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Iterative Algorithms for Split Common Fixed Point Problem Involved in Pseudo-Contractive Operators without Lipschitz Assumption

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Received: 23 July 2019; Accepted: 20 August 2019; Published: 23 August 2019



Abstract: Two iterative algorithms are suggested for approximating a solution of the split common fixed point problem involved in pseudo-contractive operators without Lipschitz assumption. We prove that the sequence generated by the first algorithm converges weakly to a solution of the split common fixed point problem and the second one converges strongly. Moreover, the sequence $\{x_n\}$ generated by Algorithm 3 strongly converges to $z = \text{proj}_S 0$, which is the minimum-norm solution of problem (1). Numerical examples are included.

Keywords: split common fixed point problem; iterative algorithms; pseudo-contractive operators; Lipschitz assumption

MSC: 47H10; 49M37; 49K35; 90C25

1. Introduction

The split common fixed point problem was investigated in 2009 by Censor Y. and Segal A. [1]. Further research on this problem discussed in works by the authors of [2–12]. More specifically, given two Hilbert space $\mathcal{H}_1$ and $\mathcal{H}_2$, nonlinear operators $U : \mathcal{H}_1 \rightarrow \mathcal{H}_1$ and $T : \mathcal{H}_2 \rightarrow \mathcal{H}_2$ and a bounded linear operator $A : \mathcal{H}_1 \rightarrow \mathcal{H}_2$. Let $x \in \mathcal{H}_1$ be a solution of split common fixed point problem satisfying

$$x \in F(U) \quad \text{and} \quad Ax \in F(T) \quad (1)$$

where $F(U)$ and $F(T)$ mean the fixed point sets. If U and T are both metric projects, problem (1) is actually problem (2) [13,14], and further development of this topic made by [15–19]. To be more specific, given two nonempty closed convex sets $C \subset \mathcal{H}_1$ and $Q \subset \mathcal{H}_2$ and A is above mentioned. Let $x \in \mathcal{H}_1$ be a solution of split feasibility problem satisfying

$$x \in C \quad \text{and} \quad Ax \in Q, \quad (2)$$

These two problems ((1) and (2)) have received much attention, and have been extensively investigated due to applications in signal processing, image reconstruction, [14], and intensity modulated radiation therapy [20]. Recently, Yen L. et al. [21] learn the problem (2) and applying it to

a model in electricity production, they successfully established a Nash–Cournot equilibrium model with minimal environmental cost. Wang J. et al. [22] study the linear convergence of CQ algorithm for solving the problem (2) and investigate an application in gene regulatory network inference.

For solving the problem (1), Censor Y. and Segal A. [1] suggested the following scheme.

$$x_{n+1} = U(x_n - \tau A^*(I - T)Ax_n), \quad (3)$$

where τ is a fixed stepsize and A^* is the adjoint operator of A . Algorithm (3) was originally designed to solved problem (1) for directed operators. Noting that if the stepsize τ is chosen in $(0, 2/\|A\|^2)$, then the iterative sequence $\{x_n\}$ generated by (3) weakly converges to a solution of the problem (1). Subsequently, iterative schemes and these variants [10,23] were explored to the demicontractive operators, quasi-nonexpansive operators and finite many directed operators.

Very recently, Wang F. [23] has been devoting himself to the study of problems (1). Accordingly, he proposed a new method for solving the problems (1) so that the variable stepsize does not need to compute the norm $\|A\|$:

$$x_{n+1} = x_n - \rho_n((I - U)x_n + A^*(I - T)Ax_n), \quad (4)$$

where $\{\rho_n\} \subset (0, \infty)$ is chosen such that

$$\rho_n = \frac{\|(I - U)x_n\|^2 + \|(I - T)Ax_n\|^2}{\|(I - U)x_n + A^*(I - T)Ax_n\|^2}, \quad (5)$$

Wang obtained the weak convergence of algorithm (4).

In this paper, we extend a previous author's results from the demicontractive operators [8,10,24], firmly-nonexpansive operators [25], quasi-nonexpansive operators [26], directed operators [1], nonexpansive operators [27], and strictly pseudo-contractive operators [28] to the more general pseudo-contractive operators. Subsequently, two algorithms are suggested based on (4) and (5) to solve the problem (1). Weak and strong convergence of the proposed algorithms are obtained.

2. Preliminaries

Let $\mathcal{H}$ be a real Hilbert space equipped up its inner product $\langle \cdot, \cdot \rangle$ and norm $\|\cdot\|$ [8]. The notation $x_n \rightharpoonup x$ means weak convergence and $x_n \rightarrow x$ means strong one. The notation $\text{Fix}(T)$ stands for the set of fixed points of the operator T . The symbol $\omega_w(x_n)$ denotes the weak ω -limit set of $\{x_n\}$, that is,

$$\omega_w(x_n) = \{x : x_{n_i} \rightharpoonup x \text{ for some subsequence } \{x_{n_i}\} \text{ of } \{x_n\}\}.$$

Let C be a nonempty closed convex subset of $\mathcal{H}$. Recall that the projection P_C from $\mathcal{H}$ onto C defined by

$$\|x - P_C x\| = \min\{\|x - y\| : y \in C, x \in \mathcal{H}\}.$$

Proposition 1 ([10]). *Given $x \in \mathcal{H}$ and $z \in C$.*

- (1) $z = P_C x \Leftrightarrow \langle x - z, y - z \rangle \leq 0$, for all $y \in C$.
- (2) $z = P_C x \Leftrightarrow \|x - z\|^2 \leq \|x - y\|^2 - \|y - z\|^2$, for all $y \in C$.
- (3) $\langle x - y, P_C x - P_C y \rangle \geq \|P_C x - P_C y\|^2$, for all $y \in \mathcal{H}$, which hence implies that P_C is nonexpansive.

Definition 1 ([4]). *Let $T : \mathcal{H} \rightarrow \mathcal{H}$ be a nonlinear operator.*

- *T is called nonexpansive if*

$$\|Tx - Ty\| \leq \|x - y\|, \quad \forall x, y \in \mathcal{H};$$

- T is called *firmly nonexpansive* if

$$\|Tx - Ty\|^2 \leq \|x - y\|^2 - \|(I - T)x - (I - T)y\|^2, \quad \forall x, y \in \mathcal{H}$$

or equivalently

$$\|Tx - Ty\|^2 \leq \langle Tx - Ty, x - y \rangle, \quad \forall x, y \in \mathcal{H}.$$

Also, the mapping $I - T$ is *firmly nonexpansive*.

- T is called *strictly pseudo-contractive* if there exists $k < 1$ such that

$$\|Tx - Ty\|^2 \leq \|x - y\|^2 + k\|(I - T)x - (I - T)y\|^2, \quad \forall x, y \in \mathcal{H}$$

- L -Lipschitzian if there exists $L > 0$ such that

$$\|Tx - Ty\| \leq L\|x - y\|, \quad \forall x, y \in \mathcal{H};$$

Definition 2 ([24]). Let $T : \mathcal{H} \rightarrow \mathcal{H}$ be a nonlinear operator with $\text{Fix}(T) \neq \emptyset$.

- T is called *demicomtractive* if there exists a constant $k \in [0, 1)$ such that

$$\|Tx - x^*\|^2 \leq \|x - x^*\|^2 + k\|x - Tx\|^2, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T)$$

or equivalently

$$\langle x - Tx, x - x^* \rangle \geq \frac{1-k}{2}\|x - Tx\|^2, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T);$$

- T is called *directed* if

$$\|Tx - x^*\|^2 \leq \|x - x^*\|^2 - \|x - Tx\|^2, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T)$$

which is equivalent to

$$\langle x - Tx, x - x^* \rangle \geq \|x - Tx\|^2, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T);$$

Definition 3 ([4]). Let $T : \mathcal{H} \rightarrow \mathcal{H}$ be a nonlinear operator.

T is called *pseudo-contractive* if

$$\langle Tx - Ty, x - y \rangle \leq \|x - y\|^2, \quad \forall x, y \in \mathcal{H}.$$

It is well known that T is a pseudo-contractive operator if and only if

$$\|Tx - Ty\|^2 \leq \|x - y\|^2 + \|(I - T)x - (I - T)y\|^2, \quad \forall x, y \in \mathcal{H}.$$

Proposition 2 ([29]). Let T be a pseudo-contractive operator with the nonempty fixed point set $\text{Fix}(T)$, then the following conclusion holds.

$$\langle Tx - x, Tx - x^* \rangle \leq \|Tx - x\|^2, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T).$$

Generally speaking, pseudo-contractive operators are also assumed to be L -Lipschitzian with $L > 1$. Next, to overcome the L -Lipschitzian property, the authors of [29] assume that the pseudo-contractive operator T satisfies the following condition.

$$\langle Tx - x, Tx - x^* \rangle \leq 0, \quad \forall (x, x^*) \in \mathcal{H} \times \text{Fix}(T). \quad (6)$$

Definition 4 ([23]). Let $T : \mathcal{H} \rightarrow \mathcal{H}$ be a nonlinear operator with $\text{Fix}(T) \neq \emptyset$. Then, $I - T$ is said to be demiclosed at zero, if, for any $\{x_n\}$ in $\mathcal{H}$, there holds the following implication:

$$\left[\begin{array}{l} x_n \rightharpoonup x \\ (I - T)x_n \rightarrow 0 \end{array} \right] \Rightarrow x \in \text{Fix}(T)$$

The demiclosedness for pseudo-contractive operators in the following will often be used.

Lemma 1 ([29]). Let $\mathcal{H}$ be a real Hilbert space, C a closed convex subset of $\mathcal{H}$. Let $T : C \rightarrow C$ be a continuous pseudo-contractive operator. Then

- (1) $\text{Fix}(T)$ is a closed convex subset of C ,
- (2) $(I - T)$ is demiclosed at zero.

To attain weak convergence result, the following result is useful.

Lemma 2 ([10]). Let $\mathcal{H}$ be a Hilbert space and $\{x_n\}$ be a bounded sequence in $\mathcal{H}$ such that there exists a nonempty closed convex set $C \in \mathcal{H}$ satisfying

- (1) for every $w \in C$, $\lim_{n \rightarrow \infty} \|x_n - w\|$ exists;
- (2) each weak cluster point of the sequence $\{x_n\}$ is in C .

Then $\{x_n\}$ converges weakly to a point in C . More specifically, $x^* = \lim_{n \rightarrow \infty} P_C x_n$.

To attain strong convergence result, we need to use the following lemmas.

Lemma 3 ([8]). Let $\{a_n\}$ be a sequence of nonnegative real numbers satisfying the property

$$a_{n+1} \leq (1 - \gamma_n)a_n + \sigma_n, n \geq 0.$$

where $\{\gamma_n\} \subset (0, 1)$ and $\{\sigma_n\}$ are such that

- (1) $\sum_{n=0}^{\infty} \gamma_n = \infty$;
- (2) either $\limsup_{n \rightarrow \infty} \frac{\sigma_n}{\gamma_n} \leq 0$ or $\sum_{n=0}^{\infty} |\sigma_n| < \infty$.

Then $\{a_n\}$ converges to zero.

Lemma 4 ([4]). Let $\{u_n\}$ be a sequence of real numbers. Assume $\{u_n\}$ does not decrease at infinity, that is, there exists at least a subsequence $\{u_{n_k}\}$ of $\{u_n\}$ such that $u_{n_k} \leq u_{n_k+1}$ for all $k \geq 0$. For every $n \geq N_0$, define an $\{\tau(n)\}$ as

$$\tau(n) = \max\{i \leq n : u_{n_i} < u_{n_i+1}\}.$$

Then $\tau(n) \rightarrow \infty$ as $n \rightarrow \infty$ and for all $n \geq N_0$,

$$\max\{u_{\tau(n)}, u_n\} \leq u_{\tau(n)+1}.$$

In the following two sections, we consider the problem (1) for pseudo-contractive operators without Lipschitz assumption. For problem (1), the standard assumptions are usually the following.

- the problem (1) is consistent, notation S means the solution set;

- both T and U are continuous pseudo-contractive operators without Lipschitz assumption.

3. Weak Convergence Theorem

Next come the iterative scheme for approximating a solution of the problem (1) involved in pseudo-contractive operators without Lipschitz assumption.

Algorithm 1. Initial guess x_0 is arbitrary chosen and assume that x_n has been constructed. If

$$\|(I - U)x_n + A^*(I - T)Ax_n\| = 0, \quad (7)$$

then stop (i.e., x_n solves the problem (1)); otherwise, calculate the next x_{n+1} by the formula [23]:

$$x_{n+1} = x_n - \rho_n((I - U)x_n + A^*(I - T)Ax_n),$$

where the stepsize sequence τ_n is chosen as

$$\rho_n = \frac{\|(I - U)x_n\|^2 + \|(I - T)Ax_n\|^2}{\|(I - U)x_n + A^*(I - T)Ax_n\|^2},$$

We need two lemmas to complete the convergence analysis of our proposed algorithm. The first lemma shows that the proposed algorithm is well defined.

Lemma 5. Assume that (7) holds for $n \geq 0$, then x_n solves the problem (1).

Proof. For any $w \in S$ and (6), we have

$$\begin{aligned} 0 &= \langle (I - U)x_n + A^*(I - T)Ax_n, x_n - w \rangle \\ &= \langle x_n - Ux_n, x_n - Ux_n \rangle + \langle (I - T)Ax_n, Ax_n - TAx_n \rangle \\ &\quad + \langle x_n - Ux_n, Ux_n - w \rangle + \langle (I - T)Ax_n, TAx_n - Aw \rangle \\ &\geq \|x_n - Ux_n\|^2 + \|Ax_n - TAx_n\|^2. \end{aligned}$$

Hence, $x_n = Ux_n$ and $Ax_n = TAx_n$, and the proof is thus complete. $\square$

Lemma 6. Assume that the sequence x_n satisfies

$$\lim_{n \rightarrow \infty} \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} = 0,$$

then it follows that

$$\lim_{n \rightarrow \infty} \|x_n - Ux_n\| = \lim_{n \rightarrow \infty} \|(I - T)Ax_n\| = 0,$$

Proof. By our hypothesis, we have

$$\begin{aligned} \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} &\geq \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{2(\|x_n - Ux_n\|^2 + \|A^*(I - T)Ax_n\|^2)} \\ &\geq \frac{\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2}{2\max(1, \|A\|^2)} \end{aligned}$$

Hence, the desired assertion follows. $\square$

The second lemma analyzes the convergence of the proposed algorithm. Now the weakly convergence of Algorithm 1 presented below.

Theorem 1. Let $\{x_n\}$ be the sequence generated by Algorithm 1. Then, $\{x_n\}$ converges weakly to a solution x^* of problem (1), where $x^* = \lim_{n \rightarrow \infty} P_S x_n$.

Proof. For any $w \in S$, by the expression of y_n , from (6), we obtain

$$\begin{aligned} \langle y_n, x_n - w \rangle &= \langle (I - U)x_n + A^*(I - T)Ax_n, x_n - w \rangle \\ &= \langle x_n - Ux_n, x_n - Ux_n \rangle + \langle (I - T)Ax_n, Ax_n - TAx_n \rangle \\ &\quad + \langle x_n - Ux_n, Ux_n - w \rangle + \langle (I - T)Ax_n, TAx_n - Aw \rangle \\ &\geq \|x_n - Ux_n\|^2 + \|Ax_n - TAx_n\|^2. \end{aligned}$$

Consequently,

$$\begin{aligned} \|x_{n+1} - w\|^2 &= \|x_n - w\|^2 - 2\rho_n \langle y_n, x_n - w \rangle + \rho_n^2 \|y_n\|^2 \\ &\leq \|x_n - w\|^2 - \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} \end{aligned}$$

In particular, $\|x_{n+1} - w\| \leq \|x_n - w\|$, so $\{x_n\}$ is Féjer-monotone w.r.s. S.

Since $\{x_n\}$ is Féjer-monotone, so $\{\|x_n - z\|\}$ is nonincreasing. Hence, $\{x_n\}$ is bounded, and so is the sequence $\{Ax_n\}$. Moreover,

$$\sum_{n=0}^{\infty} \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} < \infty.$$

In particular, we have

$$\lim_{n \rightarrow \infty} \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} = 0.$$

By Lemma 6, this yields $\lim_{n \rightarrow \infty} \|x_n - Ux_n\| = \lim_{n \rightarrow \infty} \|(I - T)Ax_n\| = 0$. From Lemma 1 and Lemma 2, sequence $\{x_n\}$ weakly converges to x^* of problem (1). $\square$

Now, we use the result to solve the problem (2).

Algorithm 2. An initial guess x_0 is arbitrarily chosen and we assume that x_n has been constructed. If

$$\|(I - P_C)x_n + A^*(I - P_Q)Ax_n\| = 0,$$

then stop (i.e., $\{x_n\}$ solves the problem (2)); otherwise, calculate the next x_{n+1} by the formula [23]

$$x_{n+1} = x_n - \rho_n((I - P_C)x_n + A^*(I - P_Q)Ax_n), \quad (8)$$

where the stepsize sequence τ_n is chosen as

$$\rho_n = \frac{\|(I - P_C)x_n\|^2 + \|(I - P_Q)Ax_n\|^2}{\|(I - P_C)x_n + A^*(I - P_Q)Ax_n\|^2},$$

Theorem 2. Let $\{x_n\}$ be the sequence generated by (8). Then, $\{x_n\}$ converges weakly to a solution x^* of problem (2).

4. Strong Convergence Theorem

We proposed a damped algorithm so that the strong convergence is obtained.

Algorithm 3. Initial guess x_0 is arbitrarily chosen and we assume x_n has been constructed. If

$$\|(I - U)x_n + A^*(I - T)Ax_n\| = 0,$$

then stop (i.e., x_n solves the problem (1)); otherwise, calculate the next x_{n+1} by the formula:

$$x_{n+1} = (1 - \delta_n)x_n + \delta_n(1 - \gamma_n)(x_n - \rho_n((I - U)x_n + A^*(I - T)Ax_n)),$$

where the stepsize sequence τ_n is chosen as

$$\rho_n = \frac{\|(I - U)x_n\|^2 + \|(I - T)Ax_n\|^2}{\|(I - U)x_n + A^*(I - T)Ax_n\|^2},$$

Theorem 3. Assume the parameters satisfy the following conditions.

- (i) $\lim_{n \rightarrow \infty} \gamma_n = 0$, $\sum_{n=0}^{\infty} \gamma_n = +\infty$;
- (ii) $0 < \liminf_{n \rightarrow \infty} \delta_n(1 - \gamma_n) \leq \limsup_{n \rightarrow \infty} \delta_n(1 - \gamma_n) < 1$.

Then the sequence $\{x_n\}$ generated by Algorithm 3 strongly converges to $z = \text{proj}_S 0$, which is the minimum-norm solution of problem (1).

Proof. Let $u_n = x_n - \rho_n((I - U)x_n + A^*(I - T)Ax_n)$. Analogously,

$$\|u_n - w\|^2 \leq \|x_n - w\|^2 - \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} \quad (9)$$

By (9), we obtain

$$\begin{aligned} \|x_{n+1} - w\|^2 &\leq (1 - \delta_n)\|x_n - w\|^2 + \delta_n\|(1 - \gamma_n)u_n - w\|^2 \\ &\leq (1 - \delta_n)\|x_n - w\|^2 + \delta_n(1 - \gamma_n)\|x_n - w\|^2 + \gamma_n\delta_n\|w\|^2 \\ &\leq \max\{\|x_n - w\|^2, \|w\|^2\}, \end{aligned} \quad (10)$$

which shows the boundedness of $\{x_n\}$. Returning to (9) and (10), we have

$$\delta_n(1 - \gamma_n) \left(\frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} \right) \leq (1 - \gamma_n\delta_n)\|x_n - w\|^2 + \gamma_n\delta_n\|w\|^2 - \|x_{n+1} - w\|^2. \quad (11)$$

Two possible cases are considered.

Case one. Suppose $m > 0$ and $n \geq m$ such that $\{\|x_n - w\|\}$ is nonincreasing. So, we have the existence of $\lim_{n \rightarrow \infty} \|x_n - w\|$. This, together with (11) and conditions (i) and (ii), such that

$$\lim_{n \rightarrow \infty} \frac{(\|x_n - Ux_n\|^2 + \|(I - T)Ax_n\|^2)^2}{\|x_n - Ux_n + A^*(I - T)Ax_n\|^2} = 0.$$

By Lemma 6, this yields $\lim_{n \rightarrow \infty} \|x_n - Ux_n\| = \lim_{n \rightarrow \infty} \|(I - T)Ax_n\| = 0$. As shown in Theorem 1, we can get succession $\{x_{n_i}\}$ of $\{x_n\}$ such that $x_{n_i} \rightharpoonup p$.

By the definition of u_n , we deduce that

$$\lim_{n \rightarrow \infty} \|u_n - x_n\| = 0.$$

Let $z_n = (1 - \gamma_n)u_n = (1 - \gamma_n)(x_n - \rho_n((I - U)x_n + A^*(I - T)Ax_n))$, $n \geq 0$. Then

$$\|z_n - x_n\| \leq (1 - \gamma_n)\|u_n - x_n\| + \gamma_n\|x_n\|.$$

Furthermore, we obtain from (i) and the properties of $\{x_n\}$ that

$$\lim_{n \rightarrow \infty} \|z_n - x_n\| = 0.$$

This together with $x_{n_i} \rightharpoonup p$ implies that $z_{n_i} \rightharpoonup p$. So,

$$\limsup_{n \rightarrow \infty} \langle w, z_n - w \rangle = \lim_{i \rightarrow \infty} \langle w, z_{n_i} - w \rangle = \langle w, p - w \rangle \geq 0. \quad (12)$$

From (9) and (10):

$$\begin{aligned} \|x_{n+1} - w\|^2 &\leq (1 - \delta_n) \|x_n - w\|^2 + \delta_n \left((1 - \gamma_n) \|u_n - w\|^2 + 2 \langle \gamma_n(-w), z_n - w \rangle \right) \\ &\leq (1 - \gamma_n \delta_n) \|x_n - w\|^2 + 2 \gamma_n \delta_n \langle -w, z_n - w \rangle. \end{aligned} \quad (13)$$

We deduce $x_n \rightarrow z$ from Lemma 3 and Equations (12) and (13).

Case two. Suppose $n_0 \geq 0$, we have

$$\|x_{n_0} - w\| \leq \|x_{n_0+1} - w\|.$$

Setting $v_n = \{\|x_n - w\|\}$, then we have

$$v_{n_0} \leq v_{n_0+1}.$$

For all $n \geq n_0$, we now describe

$$\tau(n) = \max\{l \geq 1 : n_0 \leq l \leq n, v_l \leq v_{l+1}\}.$$

So $\{\tau(n)\}$ is non-decreasing satisfying

$$\lim_{n \rightarrow \infty} \tau(n) = \infty \text{ and } v_{\tau(n)} \leq v_{\tau(n)+1}.$$

As shown in Case 1, we get

$$\lim_{n \rightarrow \infty} \|z_{\tau(n)} - x_{\tau(n)}\| = 0.$$

This implies that

$$\omega_w(z_{\tau(n)}) \subset S.$$

Thus, we obtain

$$\limsup_{n \rightarrow \infty} \langle w, z_{\tau(n)} - w \rangle \geq 0. \quad (14)$$

By $v_{\tau(n)} \leq v_{\tau(n)+1}$, we have from (13) that

$$v_{\tau(n)}^2 \leq (1 - \gamma_{\tau(n)} \delta_{\tau(n)}) v_{\tau(n)}^2 + 2 \gamma_{\tau(n)} \delta_{\tau(n)} \langle -w, z_{\tau(n)} - w \rangle, \quad (15)$$

then

$$v_{\tau(n)}^2 \leq 2 \langle -w, z_{\tau(n)} - w \rangle. \quad (16)$$

Combining (14) and (16), we get

$$\limsup_{n \rightarrow \infty} v_{\tau(n)} \leq 0$$

and then

$$\lim_{n \rightarrow \infty} v_{\tau(n)} = 0. \quad (17)$$

By (15),

$$\limsup_{n \rightarrow \infty} v_{\tau(n)+1}^2 \leq \limsup_{n \rightarrow \infty} v_{\tau(n)}^2.$$

Using the above inequality and (17), we have

$$\lim_{n \rightarrow \infty} v_{\tau(n)+1} = 0.$$

By Lemma 4, this yields

$$0 \leq v_n \leq \max\{v_{\tau(n)}, v_{\tau(n)+1}\},$$

therefore, $v_n \rightarrow 0$, i.e., $x_n \rightarrow z$. $\square$

Algorithm 4. Initial guess x_0 is arbitrarily chosen and we assume x_n has been constructed. If

$$\|(I - P_C)x_n + A^*(I - P_Q)Ax_n\| = 0,$$

then stop (i.e., x_n solves problem (2)); otherwise, calculate the next x_{n+1} by the formula

$$x_{n+1} = (1 - \delta_n)x_n + \delta_n(1 - \gamma_n)(x_n - \rho_n((I - P_C)x_n + A^*(I - P_Q)Ax_n)), \quad (18)$$

where the stepsize sequence τ_n is chosen as

$$\rho_n = \frac{\|(I - P_C)x_n\|^2 + \|(I - P_Q)Ax_n\|^2}{\|(I - P_C)x_n + A^*(I - P_Q)Ax_n\|^2},$$

Theorem 4. Assume the parameters satisfy the following conditions.

- $\lim_{n \rightarrow \infty} \gamma_n = 0, \sum_{n=0}^{\infty} \gamma_n = +\infty;$
- $0 < \liminf_{n \rightarrow \infty} \delta_n(1 - \gamma_n) \leq \limsup_{n \rightarrow \infty} \delta_n(1 - \gamma_n) < 1.$

Then the sequence $\{x_n\}$ generated by (18) strongly converges to $z = \text{proj}_S 0$, which is the minimum-norm solution of the problem (2).

5. Numerical Example

Now, we illustrate the theoretical result by numerical examples.

Let $\mathcal{H} = \mathbf{R}$, inner product $\langle x, y \rangle = xy$, and norm $|\cdot|$. Let $x \in C$, $C = [0, +\infty)$ and $Ux = x + \frac{4}{x+1} - 1$. Clearly, $\text{Fix}(U) = 3$. It now

$$\langle x - y, Ux - Uy \rangle = \langle x - y, x + \frac{4}{x+1} - y - \frac{4}{y+1} \rangle \leq |x - y|^2$$

for all $x, y \in C$. Hence, U is a pseudo-contractive operator. So is $Tx = x + \frac{3}{x+2} - 1$. Truly, both U and T are satisfy the condition (6). For more detail of condition (6), please see the work by the authors of [29].

Let $x \in \mathbf{R}$, $Ax = \frac{1}{3}x$, $n \geq 1$, $\alpha_n = \frac{1}{n}$, $\beta_n = \frac{1}{8}$, then 3 is the approximation point of the Algorithm 1. Obviously, $A^* = A$, $\text{Fix}(U) = 3$, $\text{Fix}(T) = 1$ and $S = \{3\}$. Next, we rewrite Algorithm 1:

$$x_{n+1} = x_n - \frac{3(x_n + 6)(x_n - 3)}{(4x_n + 19)(x_n + 1)} - \frac{3(x_n + 1)(x_n - 3)}{(4x_n + 19)(x_n + 6)}, \quad n \geq 1.$$

Choosing initial values $x_1 = 5$ and $x_1 = 1$, respectively, we can see from Figure 1 and the numerical results in Table 1 that the theoretical result of Theorem 1 was demonstrated.

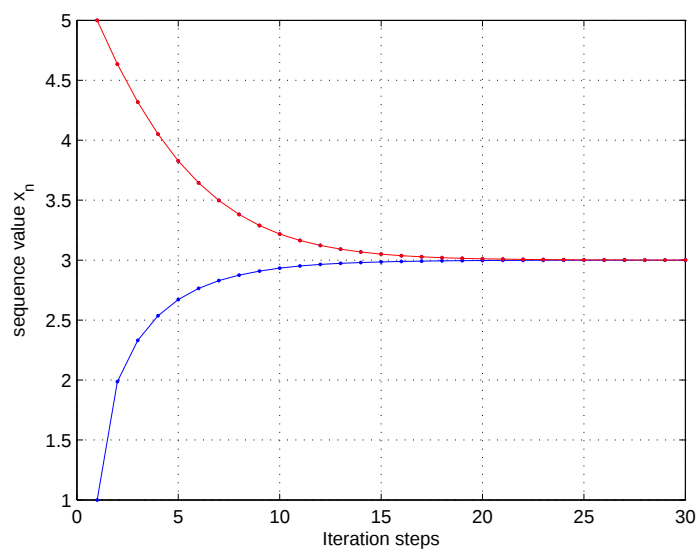


Figure 1. Weak convergence of $\{x_n\}$.

Table 1. The values of the sequence x_n .

n	x_n	x_n
1	5.000000000000000	1.000000000000000
2	4.634032634032634	1.987577639751553
3	4.318344257776301	2.331041217153439
4	4.050603724176485	2.536578840898608
5	3.827440529292961	2.671476773834111
...	...	...
27	3.001367761622588	2.999614025877529
28	3.001011222409336	2.999714678296111
29	3.000747602653792	2.999789081412396
30	3.000552695614674	2.999844081577622

Analogously, we now rewrite Algorithm 3 as follows.

$$x_{n+1} = \frac{7}{8}x_n + \frac{n-1}{8n} \left(x_n - \frac{3(x_n+6)(x_n-3)}{(4x_n+19)(x_n+1)} - \frac{3(x_n+1)(x_n-3)}{(4x_n+19)(x_n+6)} \right), \quad n \geq 1.$$

Also, choosing initial values $x_1 = 5$ and $x_1 = 1$, respectively, we can see from Figure 2 and the numerical results in Table 2 that the theoretical result of Theorem 3 was demonstrated.

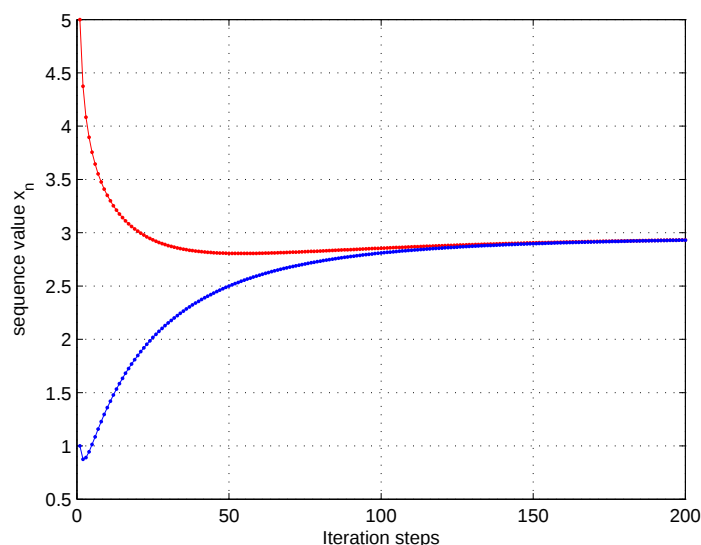


Figure 2. Strong convergence of $\{x_n\}$.

Table 2. The values of the sequence x_n .

n	x_n	x_n
1	5.000000000000000	1.000000000000000
2	4.375000000000000	0.875000000000000
3	4.084269250864560	0.890072601010101
4	3.895009196768486	0.944633146701975
5	3.754974188264774	1.012848796437149
...	...	...
197	2.930979205568247	2.929539367125689
198	2.931392607026035	2.930001859115338
199	2.931801371854751	2.930458028257748
200	2.932205567529454	2.930908000516816

We can see from Figures 1 and 2 that the rate of weak convergence may be faster than the strong one by comparing the iteration steps.

6. Conclusions

In this paper, we investigated the problem (1) involved in pseudo-contractive operators without Lipschitz assumption. By extending someone's results from [1,8,10,24–28] to the more general pseudo-contractive operators, we constructed two algorithm for solving the problem (1). Weak and strong convergence theorems are obtained under some mild hypotheses. Besides, we get the minimum-norm solution of problem (1); this is another interesting point. The results of this paper can be applied to engineering, network, and biotechnology.

Author Contributions: All authors participated in the conceptualization, validation, formal analysis, investigation, writing—original draft preparation, and writing—review and editing.

Funding: This work was supported by the Key Subject Program of Lingnan Normal University (Grant No. 1171518004), the Natural Science Foundation of Guangdong Province (2018A0303070012), and the Young Innovative Talents Project in Guangdong Universities (2017KQNCX125). Li-Jun Zhu was supported by the grants NXJG2017003, NXYLXK2017B09 and Advanced Intelligent Perception & Control Technology Innovative Team of NingXia.

Conflicts of Interest: The authors declare no conflicts of interest.

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