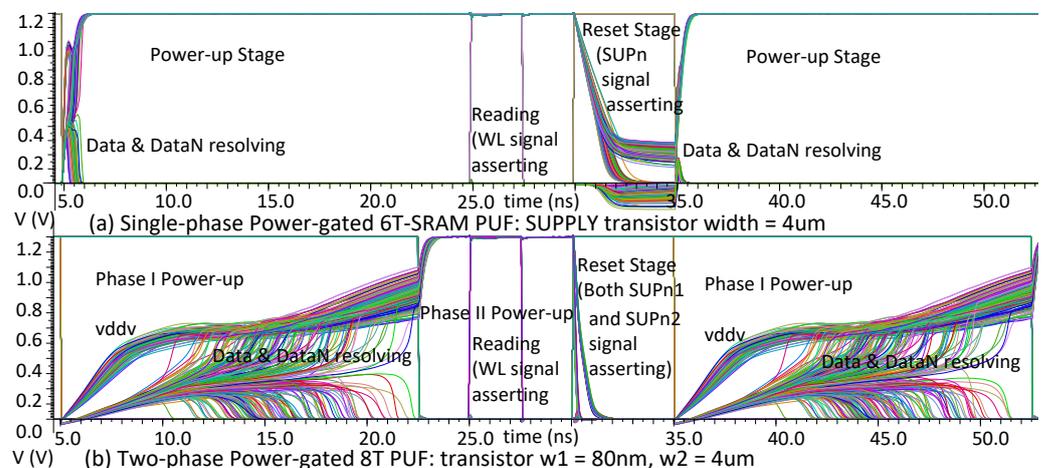


Figure 7. Waveforms of a 128-run post-layout Monte Carlo simulation.

Table 1. Reset periods in comparison.

	HOST 2012 [6]	DFT 2012 [55]	ESSCIRC 2014 [10]	DFT 2016 [20]	DATE 2023 [56]	This Work	This Work
Design	6T-SRAM	6T-SRAM	Hybrid	DFF	6T-SRAM	6T-SRAM	8T PUF
Technology	65 nm	45 nm	22 nm	45 nm	90 nm	65 nm	65 nm
Reset Time	1 ms	1 s	1 s	50 μ s	5 μ s	120 μ s	1 ns

3.3. Reset Period and Effect

Without resetting thoroughly, the assessment of the PUF characteristics can be misleading. One piece of evidence is that with the data remanence, the under-reset 6T-SRAM PUF readings do not change in more than 200,000 runs of a Monte Carlo simulation compared to the 8T PUF. This creates a deceptive deduction that the 6T-SRAM PUF cells are 100% stable in nominal conditions [57]. However, if the reset time duration is prolonged to 120 μ s with 65 nm technology, the simulation results show that the stability of the thoroughly reset 6T-SRAM PUF will be affected in the same way as the 8T PUF, as illustrated in Figure 8. Long duration transient simulations were conducted with intensified noise. In these simulations, the 8T PUF cell, with the help of two discharging NMOS transistors, can reset two internal data nodes to the absolute values below 10 mV within 1 ns, while the 6T-SRAM PUF can barely achieve the same reset level even in 100 μ s.

This reset issue has been overlooked because most research does not switch power at such high frequencies, as listed in Table 1. Consequently, the longer power-off periods result in a better reset quality under the same conditions.

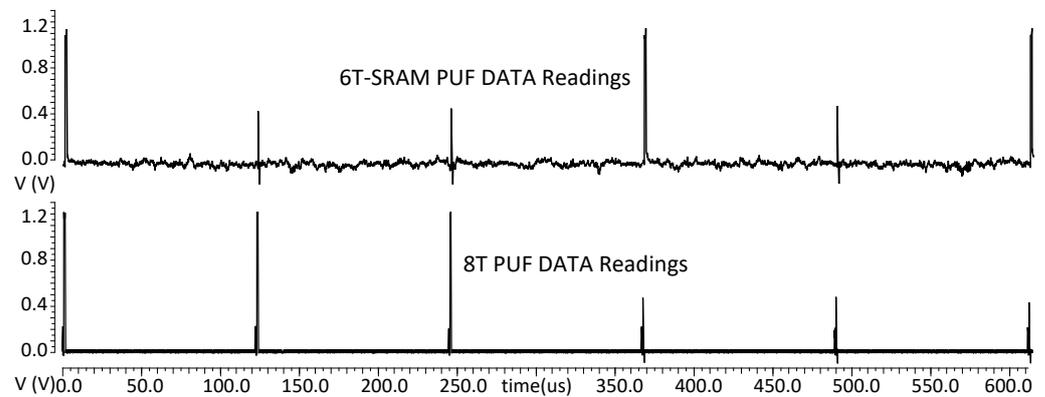


Figure 8. Transient noise simulation with a 120 μ s reset period.

4. 8-Transistor PUF Performance Measurements

To assess the PUF performance, 5G transient noise, including thermal noise, flicker noise, shot noise, etc., was added to the post-layout Monte Carlo simulations because all 8T PUF cells and 6T-SRAM PUF cells were stable in nominal conditions without any external noise. For the two-phase power gating SUPPLY transistors, an 80 nm width was chosen for phase I, and 4 μ m was chosen for phase II. Meanwhile, the 6T-SRAM PUF was examined under the same conditions for comparison. In addition to nominal conditions, the design was evaluated under different PVT corners, i.e., temperatures from -40 $^{\circ}$ C to 85 $^{\circ}$ C and supply voltages from 0.8 V to 1.6 V, and the process corners included *Slow–Slow (SS)*, *Fast–Fast (FF)* etc. For quantitative evaluation, clusters of 2048 PUF cells were read out for 100 power on–off cycles of 60 ns. This resembles 100 times the same *Challenge*. The 100 2048-bit read-out strings are the PUF *Responses*. Furthermore, the PUF clusters imitate different PUF devices. With a thorough reset, these PUF characteristics are able to be assessed with post-layout simulations. The intensive multiple simulations here aim to examine the PUF properties to prepare for prototype chip fabrication, so the extracted data were then processed with Matlab to acquire PUF robustness, uniqueness, randomness, etc.

The results were compared with some prior work in Table 2. It is worth noting that the works [8,9,25,58–61] are ASIC-based measurements, and only [62] and this work include simulation-based evaluations. In fact, post-layout simulations are very resource and time consuming, so only a few works, e.g., [62], can extract data from schematic-based simulations.

Table 2. PUF performance comparison.

Measurement	CMOS			Chip-Based			Simulation-Based		
	ISSCC 2014 [8]	ISSCC 2015 [58]	ISSCC 2016 [59]	JSSC 2017 [9]	JSSC 2020 [44]	JSSC 2022 [60]	Electronics 2023 [61]	IJCTA 2017 [62]	This Work
Technology	22 nm	65 nm	45 nm	14 nm	130 nm	180 nm	130 nm	65 nm	65 nm
Bitcell Area (μm^2)	4.66	25.35	5.3	1.84	6.3	223	72.03	-	4.19
Bitcell Area (F^2)	9632	6000	2613	9388	373	7222	4262.13	-	992
Unstable Bits	30%	1.73%	-	26.37%	2.14%	0.61%	0.586%	0.32% Δ	7.71%
Native BER	-	-	0.1%	5.76%	0.21%	0.13%	0.49%	-	1.17%
Worst BER	6%	4.56% *	2.84%	6.78%	0.34%	1.1%	3.125%	-	12.21%
Stabilised BER	0.97%	1.73%	0.21%	1.46%	0%	0.13%	-	-	0%
Intra-distance	2.58%	0.92%	-	3.4%	0.3%	0.16%	0.491%	2.25% Δ	1.45%
Inter-distance	49%	50.14%	49.8%	48.6%	49.23%	49.3%	50.12%	-	50.67%
Mask Ratio	11%	-	18.5%	20%	31.2–75%	0.61%	0%	-	21.78%
ACF @95% c.1 \dagger	0.01	0.0363	0.017	-	0.0228	0.0472	0.025	-	0.0315
Energy/bit (fj)	13	15	-	4	128	-	5.36	-	6.15

* The worst unstable bit rate. \dagger Auto-correlation Function at 95% confidence level. Δ $M_{VOUT} = 1$ mV.

By comparing this work with state-of-the-art designs, it is shown to satisfy the basic PUF requirements and can be evaluated as a PUF. From Table 2, it can be seen that some of the latest designs [44,59–61] improve the native stability of PUFs and achieve the native BER very close to 0%. Ref. [44] realises a 0% stabilised BER with relatively high mask ratios and energy consumption per bit, while [59–61] utilises a novel circuit implementation. This work tries to reach the same level of performance with high-frequency evaluations enabled by the power-gated 8T PUF structure. Since the post-layout simulation is extremely time and resource consuming, there are only a few previous works that attain the same level of results by simulation, e.g., the results of [62] are extracted from schematic-based simulations. The main achievement of the proposed design and structure is the speed, which saves time and hence conserves the total power consumption, etc.

4.1. Uniformity

For PUF uniformity, the proportion of ones and zeros of the 8T PUF readings should be evenly distributed. Our results of stable zeroes are 46.22% and stable ones are 46.06%.

4.2. Robustness: Intra-Distance and BER

For an ideal PUF cell, the Responses to the same Challenge should always be the same. However, the existence of unstable cells makes this impossible. Thus, the intra-distance, which is the Hamming distance between bit strings of the Responses from repeated measurements of the same PUF, is used to evaluate the PUF robustness. To minimise the area overheads for error correction in the key regeneration phase, the intra-distance is expected to be close to zero. By comparing every pair of two Response bit strings and adding up the total numbers of different bits, the intra-distances are calculated. The average intra-distance of all the simulated 8T PUF groups is 1.45%.

Under nominal conditions, the average percentage of total unstable bits in 20 groups of 2048 8T PUF cells is 7.71%; the raw BER is 1.17%. The corresponding BER of the worst corner is 12.21%. Most of the unstable readings can be fixed by the Data Processing block using TMV, which can be visualised in the golden bitmap generated in the key enrolment phase. For example, in the golden bitmap shown in Figure 9, squares with different colours represent the statistical results of the PUF stability, i.e., white indicates stable ‘0’, black squares show stable ‘1’ and different grey shades present the degree of bias of the cells. Most impressively, the worst corner BER can be decreased to 0% with Dark-bit

masking coordinated with TMV in the worst voltage and temperature corners and in the field. The mask ratio is 21.78%. As a result, the hardware penalty for error correction is minimised.

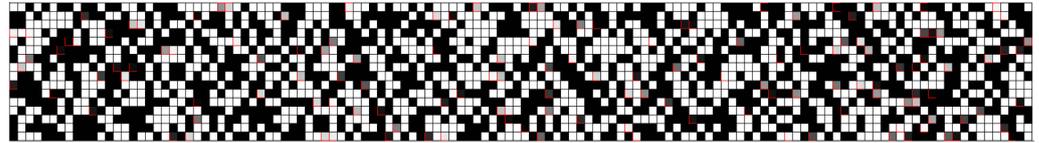


Figure 9. Golden bitmap of 2048 8T PUFs.

4.3. Uniqueness: Inter-Distance

Since PUFs are the hardware sources for identifying individual devices or generating *Secret Keys*, the extracted information from each PUF should be unique. The fractional Hamming distance was computed between different PUF devices. In order to successfully identify each PUF device, this value should be close to 50%. In the experiments, the *Response* bit strings of different PUF clusters were compared to calculate the inter-distance. The average inter-distance of the 20 8T PUF clusters is 50.67%.

4.4. Randomness

Firstly, the Applied Auto-correlation Function (ACF) was used to analyse the spatial correlations among 64 groups of 2048 8T PUFs. The result is 0.0315, which is close to zero, within the 95% confidence bound. This indicates a low spatial correlation of the 8T PUFs based on the physical extraction from the layout design. Then, the randomness of 64 groups of 2048 8T PUFs was assessed with the NIST SP 800-22 [63] statistical test suite. The NIST results are listed in Table 3. The simulation results passed most of the tests. However, the randomness generated by the simulator has limitations, especially for the large amount of data which is required by both the ACF and NIST test suite. Moreover, this practical simulation method is time and resource consuming. A test chip for facilitating on-chip statistical experiments is under development to improve the evaluation efficiency and verify the PUF behaviour in the real world.

Table 3. NIST analysis for 64 PUF groups.

	String Length	<i>p</i> -Value	Proportion
Frequency	2048	0.73856	64/64
Block Frequency	2048	0.91013	64/64
Cumulative Sums	2048	0.87698	64/64
Runs	2048	0.0204	11/64
Longest Run	2048	0.005	14/64
Rank	2048	0.44962	64/64
FFT	2048	0.18139	52/64
Approx. Entropy	2048	0.009	28/64
Serial	2048	0.18290	40/64
Linear Complexity	2048	0.63316	61/64

5. Conclusions and Future Work

In this paper, the design, architecture and evaluation of a rapid reset 8T PUF utilising the power gating technique are presented. Its purpose is to enhance PUF stability and minimise the hardware penalty for error corrections. The design can apply on–off power cycles repeatedly to PUF clusters to facilitate fast multiple evaluations for extracting the bias probability of PUF cells. The ultra-fast speed enables TMV in different voltage and temperature corners or in the field. Dark-bit masking based on these extracted data can reduce the BER to close to zero in experiments. In the design, an SRAM-based 8T PUF cell with the ability to eliminate data retention is built, and a two-phase power gating method is devised and evaluated. Besides switching the power supply swiftly and saving

power, the power-up process can be manipulated via varying the power gating method and parameters so as to decrease the interference between sensitive PUF cells and limit the in-rush current during power-up. Consequently, PUF stability and security are enhanced. In addition, the clean and fast reset makes it possible for swift and accurate PUF measurements either in simulation or in a fabricated silicon chip. The 8T PUF characteristics, including robustness, uniqueness and randomness, are thus qualitatively confirmed.

Future work will include quantitative evaluations of the 8T PUF performance on various PVT corners and assessments of different power gating settings. The extracted unstable degrees of PUF cells will be post-processed and added to the PUF *Response* to increase its entropy. In addition to this, a test chip is currently being fabricated. Afterwards, the fabricated test chip will enable on-chip statistical experiments, which serve as a platform to extract analogue secrets. Finally, the differences in the power-gated 8T PUF's entropy compared to its 6T-SRAM PUF counterpart will be assessed, measured and reported.

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Abbreviations

The following abbreviations are used in this manuscript:

8T	8-Transistor
PUF	Physically Unclonable Function
6T	6-Transistor
SRAM	Static Random-Access Memory
TMV	Temporal Majority Voting
PVT	Process, Voltage and Temperature
BER	Bit Error Rate
ASIC	Application Specific Integrated Circuit
IoT	Internet of Things
IC	Integrated Circuit
MCU	Microcontroller
vddv	virtual power supply
DPA	Differential Power Analysis
FPGA	Field-Programmable Gate Array
CRP	Challenge–Response Pair
DRV	Data Retention Voltage
WL	Word Line
TT	Typical–Typical
SS	Slow–Slow
FF	Fast–Fast
CMOS	Complementary Metal Oxide Semiconductor
ACF	Applied Autocorrelation Function

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