

Supplementary Materials

An Inverse QSAR Method Based on a Two-layered Model and Integer Programming

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1 An MILP Formulation for Inferring a Target Chemical Graph in Stage 4

1.1 Constructing Target Chemical Graphs

This section describes how to construct a target chemical graph in Stages 4 and 5.

1.1.1 Formulating an MILP for a prediction function in Stage 4

In Stage 3, we construct a prediction function $\eta_{\mathcal{N}} : \mathbb{R}^K \rightarrow \mathbb{R}$. It is known that the computation process of $\eta_{\mathcal{N}}(x)$ from a vector $x^* \in \mathbb{R}^K$ can be formulated as an MILP with the following property.

Theorem 1. ([1, 2]) *Let \mathcal{N} be an ANN with a piecewise-linear activation function for an input vector $x \in \mathbb{R}^K$, n_A denote the number of nodes in the architecture and n_B denote the total number of break-points over all activation functions. Then there is an MILP $\mathcal{M}(x, y; \mathcal{C}_1)$ that consists of variable vectors $x \in \mathbb{R}^K$, $y \in \mathbb{R}$, and an auxiliary variable vector $z \in \mathbb{R}^p$ for some integer $p = O(n_A + n_B)$ and a set \mathcal{C}_1 of $O(n_A + n_B)$ constraints on these variables such that: $\eta_{\mathcal{N}}(x^*) = y^*$ if and only if there is a vector (x^*, y^*) feasible to $\mathcal{M}(x, y; \mathcal{C}_1)$.*

Solving this MILP delivers a vector $x^* \in \mathbb{R}^K$ such that $\eta_{\mathcal{N}}(x^*) = y^*$ for a target value y^* . However, the resulting vector x^* may not admit a chemical graph G^* such that $f(G^*) = x^*$. To ensure that such chemical graph always exists in Stage 4, we further introduce some more constraints for a set of new variables in the next section.

1.1.2 Formulating an MILP for a feature vector and a target specification in Stage 4

In this section, we show an outline of formulation of an MILP that represents the computation process of a feature function $f(G)$ from a chemical graph G and a construction of a target chemical graph $G \in \mathcal{G}(G_C, \sigma_{\text{int}}, \sigma_{\text{ce}})$. Recall that the number of vertices in a target chemical graph is bounded by an upper bound n^* in a specification $(G_C, \sigma_{\text{int}}, \sigma_{\text{ce}})$. However, if we introduce a set of $(n^*)^2$ variables for all pairs of n^* vertices to present all possible graphs for a target chemical graph, then the resulting MILP formulation is hard to solve for $n^* > 20$ due to a larger number of variables and constraints. To overcome this, a sparse representation of chemical graphs has been proposed in the previous applications of the framework for acyclic graphs [3] and ρ -lean graphs [4]. We also define a similar sparse representation to formulate an MILP for our two-layered model.

Scheme Graphs We first regard a given seed graph G_C as a digraph and then add some more vertices and edges to construct a digraph, called a *scheme graph* $SG = (\mathcal{V}, \mathcal{E})$ so that any $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension H of G_C can be chosen as a subgraph of SG .

For a given target specification $(G_C, \sigma_{\text{int}}, \sigma_{\text{ce}})$, define integers that determine the size of a scheme graph SG as follows. $m_C := |E_C|$, $t_C := |V_C|$, $t_T := n_{\text{UB}}^{\text{int}} - |V_C|$, and $t_F := n^* - n_{\text{LB}}^{\text{int}}$.

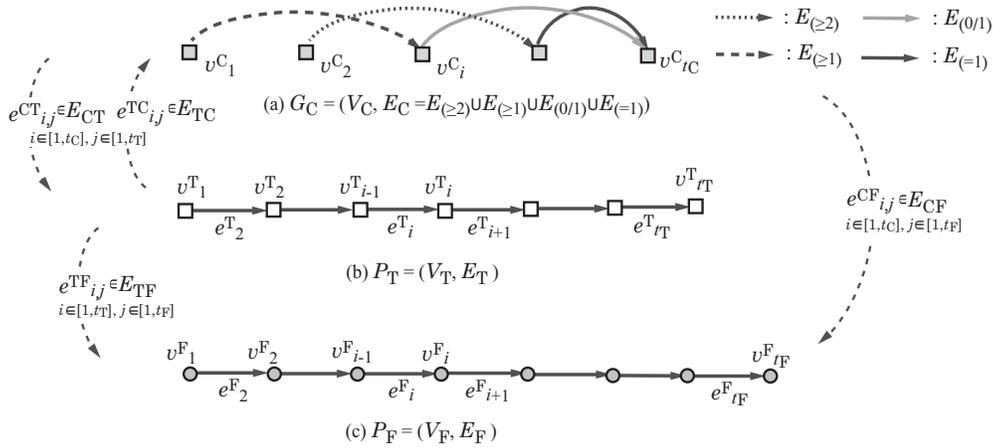


Figure 1: An illustration of a scheme graph SG : (a) A seed graph G_C ; (b) A path P_T of length $t_T - 1$; (c) A path P_F of length $t_F - 1$.

Formally the scheme graph $SG = (\mathcal{V}, \mathcal{E})$ is defined with a vertex set $\mathcal{V} = V_C \cup V_T \cup V_F$ and an edge set $\mathcal{E} = E_C \cup E_T \cup E_F \cup E_{CT} \cup E_{TC} \cup E_{CF} \cup E_{TF}$ that consist of the following sets. See Figure 1 for an illustration of these sets.

Construction of a σ_{int} -extension H^* of G_C : Denote the vertex set V_C and the edge set E_C in the seed graph G_C by $V_C = \{v^C_i \mid i \in [1, t_C]\}$ and $E_C = \{a_i \mid i \in [1, m_C]\}$, respectively, where V_C is always included in H^* . For including additional interior-vertices in H^* , introduce a path $P_T = (V_T = \{v^T_1, v^T_2, \dots, v^T_{t_T}\}, E_T = \{e^T_2, e^T_3, \dots, e^T_{t_T}\})$ of length $t_T - 1$ and a set E_{CT} (resp., E_{TC}) of directed edges $e^{CT}_{i,j} = (v^C_i, v^T_j)$ (resp., $e^{TC}_{i,j} = (v^T_j, v^C_i)$) $i \in [1, t_C]$, $j \in [1, t_T]$. In H^* , an edge $a_k = (v^C_i, v^C_{i'}) \in E_{(\ge 2)} \cup E_{(\ge 1)}$ is allowed to be replaced with a pure path P_k from vertex v^C_i to vertex $v^C_{i'}$ that visits a set of consecutive vertices $v^T_j, v^T_{j+1}, \dots, v^T_{j+p} \in V_T$ and edge $e^{TC}_{i,j} = (v^C_i, v^T_j) \in E_{CT}$, then edges $e^T_{j+1}, e^T_{j+2}, \dots, e^T_{j+p} \in E_T$ and finally edge $e^{TC}_{i',j+p} = (v^T_{j+p}, v^C_{i'}) \in E_{TC}$. The vertices in V_T selected in the path will be vertices in H^* .

Appending leaf paths with additional interior-edges in a $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension H of G_C : Introduce a path $P_F = (V_F = \{v^F_1, v^F_2, \dots, v^F_{t_F}\}, E_F = \{e^F_2, e^F_3, \dots, e^F_{t_F}\})$ of length $t_F - 1$, a set E_{CF} of directed edges $e^{CF}_{i,j} = (v^C_i, v^F_j)$, $i \in [1, t_C]$, $j \in [1, t_F]$, and a set E_{TF} of directed edges $e^{TF}_{i,j} = (v^T_i, v^F_j)$, $i \in [1, t_T]$, $j \in [1, t_F]$. In H , a leaf path Q with interior-edges that starts from a vertex $v^C_i \in V_C$ (resp., $v^T_i \in V_T$) visits a set of consecutive vertices $v^F_j, v^F_{j+1}, \dots, v^F_{j+p} \in V_F$ and edge $e^{CF}_{i,j} = (v^C_i, v^F_j) \in E_{CF}$ (resp., $e^{TF}_{i,j} = (v^T_i, v^F_j) \in E_{TF}$) and edges $e^F_{j+1}, e^F_{j+2}, \dots, e^F_{j+p} \in E_F$. In H , the edges and the vertices selected in the path Q are regarded as interior-edges and interior-vertices, respectively.

Construction of ρ -fringe-trees in a $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension G of G_C : In H , the root of a ρ -fringe-tree can be any vertex in $V_C \cup V_T \cup V_F$. For each vertex $v = v^C_i$ (resp., $v = v^T_i$ or v^F_i), we choose a chemical rooted tree T from the specified set $\mathcal{F}(v)$ (resp., \mathcal{F}_E).

Recall that the dimension K of a feature vector $x = f(G)$ used in constructing a prediction function $\eta_{\mathcal{N}}$ over a set of chemical graphs G is $K = 17 + |\Lambda^{\text{int}}(D_{\pi})| + |\Lambda^{\text{ex}}(D_{\pi})| + |\Gamma^{\text{int}}(D_{\pi})| + |\mathcal{F}(D_{\pi})|$. For a target specification $(G_{\text{C}}, \sigma_{\text{int}}, \sigma_{\text{ce}})$, let \mathcal{F}^* denote the set of chemical rooted trees ψ in the sets $\mathcal{F}(v)$, $v \in V_{\text{C}}$ and \mathcal{F}_E and $K^* := 17 + |\Lambda^{\text{int}}(D_{\pi})| + |\Lambda^{\text{ex}}(D_{\pi})| + |\Gamma^{\text{int}}(D_{\pi})| + |\mathcal{F}^*|$. Based on the scheme graph SG, we obtain the following MILP formulation $\mathcal{M}(x, g; \mathcal{C}_2)$.

Theorem 2. *Let $(G_{\text{C}}, \sigma_{\text{int}}, \sigma_{\text{ce}})$ be a target specification and $\varphi^* = |\Lambda^{\text{int}}(D_{\pi})| + |\Lambda^{\text{ex}}(D_{\pi})| + |\Gamma^{\text{int}}(D_{\pi})| + |\mathcal{F}^*|$ for sets of chemical elements, edge-configurations and fringe-configurations in σ_{ce} . Then there is an MILP $\mathcal{M}(x, g; \mathcal{C}_2)$ that consists of variable vectors $x \in \mathbb{R}^{K^*}$ and $g \in \mathbb{R}^q$ for an integer $q = O(n_{\text{UB}}^{\text{int}}(|E_{\text{C}}| + n^*) + (|E_{\text{C}}| + |\mathcal{V}|)\varphi^*)$ and a set \mathcal{C}_2 of $O(n_{\text{UB}}^{\text{int}}(|E_{\text{C}}| + n^*) + |\mathcal{V}|\varphi^*)$ constraints on x and g such that: (x^*, g^*) is feasible to $\mathcal{M}(x, g; \mathcal{C}_2)$ if and only if g^* forms a chemical graph $G \in \mathcal{G}(G_{\text{C}}, \sigma_{\text{int}}, \sigma_{\text{ce}})$ such that $f(G) = x^*$.*

Note that our MILP requires only $O(n^*)$ variables and constraints when the branch-parameter ρ , integers $|E_{\text{C}}|$, $n_{\text{UB}}^{\text{int}}$ and φ^* are constant. We explain the basic idea of our MILP that satisfies Theorem 2. The MILP mainly consists of the following three types of constraints.

- C1. Constraints for selecting an underlying graph H of a chemical graph $G \in \mathcal{G}(G_{\text{C}}, \sigma_{\text{int}}, \sigma_{\text{ce}})$ as a subgraph of the scheme graph SG;
- C2. Constraints for assigning chemical elements to interior-vertices and multiplicity to interior-edges to determine a chemical graph $G = (H, \alpha, \beta)$; and
- C3. Constraints for computing descriptors in the feature vector $f(G)$ of the selected chemical graph G .

In the constraints of C1, more formally we prepare the following.

Variables:

- a binary variable $v^{\text{X}}(i) \in \{0, 1\}$ for each vertex $v^{\text{X}}_i \in V_{\text{X}}$, $\text{X} \in \{\text{C}, \text{T}, \text{F}\}$ so that $v^{\text{X}}(i) = 1 \Leftrightarrow$ vertex v^{X}_i is used in a graph H selected from SG;
- a binary variable $e^{\text{X}}(i) \in \{0, 1\}$ (resp., $e^{\text{C}}(i) \in \{0, 1\}$) for each edge $e^{\text{X}}_i \in E_{\text{T}} \cup E_{\text{F}}$ (resp., $e^{\text{C}}_i = a_i \in E_{(\geq 2)} \cup E_{(\geq 1)} \cup E_{(0/1)}$) so that $e^{\text{X}}(i) = 1 \Leftrightarrow$ edge e^{X}_i is used in a graph H selected from SG. To save the number of variables in our MILP formulation, we do not prepare a binary variable $e^{\text{X}}(i, j) \in \{0, 1\}$ for any edge $e^{\text{X}}_{i,j} \in E_{\text{CT}} \cup E_{\text{TC}} \cup E_{\text{CF}} \cup E_{\text{TC}}$, where we represent a choice of edges in these sets by a set of $O(n^*|E_{\text{C}}|)$ variables (see Supplementary Materials for the details);
- binary variables $\delta_{\text{fr}}^{\text{C}}(i, \psi) \in \{0, 1\}$, $i \in [1, t_{\text{C}}]$, $\psi \in \mathcal{F}(v)$, $v = v^{\text{C}}_i \in V_{\text{C}}$ and $\delta_{\text{fr}}^{\text{T}}(i, \psi) \in \{0, 1\}$, $i \in [1, t_{\text{T}}]$, $\delta_{\text{fr}}^{\text{F}}(i, \psi) \in \{0, 1\}$, $i \in [1, t_{\text{F}}]$, $\psi \in \mathcal{F}_E$, where $\delta_{\text{fr}}^{\text{X}}(i, \psi) = 1$ ($\text{X} \in \{\text{C}, \text{T}, \text{F}\}$) if and only if the ρ -fringe-tree rooted at vertex v^{X}_i is r-isomorphic to ψ .

Constraints:

- linear constraints so that each ρ -fringe-tree rooted at a vertex v^{X}_i in a graph H from SG is selected from the given set $\mathcal{F}(v^{\text{C}}_i)$ for $\text{X}=\text{C}$ (or \mathcal{F}_E for $\text{X} \in \{\text{T}, \text{F}\}$);
- linear constraints such that each edge $e^{\text{C}}_i = a_i \in E_{(=1)}$ is always used as an edge in H and each edge $e^{\text{C}}_i = a_i \in E_{(0/1)}$ is used as an edge in H if necessary;

- linear constraints such that for each edge $a_k = (v^C_i, v^C_{i'}) \in E_{(\geq 2)}$, vertex $v^C_i \in V_C$ is connected to vertex $v^C_{i'} \in V_C$ in H by a pure path P_k that passes through some vertices in V_T and edges $e^{CT}_{i,j}, e^T_{j+1}, e^T_{j+2}, \dots, e^T_{j+p}, e^{TC}_{i',j+p}$ for some integers j and p ;
- linear constraints such that for each edge $a_k = (v^C_i, v^C_{i'}) \in E_{(\geq 1)}$, either the edge a_k is used as an edge in H or vertex $v^C_i \in V_C$ is connected to vertex $v^C_{i'} \in V_C$ in H by a pure path P_k as in the case of edges in $E_{(\geq 2)}$;
- linear constraints for selecting a leaf path Q_v rooted at a vertex $v = v^C_i$ (resp., $v = v^T_i$) with ρ -internal edges $e^{CF}_{i,j}$ (resp., $e^{TF}_{i,j}$), $e^F_{j+1}, e^F_{j+2}, \dots, e^F_{j+p}$ for some integers j and p .

In the constraints of C2, we prepare an integer variable $\alpha^X(i)$ for each vertex $v^X_i \in \mathcal{V}$, $X \in \{C, T, F\}$ in the scheme graph that represents the chemical element $\alpha(v^X_i) \in \Lambda$ if v^X_i is in a selected graph H (or $\alpha(v^X_i) = 0$ otherwise); integer variables $\beta^C : E_C \rightarrow [0, 3]$, $\beta^T : E_T \rightarrow [0, 3]$ and $\beta^F : E_F \rightarrow [0, 3]$ that represent the bond-multiplicity of edges in $E_C \cup E_T \cup E_F$; and integer variables $\beta^+, \beta^- : E_{(\geq 2)} \cup E_{(\geq 1)} \rightarrow [0, 3]$ and $\beta^{\text{in}} : V_C \cup V_T \rightarrow [0, 3]$ that represent the bond-multiplicity of edges in $E_{CT} \cup E_{TC} \cup E_{CF} \cup E_{TF}$. This determines a chemical graph $G = (H, \alpha, \beta)$. Also we include constraints for a selected chemical graph G to satisfy the valence condition at each interior-vertex v with the edge-configurations $\text{ec}(e)$ of the edges e incident to v and the chemical specification σ_{ce} .

In the constraints of C3, we introduce a variable for each descriptor and constraints with some more variables to compute the value of each descriptor in $f(G)$ for a selected chemical graph G .

The details of the MILP can be found in Section 3.

2 A Dynamic Programming Algorithm for Generating Isomers in Stage 5

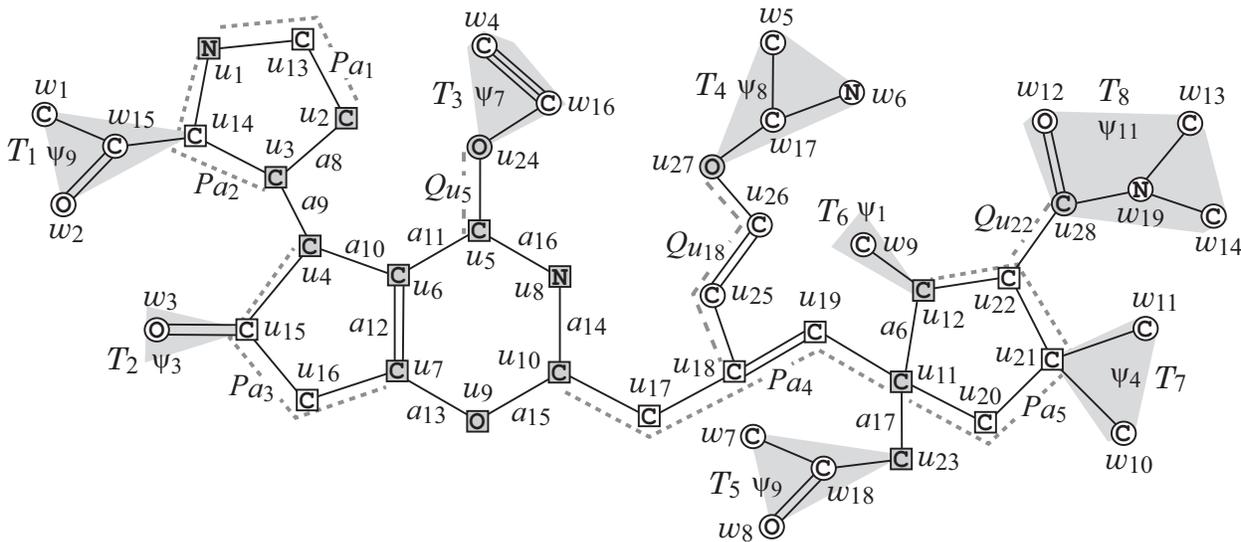


Figure 2: An illustration of a chemical graph G , where for $\rho = 2$, the exterior-vertices are w_1, w_2, \dots, w_{19} and the interior-vertices are u_1, u_2, \dots, u_{28} .

This section briefly reviews the method [4] for Stage 5. Let G^\dagger be a chemical graph that is a $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension of a seed graph $G_C = (V_C, E_C)$, where we denote by $E_{(=0)}$ the set of the edges

in $E_{(0/1)}$ that are not used in G^\dagger . We define a *base-graph* $G_B = (V_B, E_B)$ to be the seed graph $(V_C, E_C \setminus E_{(=0)})$ after removing the edges in $E_{(=0)}$. We call a chemical graph G^* a *chemical isomer* of G^\dagger if $f(G^*) = f(G^\dagger)$ and G^* is also a $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension of G_B .

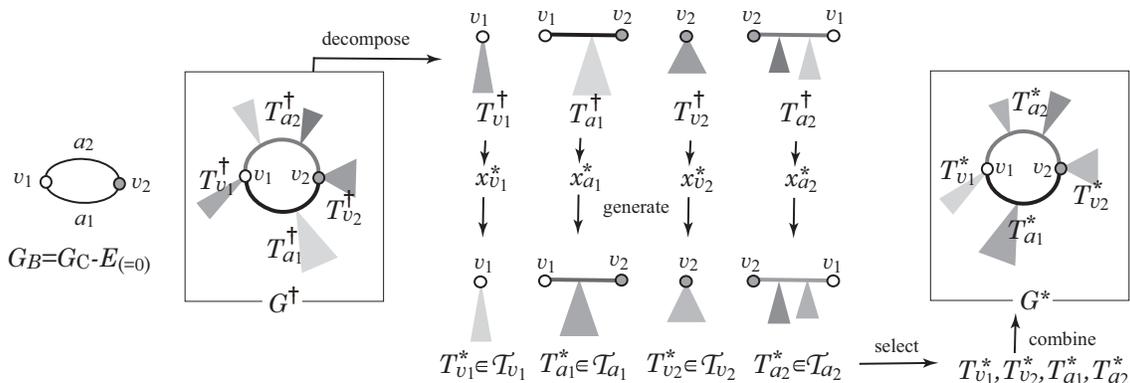


Figure 3: An illustration of generating a chemical isomer G^* of a chemical graph G^\dagger with a base-graph $G_B = (V_B, E_B)$.

The method generates chemical isomers G^* of G^\dagger in the following way, where Figure 3 illustrates the whole process in the case of $V_B = \{v_1, v_2\}$ and $E_B = \{a_1, a_2\}$.

1. We first decompose a given chemical graph G^\dagger into a collection of chemical rooted or bi-rooted trees.
 - For each vertex $v \in V_B$, let T_v^\dagger denote the chemical rooted tree rooted at v in G that is constructed with a leaf path Q_v and fringe-trees attached to Q_v . Possibly T_v^\dagger consists of a single vertex v and we call such a tree *trivial*.
 - For each edge $a = uv \in E_{(\geq 2)} \cup E_{(\geq 1)}$, let T_a^\dagger denote the chemical bi-rooted tree rooted at vertices u and v in G that consists of a pure u, v -path P_a , leaf paths rooted at internal vertices in P_a and fringe-trees attached to these leaf paths. Possibly T_a^\dagger consists of a single edge a and we call such a tree *trivial*.

Figure 4 illustrates the non-trivial chemical trees $T_t^\dagger, t \in V_B^* \cup E_B^*$ of the $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension $G^\dagger = G$ in Figure 2.

2. Let V_B^* (resp., E_B^*) denote the set of vertices $v \in V_B$ (resp., $a \in E_B$) such that T_v^\dagger (resp., T_a^\dagger) is not trivial. For each vertex or edge $t \in V_B^* \cup E_B^*$, compute the feature vector $x_t^* = f(T_t^\dagger)$ and then generate a set \mathcal{T}_t of all (or a limited number of) chemical acyclic graphs T_t^* such that $f(T_t^*) = x_t^*$ and the structure of T_t^* satisfies the lower and upper bounds in the interior-specification σ_{int} by using the dynamic programming algorithm for chemical acyclic graphs [3].
3. For each combination of chemical trees $T_t^* \in \mathcal{T}_t, t \in V_B^* \cup E_B^*$, a chemical graph G^* such that $f(G^*) = f(G^\dagger)$ is obtained from G^\dagger by replacing each tree T_t^\dagger with a new tree T_t^* . The number of such combinations is $\prod_{t \in V_B^* \cup E_B^*} |\mathcal{T}_t|$, where we ignore a possible automorphism of the resulting graphs G^* .

The above method [4] can be used to generate chemical isomers in Stage 5 in our two-layered model by making a minor modification to the definition of a feature vector $f(G)$.

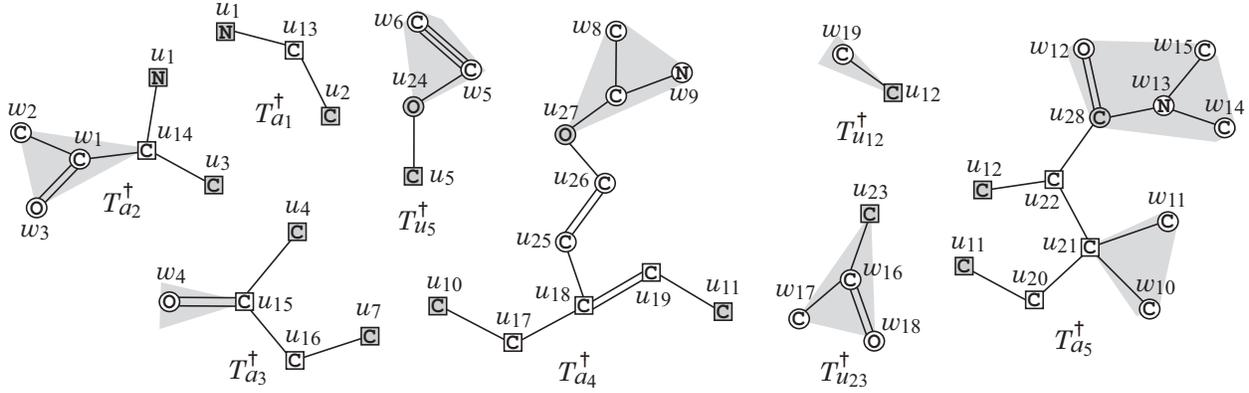


Figure 4: The non-trivial chemical rooted trees T_v^\dagger for $v \in \{u_5, u_{12}, u_{23}\} = V_B^*$ and the non-trivial chemical bi-rooted trees T_a^\dagger for $a \in \{a_1 = u_1u_2, a_2 = u_1u_3, a_3 = u_4u_7, a_4 = u_{10}u_{11}, a_5 = u_{11}u_{12}\} = E_B^*$ for the $(\sigma_{\text{int}}, \sigma_{\text{ce}})$ -extension $G^\dagger = G$ in Figure 2, where the gray squares indicate the roots of these rooted and bi-rooted trees.

3 All Constraints in an MILP Formulation for Chemical Graphs

We define a standard encoding of a finite set A of elements to be a bijection $\sigma : A \rightarrow [1, |A|]$, where we denote by $[A]$ the set $[1, |A|]$ of integers and by $[e]$ the encoded element $\sigma(e)$. Let ϵ denote *null*, a fictitious chemical element that does not belong to any set of chemical elements, chemical symbols, adjacency-configurations and edge-configurations in the following formulation. Given a finite set A , let A_ϵ denote the set $A \cup \{\epsilon\}$ and define a standard encoding of A_ϵ to be a bijection $\sigma : A_\epsilon \rightarrow [0, |A|]$ such that $\sigma(\epsilon) = 0$, where we denote by $[A_\epsilon]$ the set $[0, |A|]$ of integers and by $[e]$ the encoded element $\sigma(e)$, where $[\epsilon] = 0$.

3.1 Selecting a Cyclical-base

Recall that

$$E_{(=1)} = \{e \in E_C \mid \ell_{\text{LB}}(e) = \ell_{\text{UB}}(e) = 1\}; \quad E_{(0/1)} = \{e \in E_C \mid \ell_{\text{LB}}(e) = 0, \ell_{\text{UB}}(e) = 1\}; \\ E_{(\geq 1)} = \{e \in E_C \mid \ell_{\text{LB}}(e) = 1, \ell_{\text{UB}}(e) \geq 2\}; \quad E_{(\geq 2)} = \{e \in E_C \mid \ell_{\text{LB}}(e) \geq 2\};$$

- Every edge $a_i \in E_{(=1)}$ is included in G ;
- Each edge $a_i \in E_{(0/1)}$ is included in G if necessary;
- For each edge $a_i \in E_{(\geq 2)}$, edge a_i is not included in G and instead a path

$$P_i = (v_{\text{tail}(i)}^{\text{C}}, v_{j-1}^{\text{T}}, v_j^{\text{T}}, \dots, v_{j+t}^{\text{T}}, v_{\text{head}(i)}^{\text{C}})$$

of length at least 2 from vertex $v_{\text{tail}(i)}^{\text{C}}$ to vertex $v_{\text{head}(i)}^{\text{C}}$ visiting some vertices in V_{T} is constructed in G ; and

- For each edge $a_i \in E_{(\geq 1)}$, either edge a_i is directly used in G or the above path P_i of length at least 2 is constructed in G .

Let $t_C \triangleq |V_C|$ and denote V_C by $\{v^C_i \mid i \in [1, t_C]\}$. Regard the seed graph G_C as a digraph such that each edge a_i with end-vertices v^C_j and $v^C_{j'}$ is directed from v^C_j to $v^C_{j'}$ when $j < j'$. For each directed edge $a_i \in E_C$, let $\text{head}(i)$ and $\text{tail}(i)$ denote the head and tail of $e^C(i)$; i.e., $a_i = (v^C_{\text{tail}(i)}, v^C_{\text{head}(i)})$.

Assume that $E_C = \{a_i \mid i \in [1, m_C]\}$, $E_{(\geq 2)} = \{a_k \mid k \in [1, p]\}$, $E_{(\geq 1)} = \{a_k \mid k \in [p+1, q]\}$, $E_{(0/1)} = \{a_i \mid i \in [q+1, t]\}$ and $E_{(=1)} = \{a_i \mid i \in [t+1, m_C]\}$ for integers p, q and t . Let $I_{(=1)}$ denote the set of indices i of edges $a_i \in E_{(=1)}$. Similarly for $I_{(0/1)}$, $I_{(\geq 1)}$ and $I_{(\geq 2)}$.

Define

$$k_C \triangleq |E_{(\geq 2)} \cup E_{(\geq 1)}|, \quad \widetilde{k}_C \triangleq |E_{(\geq 2)}|.$$

To control the construction of such a path P_i for each edge $a_k \in E_{(\geq 2)} \cup E_{(\geq 1)}$, we regard the index $k \in [1, k_C]$ of each edge $a_k \in E_{(\geq 2)} \cup E_{(\geq 1)}$ as the ‘‘color’’ of the edge. To introduce necessary linear constraints that can construct such a path P_k properly in our MILP, we assign the color k to the vertices $v^T_{j-1}, v^T_j, \dots, v^T_{j+t}$ in V_T when the above path P_k is used in G .

For each index $s \in [1, t_C]$, let $I_C(s)$ denote the set of edges $e \in E_C$ incident to vertex v^C_s , and $E_{(=1)}^+(s)$ (resp., $E_{(=1)}^-(s)$) denote the set of edges $a_i \in E_{(=1)}$ such that the tail (resp., head) of a_i is vertex v^C_s . Similarly for $E_{(0/1)}^+(s)$, $E_{(0/1)}^-(s)$, $E_{(\geq 1)}^+(s)$, $E_{(\geq 1)}^-(s)$, $E_{(\geq 2)}^+(s)$ and $E_{(\geq 2)}^-(s)$. Let $I_C(s)$ denote the set of indices i of edges $a_i \in I_C(s)$. Similarly for $I_{(=1)}^+(s)$, $I_{(=1)}^-(s)$, $I_{(0/1)}^+(s)$, $I_{(0/1)}^-(s)$, $I_{(\geq 1)}^+(s)$, $I_{(\geq 1)}^-(s)$, $I_{(\geq 2)}^+(s)$ and $I_{(\geq 2)}^-(s)$. Note that $[1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$ and $[\widetilde{k}_C + 1, m_C] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}$.

constants:

- $t_C = |V_C|$, $\widetilde{k}_C = |E_{(\geq 2)}|$, $k_C = |E_{(\geq 2)} \cup E_{(\geq 1)}|$, $t_T = n_{\text{UB}}^{\text{int}} - |V_C|$, $m_C = |E_C|$. Note that $a_i \in E_C \setminus (E_{(\geq 2)} \cup E_{(\geq 1)})$ holds $i \in [k_C + 1, m_C]$;
- $\ell_{\text{LB}}(k), \ell_{\text{UB}}(k) \in [1, t_T]$, $k \in [1, k_C]$: lower and upper bounds on the length of path P_k ;

variables:

- $e^C(i) \in [0, 1]$, $i \in [1, m_C]$: $e^C(i)$ represents edge $a_i \in E_C$, $i \in [1, m_C]$ ($e^C(i) = 1$, $i \in I_{(=1)}$; $e^C(i) = 0$, $i \in I_{(\geq 2)}$) ($e^C(i) = 1 \Leftrightarrow$ edge a_i is used in G);
- $v^T(i) \in [0, 1]$, $i \in [1, t_T]$: $v^T(i) = 1 \Leftrightarrow$ vertex v^T_i is used in G ;
- $e^T(i) \in [0, 1]$, $i \in [1, t_T + 1]$: $e^T(i)$ represents edge $e^T_i = (v^T_{i-1}, v^T_i) \in E_T$, where e^T_1 and $e^T_{t_T+1}$ are fictitious edges ($e^T(i) = 1 \Leftrightarrow$ edge e^T_i is used in G);
- $\chi^T(i) \in [0, k_C]$, $i \in [1, t_T]$: $\chi^T(i)$ represents the color assigned to vertex v^T_i ($\chi^T(i) = k > 0 \Leftrightarrow$ vertex v^T_i is assigned color k ; $\chi^T(i) = 0$ means that vertex v^T_i is not used in G);
- $\text{clr}^T(k) \in [\ell_{\text{LB}}(k) - 1, \ell_{\text{UB}}(k) - 1]$, $k \in [1, k_C]$, $\text{clr}^T(0) \in [0, t_T]$: the number of vertices $v^T_i \in V_T$ with color c ;
- $\delta_\chi^T(k) \in [0, 1]$, $k \in [0, k_C]$: $\delta_\chi^T(k) = 1 \Leftrightarrow \chi^T(i) = k$ for some $i \in [1, t_T]$;
- $\chi^T(i, k) \in [0, 1]$, $i \in [1, t_T]$, $k \in [0, k_C]$ ($\chi^T(i, k) = 1 \Leftrightarrow \chi^T(i) = k$);
- $\widetilde{\text{deg}}_C^+(i) \in [0, 4]$, $i \in [1, t_C]$: the out-degree of vertex v^C_i with the used edges e^C in E_C ;
- $\widetilde{\text{deg}}_C^-(i) \in [0, 4]$, $i \in [1, t_C]$: the in-degree of vertex v^C_i with the used edges e^C in E_C ;

constraints:

$$e^C(i) = 1, \quad i \in I_{(=1)}, \quad (1)$$

$$e^C(i) = 0, \quad \text{clr}^T(i) \geq 1, \quad i \in I_{(\geq 2)}, \quad (2)$$

$$e^C(i) + \text{clr}^T(i) \geq 1, \quad \text{clr}^T(i) \leq t_T \cdot (1 - e^C(i)), \quad i \in I_{(\geq 1)}, \quad (3)$$

$$\sum_{c \in I_{(\geq 1)}^-(i) \cup I_{(0/1)}^-(i) \cup I_{(=1)}^-(i)} e^C(c) = \widetilde{\text{deg}}_C^-(i), \quad \sum_{c \in I_{(\geq 1)}^+(i) \cup I_{(0/1)}^+(i) \cup I_{(=1)}^+(i)} e^C(c) = \widetilde{\text{deg}}_C^+(i), \quad i \in [1, t_C], \quad (4)$$

$$\chi^T(i, 0) = 1 - v^T(i), \quad \sum_{k \in [0, k_C]} \chi^T(i, k) = 1, \quad \sum_{k \in [0, k_C]} k \cdot \chi^T(i, k) = \chi^T(i), \quad i \in [1, t_T], \quad (5)$$

$$\sum_{i \in [1, t_T]} \chi^T(i, k) = \text{clr}^T(k), \quad t_T \cdot \delta_\chi^T(k) \geq \sum_{i \in [1, t_T]} \chi^T(i, k) \geq \delta_\chi^T(k), \quad k \in [0, k_C], \quad (6)$$

$$v^T(i-1) \geq v^T(i), \quad k_C \cdot (v^T(i-1) - e^T(i)) \geq \chi^T(i-1) - \chi^T(i) \geq v^T(i-1) - e^T(i), \quad i \in [2, t_T]. \quad (7)$$

3.2 Constraints for Including Leaf Paths

Let \tilde{t}_C denote the number of vertices $u \in V_C$ such that $\text{bl}_{\text{UB}}(u) = 1$ and assume that $V_C = \{u_1, u_2, \dots, u_p\}$ so that

$$\text{bl}_{\text{UB}}(u_i) = 1, \quad i \in [1, \tilde{t}_C] \text{ and } \text{bl}_{\text{UB}}(u_i) = 0, \quad i \in [\tilde{t}_C + 1, t_C].$$

Define the set of colors for the vertex set $\{u_i \mid i \in [1, \tilde{t}_C]\} \cup V_T$ to be $[1, c_F]$ with

$$c_F \triangleq \tilde{t}_C + t_T = |\{u_i \mid i \in [1, \tilde{t}_C]\} \cup V_T|.$$

Let each vertex v^C_i , $i \in [1, \tilde{t}_C]$ (resp., $v^T_i \in V_T$) correspond to a color $i \in [1, c_F]$ (resp., $i + \tilde{t}_C \in [1, c_F]$). When a path $P = (u, v^F_j, v^F_{j+1}, \dots, v^F_{j+t})$ from a vertex $u \in V_C \cup V_T$ is used in G , we assign the color $i \in [1, c_F]$ of the vertex u to the vertices $v^F_j, v^F_{j+1}, \dots, v^F_{j+t} \in V_F$.

constants:

- c_F : the maximum number of different colors assigned to the vertices in V_F ;
- $n_{\text{LB}}^{\text{int}}, n_{\text{UB}}^{\text{int}} \in [2, n^*]$: lower and upper bounds on the number of interior-vertices in G ;
- $\text{bl}_{\text{LB}}(i) \in [0, 1]$, $i \in [1, \tilde{t}_C]$: a lower bound on the number of leaf ρ -branches in the leaf path rooted at a vertex v^C_i ;
- $\text{bl}_{\text{LB}}(k), \text{bl}_{\text{UB}}(k) \in [0, \ell_{\text{UB}}(k) - 1]$, $k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$: lower and upper bounds on the number of leaf ρ -branches in the trees rooted at internal vertices of a pure path P_k for an edge $a_k \in E_{(\geq 1)} \cup E_{(\geq 2)}$;

variables:

- $n_G^{\text{int}} \in [n_{\text{LB}}^{\text{int}}, n_{\text{UB}}^{\text{int}}]$: the number of interior-vertices in G ;
- $v^{\text{F}}(i) \in [0, 1]$, $i \in [1, t_{\text{F}}]$: $v^{\text{F}}(i) = 1 \Leftrightarrow$ vertex v_i^{F} is used in G ;
- $e^{\text{F}}(i) \in [0, 1]$, $i \in [1, t_{\text{F}} + 1]$: $e^{\text{F}}(i)$ represents edge $e_i^{\text{F}} = v_{i-1}^{\text{F}}v_i^{\text{F}}$, where e_1^{F} and $e_{t_{\text{F}}+1}^{\text{F}}$ are fictitious edges ($e^{\text{F}}(i) = 1 \Leftrightarrow$ edge e_i^{F} is used in G);
- $\chi^{\text{F}}(i) \in [0, c_{\text{F}}]$, $i \in [1, t_{\text{F}}]$: $\chi^{\text{F}}(i)$ represents the color assigned to vertex v_i^{F} ($\chi^{\text{F}}(i) = c \Leftrightarrow$ vertex v_i^{F} is assigned color c);
- $\text{clr}^{\text{F}}(c) \in [0, t_{\text{F}}]$, $c \in [0, c_{\text{F}}]$: the number of vertices v_i^{F} with color c ;
- $\delta_{\chi}^{\text{F}}(c) \in [\text{bl}_{\text{LB}}(c), 1]$, $c \in [1, \tilde{t}_{\text{C}}]$: $\delta_{\chi}^{\text{F}}(c) = 1 \Leftrightarrow \chi^{\text{F}}(i) = c$ for some $i \in [1, t_{\text{F}}]$;
- $\delta_{\chi}^{\text{F}}(c) \in [0, 1]$, $c \in [\tilde{t}_{\text{C}} + 1, c_{\text{F}}]$: $\delta_{\chi}^{\text{F}}(c) = 1 \Leftrightarrow \chi^{\text{F}}(i) = c$ for some $i \in [1, t_{\text{F}}]$;
- $\chi^{\text{F}}(i, c) \in [0, 1]$, $i \in [1, t_{\text{F}}]$, $c \in [0, c_{\text{F}}]$: $\chi^{\text{F}}(i, c) = 1 \Leftrightarrow \chi^{\text{F}}(i) = c$;
- $\text{bl}(k, i) \in [0, 1]$, $k \in [1, k_{\text{C}}] = I_{(\geq 2)} \cup I_{(\geq 1)}$, $i \in [1, t_{\text{T}}]$: $\text{bl}(k, i) = 1 \Leftrightarrow$ path P_k contains vertex v_i^{T} as an internal vertex and the ρ -fringe-tree rooted at v_i^{T} contains a leaf ρ -branch;

constraints:

$$\chi^{\text{F}}(i, 0) = 1 - v^{\text{F}}(i), \quad \sum_{c \in [0, c_{\text{F}}]} \chi^{\text{F}}(i, c) = 1, \quad \sum_{c \in [0, c_{\text{F}}]} c \cdot \chi^{\text{F}}(i, c) = \chi^{\text{F}}(i), \quad i \in [1, t_{\text{F}}], \quad (8)$$

$$\sum_{i \in [1, t_{\text{F}}]} \chi^{\text{F}}(i, c) = \text{clr}^{\text{F}}(c), \quad t_{\text{F}} \cdot \delta_{\chi}^{\text{F}}(c) \geq \sum_{i \in [1, t_{\text{F}}]} \chi^{\text{F}}(i, c) \geq \delta_{\chi}^{\text{F}}(c), \quad c \in [0, c_{\text{F}}], \quad (9)$$

$$e^{\text{F}}(1) = e^{\text{F}}(t_{\text{F}} + 1) = 0, \quad (10)$$

$$\begin{aligned} v^{\text{F}}(i-1) &\geq v^{\text{F}}(i), \\ c_{\text{F}} \cdot (v^{\text{F}}(i-1) - e^{\text{F}}(i)) &\geq \chi^{\text{F}}(i-1) - \chi^{\text{F}}(i) \geq v^{\text{F}}(i-1) - e^{\text{F}}(i), \end{aligned} \quad i \in [2, t_{\text{F}}], \quad (11)$$

$$\text{bl}(k, i) \geq \delta_{\chi}^{\text{F}}(\tilde{t}_{\text{C}} + i) + \chi^{\text{T}}(i, k) - 1, \quad k \in [1, k_{\text{C}}], i \in [1, t_{\text{T}}], \quad (12)$$

$$\sum_{k \in [1, k_{\text{C}}], i \in [1, t_{\text{T}}]} \text{bl}(k, i) \leq \sum_{i \in [1, t_{\text{T}}]} \delta_{\chi}^{\text{F}}(\tilde{t}_{\text{C}} + i), \quad (13)$$

$$\text{bl}_{\text{LB}}(k) \leq \sum_{i \in [1, t_{\text{T}}]} \text{bl}(k, i) \leq \text{bl}_{\text{UB}}(k), \quad k \in [1, k_{\text{C}}], \quad (14)$$

$$t_{\text{C}} + \sum_{i \in [1, t_{\text{T}}]} v^{\text{T}}(i) + \sum_{i \in [1, t_{\text{F}}]} v^{\text{F}}(i) = n_G^{\text{int}}. \quad (15)$$

3.3 Constraints for Including Fringe-trees

To express the condition that the ρ -fringe-tree is chosen from a rooted tree C_i , T_i or F_i , we introduce the following set of variables and constraints.

constants:

- n_{LB}, n^* : lower and upper bounds on $n(G)$, where $n_{\text{LB}}, n^* \geq n_{\text{LB}}^{\text{int}}$;
- $\text{ch}_{\text{LB}}(i), \text{ch}_{\text{UB}}(i) \in [0, n^*]$, $i \in [1, t_{\text{T}}]$: lower and upper bounds on $\text{ht}(T_i)$ of the tree T_i rooted at a vertex v_i^{C} ;
- $\text{ch}_{\text{LB}}(k), \text{ch}_{\text{UB}}(k) \in [0, n^*]$, $k \in [1, k_{\text{C}}] = I_{(\geq 2)} \cup I_{(\geq 1)}$: lower and upper bounds on the maximum height $\text{ht}(T)$ of the tree $T \in \mathcal{F}(P_k)$ rooted at an internal vertex of a path P_k for an edge $a_k \in E_{(\geq 1)} \cup E_{(\geq 2)}$;
- Let \mathcal{F}_{Λ} denote the set of chemical rooted trees $\psi = (\{v\}, \emptyset)$ with $\text{ht}(\psi) = 0$ and $\alpha(v) = \mathbf{a}$ for each chemical element $\mathbf{a} \in \Lambda$;
- Prepare a coding of the set $\mathcal{F}(D_{\pi})$ and let $[\psi]$ denote the coded integer of an element ψ in $\mathcal{F}(D_{\pi})$;
- Sets $\mathcal{F}(v) \subseteq \mathcal{F}(D_{\pi})$, $v \in V_{\text{C}}$ and $\mathcal{F}_E \subseteq \mathcal{F}(D_{\pi})$ of chemical rooted trees T with $\text{ht}(T) \in [1, \rho]$;
- Define $\mathcal{F}^* := \bigcup_{v \in V_{\text{C}}} \mathcal{F}(v) \cup \mathcal{F}_E$, $\mathcal{F}_i^{\text{C}} := \mathcal{F}(v_i^{\text{C}})$, $i \in [1, t_{\text{C}}]$, $\mathcal{F}_i^{\text{T}} := \mathcal{F}_E$, $i \in [1, t_{\text{T}}]$ and $\mathcal{F}_i^{\text{F}} := \mathcal{F}_E$, $i \in [1, t_{\text{F}}]$;
- $\mathcal{F}_i^{\text{X}}[p]$, $p \in [1, \rho]$, $\text{X} \in \{\text{C}, \text{T}, \text{F}\}$: the set of chemical rooted trees $T \in \mathcal{F}_i^{\text{X}}$ with $\text{ht}(T) = p$;
- $n([\psi]) \in [0, 3^{\rho}]$, $\psi \in \mathcal{F}^*$: the number of non-root vertices in a chemical rooted tree ψ ;
- $\text{ht}([\psi]) \in [0, \rho]$, $\psi \in \mathcal{F}^*$: the height of a chemical rooted tree ψ ;
- $\text{deg}_r([\psi]) \in [0, 4]$, $\psi \in \mathcal{F}^*$: the number of children of the root r of a chemical rooted tree ψ ;

variables:

- $n_G \in [n_{\text{LB}}, n^*]$: $n(G)$;
- $v^{\text{X}}(i) \in [0, 1]$, $i \in [1, t_{\text{X}}]$, $\text{X} \in \{\text{T}, \text{F}\}$: $v^{\text{X}}(i) = 1 \Leftrightarrow$ vertex v_i^{X} is used in G ;
- $h^{\text{X}}(i) \in [0, \rho]$, $i \in [1, t_{\text{X}}]$, $\text{X} \in \{\text{C}, \text{T}, \text{F}\}$: the height of the ρ -fringe-tree rooted at vertex v_i^{X} in G ;
- $\delta_{\text{fr}}^{\text{X}}(i, [\psi]) \in [0, 1]$, $i \in [1, t_{\text{X}}]$, $\psi \in \mathcal{F}_i^{\text{X}} \cup \mathcal{F}_{\Lambda}$, $\text{X} \in \{\text{T}, \text{F}\}$: $\delta_{\text{fr}}^{\text{X}}(i, [\psi]) = 1 \Leftrightarrow \psi$ is the ρ -fringe-tree at vertex v_i^{X} , where $\psi \in \mathcal{F}_{\Lambda}$ means that the height of the ρ -fringe-tree is 0;
- $\text{deg}_X^{\text{ex}}(i) \in [0, 3]$, $i \in [1, t_{\text{X}}]$, $\text{X} \in \{\text{C}, \text{T}, \text{F}\}$: the number of children of the root of the ρ -fringe-tree rooted at vertex v_i^{X} in G ;
- $\sigma(k, i) \in [0, 1]$, $k \in [1, k_{\text{C}}] = I_{(\geq 2)} \cup I_{(\geq 1)}$, $i \in [1, t_{\text{T}}]$: $\sigma(k, i) = 1 \Leftrightarrow$ the ρ -fringe-tree T_v rooted at vertex $v = v_i^{\text{T}}$ with color k has the largest height among such trees;

constraints:

$$\begin{aligned}
 \sum_{\psi \in \mathcal{F}_i^{\text{C}} \cup \mathcal{F}_{\Lambda}} \delta_{\text{fr}}^{\text{C}}(i, [\psi]) &= 1, & \sum_{\psi \in \mathcal{F}_i^{\text{C}} \cup \mathcal{F}_{\Lambda}} \text{deg}_r([\psi]) \cdot \delta_{\text{fr}}^{\text{C}}(i, [\psi]) &= \text{deg}_{\text{C}}^{\text{ex}}(i), & i \in [1, t_{\text{C}}], \\
 \sum_{\psi \in \mathcal{F}_i^{\text{X}} \cup \mathcal{F}_{\Lambda}} \delta_{\text{fr}}^{\text{X}}(i, [\psi]) &= v^{\text{X}}(i), & \sum_{\psi \in \mathcal{F}_i^{\text{X}} \cup \mathcal{F}_{\Lambda}} \text{deg}_r([\psi]) \cdot \delta_{\text{fr}}^{\text{X}}(i, [\psi]) &= \text{deg}_X^{\text{ex}}(i), & i \in [1, t_{\text{X}}], \text{X} \in \{\text{T}, \text{F}\}, \quad (16)
 \end{aligned}$$

$$\sum_{\psi \in \mathcal{F}_i^F[\rho]} \delta_{\text{fr}}^F(i, [\psi]) \geq v^F(i) - e^F(i+1), \quad i \in [1, t_F] \ (e^F(t_F+1) = 0), \quad (17)$$

$$\sum_{\psi \in \mathcal{F}_i^X} \text{ht}([\psi]) \cdot \delta_{\text{fr}}^X(i, [\psi]) = h^X(i), \quad i \in [1, t_X], X \in \{C, T, F\}, \quad (18)$$

$$\sum_{\substack{\psi \in \mathcal{F}_i^X \\ i \in [1, t_X], X \in \{C, T, F\}}} n([\psi]) \cdot \delta_{\text{fr}}^X(i, [\psi]) + \sum_{i \in [1, t_X], X \in \{T, F\}} v^X(i) + t_C = n_G, \quad (19)$$

$$\begin{aligned} h^C(i) &\geq \text{ch}_{\text{LB}}(i) - n^* \delta_{\chi}^F(i), \quad \text{clr}^F(i) + \rho \geq \text{ch}_{\text{LB}}(i), \\ h^C(i) &\leq \text{ch}_{\text{UB}}(i), \quad \text{clr}^F(i) + \rho \leq \text{ch}_{\text{UB}}(i) + n^*(1 - \delta_{\chi}^F(i)), \quad i \in [1, \tilde{t}_C], \end{aligned} \quad (20)$$

$$\text{ch}_{\text{LB}}(i) \leq h^C(i) \leq \text{ch}_{\text{UB}}(i), \quad i \in [\tilde{t}_C + 1, t_C], \quad (21)$$

$$\begin{aligned} h^T(i) &\leq \text{ch}_{\text{UB}}(k) + n^*(\delta_{\chi}^F(\tilde{t}_C + i) + 1 - \chi^T(i, k)), \\ \text{clr}^F(\tilde{t}_C + i) + \rho &\leq \text{ch}_{\text{UB}}(k) + n^*(2 - \delta_{\chi}^F(\tilde{t}_C + i) - \chi^T(i, k)), \\ &k \in [1, k_C], i \in [1, t_T], \end{aligned} \quad (22)$$

$$\sum_{i \in [1, t_T]} \sigma(k, i) = \delta_{\chi}^T(k), \quad k \in [1, k_C], \quad (23)$$

$$\begin{aligned} \chi^T(i, k) &\geq \sigma(k, i), \\ h^T(i) &\geq \text{ch}_{\text{LB}}(k) - n^*(\delta_{\chi}^F(\tilde{t}_C + i) + 1 - \sigma(k, i)), \\ \text{clr}^F(\tilde{t}_C + i) + \rho &\geq \text{ch}_{\text{LB}}(k) - n^*(2 - \delta_{\chi}^F(\tilde{t}_C + i) - \sigma(k, i)), \quad k \in [1, k_C], i \in [1, t_T]. \end{aligned} \quad (24)$$

3.4 Descriptor for the Number of Specified Degree

We include constraints to compute descriptors $\text{dg}_d^{\text{int}}(G)$, $d \in [1, 4]$.

variables:

- $\text{deg}^X(i) \in [0, 4]$, $i \in [1, t_X]$, $X \in \{C, T, F\}$: the degree $\text{deg}_G(v^X_i)$ of vertex v^X_i in G ;
- $\text{deg}_{\text{CT}}(i) \in [0, 4]$, $i \in [1, t_C]$: the number of edges from vertex v^C_i to vertices v^T_j , $j \in [1, t_T]$;
- $\text{deg}_{\text{TC}}(i) \in [0, 4]$, $i \in [1, t_C]$: the number of edges from vertices v^T_j , $j \in [1, t_T]$ to vertex v^C_i ;

- $\delta_{\text{dg}}^{\text{C}}(i, d) \in [0, 1]$, $i \in [1, t_{\text{C}}]$, $d \in [1, 4]$, $\delta_{\text{dg}}^{\text{X}}(i, d) \in [0, 1]$, $i \in [1, t_{\text{X}}]$, $d \in [0, 4]$, $\text{X} \in \{\text{T}, \text{F}\}$:
 $\delta_{\text{dg}}^{\text{X}}(i, d) = 1 \Leftrightarrow \text{deg}^{\text{X}}(i) = d$;
- $\text{dg}(d) \in [\text{dg}_{\text{LB}}(d), \text{dg}_{\text{UB}}(d)]$, $d \in [1, 4]$: the number of interior-vertices v with $\text{deg}_G(v) = d$;
- $\text{deg}_{\text{C}}^{\text{int}}(i) \in [1, 4]$, $i \in [1, t_{\text{C}}]$, $\text{deg}_{\text{X}}^{\text{int}}(i) \in [0, 4]$, $i \in [1, t_{\text{X}}]$, $\text{X} \in \{\text{T}, \text{F}\}$: the interior-degree $\text{deg}_{(V^{\text{int}}, E^{\text{int}})}(v^{\text{X}}_i)$; i.e., the number of interior-edges incident to vertex v^{X}_i ;
- $\delta_{\text{dg}, \text{C}}^{\text{int}}(i, d) \in [0, 1]$, $i \in [1, t_{\text{C}}]$, $d \in [1, 4]$, $\delta_{\text{dg}, \text{X}}^{\text{int}}(i, d) \in [0, 1]$, $i \in [1, t_{\text{X}}]$, $d \in [0, 4]$, $\text{X} \in \{\text{T}, \text{F}\}$:
 $\delta_{\text{dg}, \text{X}}^{\text{int}}(i, d) = 1 \Leftrightarrow \text{deg}_{\text{X}}^{\text{int}}(i) = d$;
- $\text{dg}^{\text{int}}(d) \in [\text{dg}_{\text{LB}}(d), \text{dg}_{\text{UB}}(d)]$, $d \in [1, 4]$: the number of interior-vertices v with the interior-degree $\text{deg}_{(V^{\text{int}}, E^{\text{int}})}(v) = d$;

constraints:

$$\sum_{k \in I_{(\geq 2)}^+(i) \cup I_{(\geq 1)}^+(i)} \delta_{\chi}^{\text{T}}(k) = \text{deg}_{\text{CT}}(i), \quad \sum_{k \in I_{(\geq 2)}^-(i) \cup I_{(\geq 1)}^-(i)} \delta_{\chi}^{\text{T}}(k) = \text{deg}_{\text{TC}}(i), \quad i \in [1, t_{\text{C}}], \quad (25)$$

$$\widetilde{\text{deg}}_{\text{C}}^-(i) + \widetilde{\text{deg}}_{\text{C}}^+(i) + \text{deg}_{\text{CT}}(i) + \text{deg}_{\text{TC}}(i) + \delta_{\chi}^{\text{F}}(i) = \text{deg}_{\text{C}}^{\text{int}}(i), \quad i \in [1, \widetilde{t}_{\text{C}}], \quad (26)$$

$$\widetilde{\text{deg}}_{\text{C}}^-(i) + \widetilde{\text{deg}}_{\text{C}}^+(i) + \text{deg}_{\text{CT}}(i) + \text{deg}_{\text{TC}}(i) = \text{deg}_{\text{C}}^{\text{int}}(i), \quad i \in [\widetilde{t}_{\text{C}} + 1, t_{\text{C}}], \quad (27)$$

$$\text{deg}_{\text{C}}^{\text{int}}(i) + \text{deg}_{\text{C}}^{\text{ex}}(i) = \text{deg}_{\text{C}}^{\text{C}}(i), \quad i \in [1, t_{\text{C}}], \quad (28)$$

$$\sum_{\psi \in \mathcal{F}_i^{\text{C}}[\rho]} \delta_{\text{fr}}^{\text{C}}(i, [\psi]) \geq 2 - \text{deg}_{\text{C}}^{\text{int}}(i) \quad i \in [1, t_{\text{C}}], \quad (29)$$

$$\begin{aligned} 2v^{\text{T}}(i) + \delta_{\chi}^{\text{F}}(\widetilde{t}_{\text{C}} + i) &= \text{deg}_{\text{T}}^{\text{int}}(i), \\ \text{deg}_{\text{T}}^{\text{int}}(i) + \text{deg}_{\text{T}}^{\text{ex}}(i) &= \text{deg}_{\text{T}}^{\text{T}}(i), \quad i \in [1, t_{\text{T}}] \quad (e^{\text{T}}(1) = e^{\text{T}}(t_{\text{T}} + 1) = 0), \end{aligned} \quad (30)$$

$$\begin{aligned} v^{\text{F}}(i) + e^{\text{F}}(i + 1) &= \text{deg}_{\text{F}}^{\text{int}}(i), \\ \text{deg}_{\text{F}}^{\text{int}}(i) + \text{deg}_{\text{F}}^{\text{ex}}(i) &= \text{deg}_{\text{F}}^{\text{F}}(i), \quad i \in [1, t_{\text{F}}] \quad (e^{\text{F}}(1) = e^{\text{F}}(t_{\text{F}} + 1) = 0), \end{aligned} \quad (31)$$

$$\begin{aligned} \sum_{d \in [0, 4]} \delta_{\text{dg}}^{\text{X}}(i, d) &= 1, \quad \sum_{d \in [1, 4]} d \cdot \delta_{\text{dg}}^{\text{X}}(i, d) = \text{deg}^{\text{X}}(i), \\ \sum_{d \in [0, 4]} \delta_{\text{dg}, \text{X}}^{\text{int}}(i, d) &= 1, \quad \sum_{d \in [1, 4]} d \cdot \delta_{\text{dg}, \text{X}}^{\text{int}}(i, d) = \text{deg}_{\text{X}}^{\text{int}}(i), \quad i \in [1, t_{\text{X}}], \text{X} \in \{\text{T}, \text{C}, \text{F}\}, \end{aligned} \quad (32)$$

$$\begin{aligned} \sum_{i \in [1, t_{\text{C}}]} \delta_{\text{dg}}^{\text{C}}(i, d) + \sum_{i \in [1, t_{\text{T}}]} \delta_{\text{dg}}^{\text{T}}(i, d) + \sum_{i \in [1, t_{\text{F}}]} \delta_{\text{dg}}^{\text{F}}(i, d) &= \text{dg}(d), \\ \sum_{i \in [1, t_{\text{C}}]} \delta_{\text{dg}, \text{C}}^{\text{int}}(i, d) + \sum_{i \in [1, t_{\text{T}}]} \delta_{\text{dg}, \text{T}}^{\text{int}}(i, d) + \sum_{i \in [1, t_{\text{F}}]} \delta_{\text{dg}, \text{F}}^{\text{int}}(i, d) &= \text{dg}^{\text{int}}(d), \quad d \in [1, 4]. \end{aligned} \quad (33)$$

3.5 Assigning Multiplicity

We prepare an integer variable $\beta(e)$ for each edge e in the scheme graph SG to denote the bond-multiplicity of e in a selected graph G and include necessary constraints for the variables to satisfy in G .

constants:

- $\beta_r([\psi])$: the sum of bond-multiplicities of edges incident to the root of a tree $\psi \in \mathcal{F}^*$;

variables:

- $\beta^X(i) \in [0, 3]$, $i \in [2, t_X]$, $X \in \{T, F\}$: the bond-multiplicity of edge e^X_i ;
- $\beta^C(i) \in [0, 3]$, $i \in [\widetilde{k}_C + 1, m_C] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}$: the bond-multiplicity of edge $a_i \in E_{(\geq 1)} \cup E_{(0/1)} \cup E_{(=1)}$;
- $\beta^+(k), \beta^-(k) \in [0, 3]$, $k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$: the bond-multiplicity of the first (resp., last) edge of the pure path P_k ;
- $\beta^{\text{in}}(c) \in [0, 3]$, $c \in [1, c_F]$: the bond-multiplicity of the first edge of the leaf path Q_c rooted at vertex c ;
- $\beta^X_{\text{ex}}(i) \in [0, 4]$, $i \in [1, t_X]$, $X \in \{C, T, F\}$: the sum $\beta_{T_v}(v)$ of bond-multiplicities of edges in the ρ -fringe-tree T_v rooted at interior-vertex $v = v^X_i$;
- $\delta^X_\beta(i, m) \in [0, 1]$, $i \in [2, t_X]$, $m \in [0, 3]$, $X \in \{T, F\}$: $\delta^X_\beta(i, m) = 1 \Leftrightarrow \beta^X(i) = m$;
- $\delta^C_\beta(i, m) \in [0, 1]$, $i \in [\widetilde{k}_C, m_C] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}$, $m \in [0, 3]$: $\delta^C_\beta(i, m) = 1 \Leftrightarrow \beta^C(i) = m$;
- $\delta^+_\beta(k, m), \delta^-_\beta(k, m) \in [0, 1]$, $k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$, $m \in [0, 3]$: $\delta^+_\beta(k, m) = 1$ (resp., $\delta^-_\beta(k, m) = 1$) $\Leftrightarrow \beta^+(k) = m$ (resp., $\beta^-(k) = m$);
- $\delta^{\text{in}}_\beta(c, m) \in [0, 1]$, $c \in [1, c_F]$, $m \in [0, 3]$: $\delta^{\text{in}}_\beta(c, m) = 1 \Leftrightarrow \beta^{\text{in}}(c) = m$;
- $\text{bd}^{\text{int}}(m) \in [0, 2n_{\text{UB}}^{\text{int}}]$, $m \in [1, 3]$: the number of interior-edges with bond-multiplicity m in G ;
- $\text{bd}_X(m) \in [0, 2n_{\text{UB}}^{\text{int}}]$, $X \in \{C, T, CT, TC\}$, $\text{bd}_X(m) \in [0, 2n_{\text{UB}}^{\text{int}}]$, $X \in \{F, CF, TF\}$, $m \in [1, 3]$: the number of interior-edges $e \in E_X$ with bond-multiplicity m in G ;

constraints:

$$e^C(i) \leq \beta^C(i) \leq 3e^C(i), i \in [\widetilde{k}_C + 1, m_C] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}, \quad (34)$$

$$e^X(i) \leq \beta^X(i) \leq 3e^X(i), \quad i \in [2, t_X], X \in \{T, F\}, \quad (35)$$

$$\delta^T_\chi(k) \leq \beta^+(k) \leq 3\delta^T_\chi(k), \quad \delta^T_\chi(k) \leq \beta^-(k) \leq 3\delta^T_\chi(k), \quad k \in [1, k_C], \quad (36)$$

$$\delta^F_\chi(c) \leq \beta^{\text{in}}(c) \leq 3\delta^F_\chi(c), \quad c \in [1, c_F], \quad (37)$$

$$\sum_{m \in [0,3]} \delta_{\beta}^X(i, m) = 1, \quad \sum_{m \in [0,3]} m \cdot \delta_{\beta}^X(i, m) = \beta^X(i), \quad i \in [2, t_X], X \in \{T, F\}, \quad (38)$$

$$\sum_{m \in [0,3]} \delta_{\beta}^C(i, m) = 1, \quad \sum_{m \in [0,3]} m \cdot \delta_{\beta}^C(i, m) = \beta^C(i), \quad i \in [\widetilde{k}_C + 1, m_C], \quad (39)$$

$$\begin{aligned} \sum_{m \in [0,3]} \delta_{\beta}^+(k, m) &= 1, & \sum_{m \in [0,3]} m \cdot \delta_{\beta}^+(k, m) &= \beta^+(k), & k &\in [1, k_C], \\ \sum_{m \in [0,3]} \delta_{\beta}^-(k, m) &= 1, & \sum_{m \in [0,3]} m \cdot \delta_{\beta}^-(k, m) &= \beta^-(k), & k &\in [1, k_C], \\ \sum_{m \in [0,3]} \delta_{\beta}^{\text{in}}(c, m) &= 1, & \sum_{m \in [0,3]} m \cdot \delta_{\beta}^{\text{in}}(c, m) &= \beta^{\text{in}}(c), & c &\in [1, c_F], \end{aligned} \quad (40)$$

$$\sum_{\psi \in \mathcal{F}_i^X} \beta_{\text{T}}([\psi]) \cdot \delta_{\text{fr}}^X(i, [\psi]) = \beta_{\text{ex}}^X(i), \quad i \in [1, t_X], X \in \{C, T, F\}, \quad (41)$$

$$\begin{aligned} \sum_{i \in [\widetilde{k}_C + 1, m_C]} \delta_{\beta}^C(i, m) &= \text{bd}_C(m), & \sum_{i \in [2, t_T]} \delta_{\beta}^T(i, m) &= \text{bd}_T(m), \\ \sum_{k \in [1, k_C]} \delta_{\beta}^+(k, m) &= \text{bd}_{CT}(m), & \sum_{k \in [1, k_C]} \delta_{\beta}^-(k, m) &= \text{bd}_{TC}(m), \\ \sum_{i \in [2, t_F]} \delta_{\beta}^F(i, m) &= \text{bd}_F(m), & \sum_{c \in [1, \widetilde{t}_C]} \delta_{\beta}^{\text{in}}(c, m) &= \text{bd}_{CF}(m), \\ & & \sum_{c \in [\widetilde{t}_C + 1, c_F]} \delta_{\beta}^{\text{in}}(c, m) &= \text{bd}_{TF}(m), \end{aligned}$$

$$\text{bd}_C(m) + \text{bd}_T(m) + \text{bd}_F(m) + \text{bd}_{CT}(m) + \text{bd}_{TC}(m) + \text{bd}_{TF}(m) + \text{bd}_{CF}(m) = \text{bd}^{\text{int}}(m), \quad m \in [1, 3]. \quad (42)$$

3.6 Assigning Chemical Elements and Valence Condition

We include constraints so that each vertex u in a selected graph H satisfies the valence condition; i.e., $\sum_{uv \in E(H)} \beta(uv) \leq \text{val}(\alpha(u))$. With these constraints, a chemical graph $G = (H, \alpha, \beta)$ on a selected subgraph H will be constructed.

constants:

- Subsets $\Lambda^{\text{int}}, \Lambda^{\text{ex}} \subseteq \Lambda$ of chemical elements, where we denote by $[\mathbf{e}]$ (resp., $[\mathbf{e}]^{\text{int}}$ and $[\mathbf{e}]^{\text{ex}}$) of a standard encoding of an element \mathbf{e} in the set Λ (resp., $\Lambda_{\mathbf{e}}^{\text{int}}$ and $\Lambda_{\mathbf{e}}^{\text{ex}}$);
- A valence function: $\text{val} : \Lambda \rightarrow [1, 4]$;
- A function $\text{mass}^* : \Lambda \rightarrow \mathbb{Z}$ (we let $\text{mass}(\mathbf{a})$ denote the observed mass of a chemical element $\mathbf{a} \in \Lambda$, and define $\text{mass}^*(\mathbf{a}) \triangleq \lfloor 10 \cdot \text{mass}(\mathbf{a}) \rfloor$);

- Subsets $\Lambda^*(i) \subseteq \Lambda^{\text{int}}, i \in [1, t_C]$;
- $\text{na}_{\text{LB}}(\mathbf{a}), \text{na}_{\text{UB}}(\mathbf{a}) \in [0, n^*], \mathbf{a} \in \Lambda$: lower and upper bounds on the number of vertices v with $\alpha(v) = \mathbf{a}$;
- $\text{na}_{\text{LB}}^{\text{int}}(\mathbf{a}), \text{na}_{\text{UB}}^{\text{int}}(\mathbf{a}) \in [0, n^*], \mathbf{a} \in \Lambda^{\text{int}}$: lower and upper bounds on the number of interior-vertices v with $\alpha(v) = \mathbf{a}$;
- $\alpha_r([\psi]) \in [\Lambda^{\text{ex}}], \in \mathcal{F}^* \cup \mathcal{F}_\Lambda$: the chemical element $\alpha(r)$ of the root r of ψ ;
- $\text{na}_{\mathbf{a}}^{\text{ex}}([\psi]) \in [0, n^*], \mathbf{a} \in \Lambda^{\text{ex}}, \psi \in \mathcal{F}^*$: the frequency of chemical element \mathbf{a} in the set of non-rooted vertices in ψ ;
- $n_{\text{H}}([\psi], d) \in [0, 3^\rho], \psi \in \mathcal{F}^* \cup \mathcal{F}_\Lambda, d \in [0, 3]$: the number of non-root vertices with $\text{deg}_{\text{hyd}}(v) = d$ in ψ .

variables:

- $\beta^{\text{CT}}(i), \beta^{\text{TC}}(i) \in [0, 3], i \in [1, t_T]$: the bond-multiplicity of edge $e^{\text{CT}}_{j,i}$ (resp., $e^{\text{TC}}_{j,i}$) if one exists;
- $\beta^{\text{CF}}(i), \beta^{\text{TF}}(i) \in [0, 3], i \in [1, t_F]$: the bond-multiplicity of $e^{\text{CF}}_{j,i}$ (resp., $e^{\text{TF}}_{j,i}$) if one exists;
- $\alpha^{\text{X}}(i) \in [\Lambda_\epsilon^{\text{int}}], \delta_\alpha^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) \in [0, 1], \mathbf{a} \in \Lambda_\epsilon^{\text{int}}, i \in [1, t_X], X \in \{\text{C}, \text{T}, \text{F}\}$: $\alpha^{\text{X}}(i) = [\mathbf{a}]^{\text{int}} \geq 1$ (resp., $\alpha^{\text{X}}(i) = 0$) $\Leftrightarrow \delta_\alpha^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) = 1$ (resp., $\delta_\alpha^{\text{X}}(i, 0) = 0$) $\Leftrightarrow \alpha(v^{\text{X}}_i) = \mathbf{a} \in \Lambda$ (resp., vertex v^{X}_i is not used in G);
- $\delta_\alpha^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) \in [0, 1], i \in [1, t_X], \mathbf{a} \in \Lambda^{\text{int}}, X \in \{\text{C}, \text{T}, \text{F}\}$: $\delta_\alpha^{\text{X}}(i, [\mathbf{a}]^{\text{t}}) = 1 \Leftrightarrow \alpha(v^{\text{X}}_i) = \mathbf{a}$;
- $\text{Mass} \in \mathbb{Z}_+$: $\sum_{v \in V(H)} \text{mass}^*(\alpha(v))$;
- $\text{na}([\mathbf{a}]) \in [\text{na}_{\text{LB}}(\mathbf{a}), \text{na}_{\text{UB}}(\mathbf{a})], \mathbf{a} \in \Lambda$: the number of vertices $v \in V(H)$ with $\alpha(v) = \mathbf{a}$;
- $\text{na}^{\text{int}}([\mathbf{a}]^{\text{int}}) \in [\text{na}_{\text{LB}}^{\text{int}}(\mathbf{a}), \text{na}_{\text{UB}}^{\text{int}}(\mathbf{a})], \mathbf{a} \in \Lambda, X \in \{\text{C}, \text{T}, \text{F}\}$: the number of interior-vertices $v \in V(G)$ with $\alpha(v) = \mathbf{a}$;
- $\text{na}_X^{\text{ex}}([\mathbf{a}]^{\text{ex}}), \text{na}^{\text{ex}}([\mathbf{a}]^{\text{ex}}) \in [0, \text{na}_{\text{UB}}(\mathbf{a})], \mathbf{a} \in \Lambda, X \in \{\text{C}, \text{T}, \text{F}\}$: the number of exterior-vertices rooted at vertices $v \in V_X$ and the number of exterior-vertices v such that $\alpha(v) = \mathbf{a}$;
- $\delta_{\text{hyd}}^{\text{X}}(i, d) \in [0, 1], d \in [0, 3], X \in \{\text{C}, \text{T}, \text{F}\}$: $\delta_{\text{hyd}}^{\text{X}}(i, d) \Leftrightarrow \text{deg}_{\text{hyd}}(v^{\text{X}}_i) = d$;
- $\text{hydg}(d), d \in [0, 3]$: the number of vertices v with $\text{deg}_{\text{hyd}}(v^{\text{X}}_i) = d$;

constraints:

$$\begin{aligned}
\beta^+(k) - 3(e^{\text{T}}(i) - \chi^{\text{T}}(i, k) + 1) &\leq \beta^{\text{CT}}(i) \leq \beta^+(k) + 3(e^{\text{T}}(i) - \chi^{\text{T}}(i, k) + 1), i \in [1, t_T], \\
\beta^-(k) - 3(e^{\text{T}}(i+1) - \chi^{\text{T}}(i, k) + 1) &\leq \beta^{\text{TC}}(i) \leq \beta^-(k) + 3(e^{\text{T}}(i+1) - \chi^{\text{T}}(i, k) + 1), i \in [1, t_T], \\
& k \in [1, k_C],
\end{aligned} \tag{43}$$

$$\begin{aligned}
\beta^{\text{in}}(c) - 3(e^{\text{F}}(i) - \chi^{\text{F}}(i, c) + 1) &\leq \beta^{\text{CF}}(i) \leq \beta^{\text{in}}(c) + 3(e^{\text{F}}(i) - \chi^{\text{F}}(i, c) + 1), i \in [1, t_F], \quad c \in [1, \tilde{t}_C], \\
\beta^{\text{in}}(c) - 3(e^{\text{F}}(i) - \chi^{\text{F}}(i, c) + 1) &\leq \beta^{\text{TF}}(i) \leq \beta^{\text{in}}(c) + 3(e^{\text{F}}(i) - \chi^{\text{F}}(i, c) + 1), i \in [1, t_F], \quad c \in [\tilde{t}_C + 1, c_F],
\end{aligned} \tag{44}$$

$$\begin{aligned} \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \delta_{\alpha}^{\text{C}}(i, [\mathbf{a}]^{\text{int}}) &= 1, & \sum_{\mathbf{a} \in \Lambda^{\text{int}}} [\mathbf{a}]^{\text{int}} \cdot \delta_{\alpha}^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) &= \alpha^{\text{C}}(i), & i \in [1, t_{\text{C}}], \\ \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \delta_{\alpha}^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) &= v^{\text{X}}(i), & \sum_{\mathbf{a} \in \Lambda^{\text{int}}} [\mathbf{a}]^{\text{int}} \cdot \delta_{\alpha}^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) &= \alpha^{\text{X}}(i), & i \in [1, t_{\text{X}}], \text{X} \in \{\text{T}, \text{F}\}, \end{aligned} \quad (45)$$

$$\sum_{\psi \in \mathcal{F}_i^{\text{X}} \cup \mathcal{F}_{\Lambda}} \alpha_{\text{r}}([\psi]) \cdot \delta_{\text{fr}}^{\text{X}}(i, [\psi]) = \alpha^{\text{X}}(i), \quad i \in [1, t_{\text{X}}], \text{X} \in \{\text{C}, \text{T}, \text{F}\}, \quad (46)$$

$$\begin{aligned} \sum_{j \in I_{\text{C}}(i)} \beta^{\text{C}}(j) + \sum_{k \in I_{(\geq 2)}^+(i) \cup I_{(\geq 1)}^+(i)} \beta^+(k) + \sum_{k \in I_{(\geq 2)}^-(i) \cup I_{(\geq 1)}^-(i)} \beta^-(k) \\ + \beta^{\text{in}}(i) + \beta_{\text{ex}}^{\text{C}}(i) + \sum_{d \in [0,3]} d \cdot \delta_{\text{hyd}}^{\text{C}}(i, d) &= \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \text{val}(\mathbf{a}) \delta_{\alpha}^{\text{C}}(i, [\mathbf{a}]^{\text{int}}), & i \in [1, \tilde{t}_{\text{C}}], \end{aligned} \quad (47)$$

$$\begin{aligned} \sum_{j \in I_{\text{C}}(i)} \beta^{\text{C}}(j) + \sum_{k \in I_{(\geq 2)}^+(i) \cup I_{(\geq 1)}^+(i)} \beta^+(k) + \sum_{k \in I_{(\geq 2)}^-(i) \cup I_{(\geq 1)}^-(i)} \beta^-(k) \\ + \beta_{\text{ex}}^{\text{C}}(i) + \sum_{d \in [0,3]} d \cdot \delta_{\text{hyd}}^{\text{C}}(i, d) &= \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \text{val}(\mathbf{a}) \delta_{\alpha}^{\text{C}}(i, [\mathbf{a}]^{\text{int}}), & i \in [\tilde{t}_{\text{C}} + 1, t_{\text{C}}], \end{aligned} \quad (48)$$

$$\begin{aligned} \beta^{\text{T}}(i) + \beta^{\text{T}}(i+1) + \beta_{\text{ex}}^{\text{T}}(i) + \beta^{\text{CT}}(i) + \beta^{\text{TC}}(i) \\ + \beta^{\text{in}}(\tilde{t}_{\text{C}} + i) + \sum_{d \in [0,3]} d \cdot \delta_{\text{hyd}}^{\text{T}}(i, d) &= \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \text{val}(\mathbf{a}) \delta_{\alpha}^{\text{T}}(i, [\mathbf{a}]^{\text{int}}), \\ i \in [1, t_{\text{T}}] \quad (\beta^{\text{T}}(1) = \beta^{\text{T}}(t_{\text{T}} + 1) = 0), & \end{aligned} \quad (49)$$

$$\begin{aligned} \beta^{\text{F}}(i) + \beta^{\text{F}}(i+1) + \beta^{\text{CF}}(i) + \beta^{\text{TF}}(i) \\ + \beta_{\text{ex}}^{\text{F}}(i) + \sum_{d \in [0,3]} d \cdot \delta_{\text{hyd}}^{\text{F}}(i, d) &= \sum_{\mathbf{a} \in \Lambda^{\text{int}}} \text{val}(\mathbf{a}) \delta_{\alpha}^{\text{F}}(i, [\mathbf{a}]^{\text{int}}), \\ i \in [1, t_{\text{F}}] \quad (\beta^{\text{F}}(1) = \beta^{\text{F}}(t_{\text{F}} + 1) = 0), & \end{aligned} \quad (50)$$

$$\sum_{i \in [1, t_{\text{X}}]} \delta_{\alpha}^{\text{X}}(i, [\mathbf{a}]^{\text{int}}) = \text{na}_{\text{X}}([\mathbf{a}]^{\text{int}}), \quad \mathbf{a} \in \Lambda^{\text{int}}, \text{X} \in \{\text{C}, \text{T}, \text{F}\}, \quad (51)$$

$$\sum_{\psi \in \mathcal{F}_i^{\text{X}}} \text{na}_{\mathbf{a}}^{\text{ex}}([\psi]) \cdot \delta_{\text{fr}}^{\text{X}}(i, [\psi]) = \text{na}_{\text{X}}^{\text{ex}}([\mathbf{a}]^{\text{ex}}), \quad \mathbf{a} \in \Lambda^{\text{ex}}, \text{X} \in \{\text{C}, \text{T}, \text{F}\}, \quad (52)$$

$$\begin{aligned}
na_C([\mathbf{a}]^{\text{int}}) + na_T([\mathbf{a}]^{\text{int}}) + na_F([\mathbf{a}]^{\text{int}}) &= na^{\text{int}}([\mathbf{a}]^{\text{int}}), & \mathbf{a} \in \Lambda^{\text{int}}, \\
\sum_{X \in \{C, T, F\}} na_X^{\text{ex}}([\mathbf{a}]^{\text{ex}}) &= na^{\text{ex}}([\mathbf{a}]^{\text{ex}}), & \mathbf{a} \in \Lambda^{\text{ex}}, \\
na^{\text{int}}([\mathbf{a}]^{\text{int}}) + na^{\text{ex}}([\mathbf{a}]^{\text{ex}}) &= na([\mathbf{a}]), & \mathbf{a} \in \Lambda^{\text{int}} \cap \Lambda^{\text{ex}}, \\
na^{\text{int}}([\mathbf{a}]^{\text{int}}) &= na([\mathbf{a}]), & \mathbf{a} \in \Lambda^{\text{int}} \setminus \Lambda^{\text{ex}}, \\
na^{\text{ex}}([\mathbf{a}]^{\text{ex}}) &= na([\mathbf{a}]), & \mathbf{a} \in \Lambda^{\text{ex}} \setminus \Lambda^{\text{int}},
\end{aligned} \tag{53}$$

$$\sum_{\mathbf{a} \in \Lambda} \text{mass}^*(\mathbf{a}) \cdot na([\mathbf{a}]) = \text{Mass}, \tag{54}$$

$$\begin{aligned}
\sum_{d \in [0, 3]} \delta_{\text{hyd}}^C(i, d) &= 1, \quad i \in [1, t_C], \\
\sum_{d \in [0, 3]} \delta_{\text{hyd}}^X(i, d) &= v^X(i), \quad i \in [1, t_X], X \in \{T, F\},
\end{aligned} \tag{55}$$

$$\sum_{i \in [1, t_X], X \in \{C, T, F\}} \delta_{\text{hyd}}^X(i, d) + \sum_{\psi \in \mathcal{F}_i^X, i \in [1, t_X], X \in \{C, T, F\}} n_{\text{H}}([\psi], d) \cdot \delta_{\text{fr}}^X(i, [\psi]) = \text{hydg}(d), \quad d \in [0, 3], \tag{56}$$

$$\sum_{\mathbf{a} \in \Lambda^*(i)} \delta_{\alpha}^C(i, [\mathbf{a}]^{\text{int}}) = 1, \quad i \in [1, t_C]. \tag{57}$$

3.7 Constraints for Bounds on the Number of Bonds

We include constraints for specification of lower and upper bounds bd_{LB} and bd_{UB} .

constants:

- $\text{bd}_{m, \text{LB}}(i), \text{bd}_{m, \text{UB}}(i) \in [0, n_{\text{UB}}^{\text{int}}]$, $i \in [1, m_C]$, $m \in [2, 3]$: lower and upper bounds on the number of edges $e \in E(P_i)$ with bond-multiplicity $\beta(e) = m$ in the pure path P_i for edge $e_i \in E_C$;

variables :

- $\text{bd}_T(k, i, m) \in [0, 1]$, $k \in [1, k_C]$, $i \in [2, t_T]$, $m \in [2, 3]$: $\text{bd}_T(k, i, m) = 1 \Leftrightarrow$ the pure path P_k for edge $e_k \in E_C$ contains edge e^T_i with $\beta(e^T_i) = m$;

constraints:

$$\text{bd}_{m, \text{LB}}(i) \leq \delta_{\beta}^C(i, m) \leq \text{bd}_{m, \text{UB}}(i), \quad i \in I_{(=1)} \cup I_{(0/1)}, \quad m \in [2, 3], \tag{58}$$

$$\text{bd}_T(k, i, m) \geq \delta_{\beta}^T(i, m) + \chi^T(i, k) - 1, \quad k \in [1, k_C], \quad i \in [2, t_T], \quad m \in [2, 3], \tag{59}$$

$$\sum_{j \in [2, t_T]} \delta_\beta^T(j, m) \geq \sum_{k \in [1, k_C], i \in [2, t_T]} \text{bd}_T(k, i, m), \quad m \in [2, 3], \quad (60)$$

$$\text{bd}_{m, \text{LB}}(k) \leq \sum_{i \in [2, t_T]} \text{bd}_T(k, i, m) + \delta_\beta^+(k, m) + \delta_\beta^-(k, m) \leq \text{bd}_{m, \text{UB}}(k), \quad (61)$$

$$k \in [1, k_C], m \in [2, 3].$$

3.8 Descriptor for the Number of Adjacency-configurations

We call a tuple $(\mathbf{a}, \mathbf{b}, m) \in \Lambda \times \Lambda \times [1, 3]$ an *adjacency-configuration*. The adjacency-configuration of an edge-configuration $(\mu = \mathbf{ad}, \mu' = \mathbf{bd}', m)$ is defined to be $(\mathbf{a}, \mathbf{b}, m)$. We include constraints to compute the frequency of each adjacency-configuration in an inferred chemical graph G .

constants:

- A set Γ^{int} of edge-configurations $\gamma = (\mu, \xi, m)$ with $\mu \leq \xi$;
- Let $\bar{\gamma}$ of an edge-configuration $\gamma = (\mu, \xi, m)$ denote the edge-configuration (ξ, μ, m) ;
- Let $\Gamma_{<}^{\text{int}} = \{(\mu, \xi, m) \in \Gamma^{\text{int}} \mid \mu < \xi\}$, $\Gamma_{=}^{\text{int}} = \{(\mu, \xi, m) \in \Gamma^{\text{int}} \mid \mu = \xi\}$ and $\Gamma_{>}^{\text{int}} = \{\bar{\gamma} \mid \gamma \in \Gamma_{<}^{\text{int}}\}$;
- Let $\Gamma_{\text{ac}, <}^{\text{int}}$, $\Gamma_{\text{ac}, =}^{\text{int}}$ and $\Gamma_{\text{ac}, >}^{\text{int}}$ denote the sets of the adjacency-configurations of edge-configurations in the sets $\Gamma_{<}^{\text{int}}$, $\Gamma_{=}^{\text{int}}$ and $\Gamma_{>}^{\text{int}}$, respectively;
- Let $\bar{\nu}$ of an adjacency-configuration $\nu = (\mathbf{a}, \mathbf{b}, m)$ denote the adjacency-configuration $(\mathbf{b}, \mathbf{a}, m)$;
- Prepare a coding of the set $\Gamma_{\text{ac}}^{\text{int}} \cup \Gamma_{\text{ac}, >}^{\text{int}}$ and let $[\nu]^{\text{int}}$ denote the coded integer of an element ν in $\Gamma_{\text{ac}}^{\text{int}} \cup \Gamma_{\text{ac}, >}^{\text{int}}$;
- Choose subsets $\tilde{\Gamma}_{\text{ac}}^{\text{C}}, \tilde{\Gamma}_{\text{ac}}^{\text{T}}, \tilde{\Gamma}_{\text{ac}}^{\text{CT}}, \tilde{\Gamma}_{\text{ac}}^{\text{TC}}, \tilde{\Gamma}_{\text{ac}}^{\text{F}}, \tilde{\Gamma}_{\text{ac}}^{\text{CF}}, \tilde{\Gamma}_{\text{ac}}^{\text{TF}} \subseteq \Gamma_{\text{ac}}^{\text{int}} \cup \Gamma_{\text{ac}, >}^{\text{int}}$; To compute the frequency of adjacency-configurations exactly, set $\tilde{\Gamma}_{\text{ac}}^{\text{C}} := \tilde{\Gamma}_{\text{ac}}^{\text{T}} := \tilde{\Gamma}_{\text{ac}}^{\text{CT}} := \tilde{\Gamma}_{\text{ac}}^{\text{TC}} := \tilde{\Gamma}_{\text{ac}}^{\text{F}} := \tilde{\Gamma}_{\text{ac}}^{\text{CF}} := \tilde{\Gamma}_{\text{ac}}^{\text{TF}} := \Gamma_{\text{ac}}^{\text{int}} \cup \Gamma_{\text{ac}, >}^{\text{int}}$;
- $\text{ac}_{\text{LB}}^{\text{int}}(\nu), \text{ac}_{\text{UB}}^{\text{int}}(\nu) \in [0, 2n_{\text{UB}}^{\text{int}}], \nu = (\mathbf{a}, \mathbf{b}, m) \in \Gamma_{\text{ac}}^{\text{int}}$: lower and upper bounds on the number of interior-edges $e = uv$ with $\alpha(u) = \mathbf{a}$, $\alpha(v) = \mathbf{b}$ and $\beta(e) = m$;

variables:

- $\text{ac}^{\text{int}}([\nu]^{\text{int}}) \in [\text{ac}_{\text{LB}}^{\text{int}}(\nu), \text{ac}_{\text{UB}}^{\text{int}}(\nu)], \nu \in \Gamma_{\text{ac}}^{\text{int}}$: the number of interior-edges with adjacency-configuration ν ;
- $\text{ac}_{\text{C}}([\nu]^{\text{int}}) \in [0, m_{\text{C}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{C}}, \text{ac}_{\text{T}}([\nu]^{\text{int}}) \in [0, t_{\text{T}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{T}}, \text{ac}_{\text{F}}([\nu]^{\text{int}}) \in [0, t_{\text{F}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}$: the number of edges $e^{\text{C}} \in E_{\text{C}}$ (resp., edges $e^{\text{T}} \in E_{\text{T}}$ and edges $e^{\text{F}} \in E_{\text{F}}$) with adjacency-configuration ν ;
- $\text{ac}_{\text{CT}}([\nu]^{\text{int}}) \in [0, \min\{k_{\text{C}}, t_{\text{T}}\}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{CT}}, \text{ac}_{\text{TC}}([\nu]^{\text{int}}) \in [0, \min\{k_{\text{C}}, t_{\text{T}}\}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{TC}}, \text{ac}_{\text{CF}}([\nu]^{\text{int}}) \in [0, \tilde{t}_{\text{C}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}, \text{ac}_{\text{TF}}([\nu]^{\text{int}}) \in [0, t_{\text{T}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}$: the number of edges $e^{\text{CT}} \in E_{\text{CT}}$ (resp., edges $e^{\text{TC}} \in E_{\text{TC}}$ and edges $e^{\text{CF}} \in E_{\text{CF}}$ and $e^{\text{TF}} \in E_{\text{TF}}$) with adjacency-configuration ν ;
- $\delta_{\text{ac}}^{\text{C}}(i, [\nu]^{\text{int}}) \in [0, 1], i \in [\tilde{k}_{\text{C}} + 1, m_{\text{C}}] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}, \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{C}}, \delta_{\text{ac}}^{\text{T}}(i, [\nu]^{\text{int}}) \in [0, 1], i \in [2, t_{\text{T}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{T}}, \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}) \in [0, 1], i \in [2, t_{\text{F}}], \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}$: $\delta_{\text{ac}}^{\text{X}}(i, [\nu]^{\text{int}}) = 1 \Leftrightarrow$ edge e^{X}_i has adjacency-configuration ν ;

- $\delta_{ac}^{CT}(k, [\nu]^{int}), \delta_{ac}^{TC}(k, [\nu]^{int}) \in [0, 1], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}, \nu \in \tilde{\Gamma}_{ac}^{CT}$: $\delta_{ac}^{CT}(k, [\nu]^{int}) = 1$ (resp., $\delta_{ac}^{TC}(k, [\nu]^{int}) = 1$) \Leftrightarrow edge $e_{tail(k),j}^{CT}$ (resp., $e_{head(k),j}^{TC}$) for some $j \in [1, t_T]$ has adjacency-configuration ν ;
- $\delta_{ac}^{CF}(c, [\nu]^{int}) \in [0, 1], c \in [1, \tilde{t}_C], \nu \in \tilde{\Gamma}_{ac}^{CF}$: $\delta_{ac}^{CF}(c, [\nu]^{int}) = 1 \Leftrightarrow$ edge $e_{c,i}^{CF}$ for some $i \in [1, t_F]$ has adjacency-configuration ν ;
- $\delta_{ac}^{TF}(i, [\nu]^{int}) \in [0, 1], i \in [1, t_T], \nu \in \tilde{\Gamma}_{ac}^{TF}$: $\delta_{ac}^{TF}(i, [\nu]^{int}) = 1 \Leftrightarrow$ edge $e_{i,j}^{TF}$ for some $j \in [1, t_F]$ has adjacency-configuration ν ;
- $\alpha^{CT}(k), \alpha^{TC}(k) \in [0, |\Lambda^{int}|], k \in [1, k_C]$: $\alpha(v)$ of the edge $(v_{tail(k)}^C, v) \in E_{CT}$ (resp., $(v, v_{head(k)}^C) \in E_{TC}$) if any;
- $\alpha^{CF}(c) \in [0, |\Lambda^{int}|], c \in [1, \tilde{t}_C]$: $\alpha(v)$ of the edge $(v_c^C, v) \in E_{CF}$ if any;
- $\alpha^{TF}(i) \in [0, |\Lambda^{int}|], i \in [1, t_T]$: $\alpha(v)$ of the edge $(v_i^T, v) \in E_{TF}$ if any;
- $\Delta_{ac}^{C+}(i), \Delta_{ac}^{C-}(i) \in [0, |\Lambda^{int}|], i \in [\tilde{k}_C+1, m_C], \Delta_{ac}^{T+}(i), \Delta_{ac}^{T-}(i) \in [0, |\Lambda^{int}|], i \in [2, t_T], \Delta_{ac}^{F+}(i), \Delta_{ac}^{F-}(i) \in [0, |\Lambda^{int}|], i \in [2, t_F]$: $\Delta_{ac}^{X+}(i) = \Delta_{ac}^{X-}(i) = 0$ (resp., $\Delta_{ac}^{X+}(i) = \alpha(u)$ and $\Delta_{ac}^{X-}(i) = \alpha(v)$) \Leftrightarrow edge $e_{X,i}^X = (u, v) \in E_X$ is used in G (resp., $e_{X,i}^X \notin E(G)$);
- $\Delta_{ac}^{CT+}(k), \Delta_{ac}^{CT-}(k) \in [0, |\Lambda^{int}|], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$: $\Delta_{ac}^{CT+}(k) = \Delta_{ac}^{CT-}(k) = 0$ (resp., $\Delta_{ac}^{CT+}(k) = \alpha(u)$ and $\Delta_{ac}^{CT-}(k) = \alpha(v)$) \Leftrightarrow edge $e_{tail(k),j}^{CT} = (u, v) \in E_{CT}$ for some $j \in [1, t_T]$ is used in G (resp., otherwise);
- $\Delta_{ac}^{TC+}(k), \Delta_{ac}^{TC-}(k) \in [0, |\Lambda^{int}|], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}$: Analogous with $\Delta_{ac}^{CT+}(k)$ and $\Delta_{ac}^{CT-}(k)$;
- $\Delta_{ac}^{CF+}(c) \in [0, |\Lambda^{int}|], \Delta_{ac}^{CF-}(c) \in [0, |\Lambda^{int}|], c \in [1, \tilde{t}_C]$: $\Delta_{ac}^{CF+}(c) = \Delta_{ac}^{CF-}(c) = 0$ (resp., $\Delta_{ac}^{CF+}(c) = \alpha(u)$ and $\Delta_{ac}^{CF-}(c) = \alpha(v)$) \Leftrightarrow edge $e_{c,i}^{CF} = (u, v) \in E_{CF}$ for some $i \in [1, t_F]$ is used in G (resp., otherwise);
- $\Delta_{ac}^{TF+}(i) \in [0, |\Lambda^{int}|], \Delta_{ac}^{TF-}(i) \in [0, |\Lambda^{int}|], i \in [1, t_T]$: Analogous with $\Delta_{ac}^{CF+}(c)$ and $\Delta_{ac}^{CF-}(c)$;

constraints:

$$\begin{aligned}
ac_C([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^C, \\
ac_T([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^T, \\
ac_F([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^F, \\
ac_{CT}([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^{CT}, \\
ac_{TC}([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^{TC}, \\
ac_{CF}([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^{CF}, \\
ac_{TF}([\nu]^{int}) &= 0, & \nu &\in \Gamma_{ac}^{int} \setminus \tilde{\Gamma}_{ac}^{TF},
\end{aligned}$$

(62)

$$\begin{aligned}
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_C([\nu]^{\text{int}}) &= \sum_{i \in [\widetilde{k}_C+1, m_C]} \delta_\beta^C(i, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_T([\nu]^{\text{int}}) &= \sum_{i \in [2, t_T]} \delta_\beta^T(i, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_F([\nu]^{\text{int}}) &= \sum_{i \in [2, t_F]} \delta_\beta^F(i, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_{CT}([\nu]^{\text{int}}) &= \sum_{k \in [1, k_C]} \delta_\beta^+(k, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_{TC}([\nu]^{\text{int}}) &= \sum_{k \in [1, k_C]} \delta_\beta^-(k, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_{CF}([\nu]^{\text{int}}) &= \sum_{c \in [1, \widetilde{t}_C]} \delta_\beta^{\text{in}}(c, m), & m \in [1, 3], \\
\sum_{(a,b,m)=\nu \in \Gamma_{ac}^{\text{int}}} \text{ac}_{TF}([\nu]^{\text{int}}) &= \sum_{c \in [\widetilde{t}_C+1, c_F]} \delta_\beta^{\text{in}}(c, m), & m \in [1, 3],
\end{aligned} \tag{63}$$

$$\begin{aligned}
&\sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^C} m \cdot \delta_{ac}^C(i, [\nu]^{\text{int}}) = \beta^C(i), \\
\Delta_{ac}^{C+}(i) + \sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^C} [\mathbf{a}]^{\text{int}} \delta_{ac}^C(i, [\nu]^{\text{int}}) &= \alpha^C(\text{tail}(i)), \\
\Delta_{ac}^{C-}(i) + \sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^C} [\mathbf{b}]^{\text{int}} \delta_{ac}^C(i, [\nu]^{\text{int}}) &= \alpha^C(\text{head}(i)), \\
\Delta_{ac}^{C+}(i) + \Delta_{ac}^{C-}(i) &\leq 2|\Lambda^{\text{int}}|(1 - e^C(i)), & i \in [\widetilde{k}_C + 1, m_C], \\
\sum_{i \in [\widetilde{k}_C+1, m_C]} \delta_{ac}^C(i, [\nu]^{\text{int}}) &= \text{ac}_C([\nu]^{\text{int}}), & \nu \in \widetilde{\Gamma}_{ac}^C,
\end{aligned} \tag{64}$$

$$\begin{aligned}
&\sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^T} m \cdot \delta_{ac}^T(i, [\nu]^{\text{int}}) = \beta^T(i), \\
\Delta_{ac}^{T+}(i) + \sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^T} [\mathbf{a}]^{\text{int}} \delta_{ac}^T(i, [\nu]^{\text{int}}) &= \alpha^T(i-1), \\
\Delta_{ac}^{T-}(i) + \sum_{\nu=(a,b,m) \in \widetilde{\Gamma}_{ac}^T} [\mathbf{b}]^{\text{int}} \delta_{ac}^T(i, [\nu]^{\text{int}}) &= \alpha^T(i), \\
\Delta_{ac}^{T+}(i) + \Delta_{ac}^{T-}(i) &\leq 2|\Lambda^{\text{int}}|(1 - e^T(i)), & i \in [2, t_T], \\
\sum_{i \in [2, t_T]} \delta_{ac}^T(i, [\nu]^{\text{int}}) &= \text{ac}_T([\nu]^{\text{int}}), & \nu \in \widetilde{\Gamma}_{ac}^T,
\end{aligned} \tag{65}$$

$$\begin{aligned}
& \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}} m \cdot \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}) = \beta^{\text{F}}(i), \\
\Delta_{\text{ac}}^{\text{F}+}(i) + & \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}} [\mathbf{a}]^{\text{int}} \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}) = \alpha^{\text{F}}(i-1), \\
\Delta_{\text{ac}}^{\text{F}-}(i) + & \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}} [\mathbf{b}]^{\text{int}} \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}) = \alpha^{\text{F}}(i), \\
\Delta_{\text{ac}}^{\text{F}+}(i) + \Delta_{\text{ac}}^{\text{F}-}(i) & \leq 2|\Lambda^{\text{ex}}|(1 - e^{\text{F}}(i)), & i \in [2, t_{\text{F}}], \\
\sum_{i \in [2, t_{\text{F}}]} \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}) & = \text{ac}_{\text{F}}([\nu]^{\text{int}}), & \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{F}}, \tag{66}
\end{aligned}$$

$$\begin{aligned}
\alpha^{\text{T}}(i) + |\Lambda^{\text{int}}|(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i)) & \geq \alpha^{\text{CT}}(k), \\
\alpha^{\text{CT}}(k) \geq \alpha^{\text{T}}(i) - |\Lambda^{\text{int}}|(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i)), & i \in [1, t_{\text{T}}], \\
\sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CT}}} m \cdot \delta_{\text{ac}}^{\text{CT}}(k, [\nu]^{\text{int}}) & = \beta^+(k), \\
\Delta_{\text{ac}}^{\text{CT}+}(k) + \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CT}}} [\mathbf{a}]^{\text{int}} \delta_{\text{ac}}^{\text{CT}}(k, [\nu]^{\text{int}}) & = \alpha^{\text{C}}(\text{tail}(k)), \\
\Delta_{\text{ac}}^{\text{CT}-}(k) + \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CT}}} [\mathbf{b}]^{\text{int}} \delta_{\text{ac}}^{\text{CT}}(k, [\nu]^{\text{int}}) & = \alpha^{\text{CT}}(k), \\
\Delta_{\text{ac}}^{\text{CT}+}(k) + \Delta_{\text{ac}}^{\text{CT}-}(k) & \leq 2|\Lambda^{\text{int}}|(1 - \delta_{\chi}^{\text{T}}(k)), & k \in [1, k_{\text{C}}], \\
\sum_{k \in [1, k_{\text{C}}]} \delta_{\text{ac}}^{\text{CT}}(k, [\nu]^{\text{int}}) & = \text{ac}_{\text{CT}}([\nu]^{\text{int}}), & \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{CT}}, \tag{67}
\end{aligned}$$

$$\begin{aligned}
\alpha^{\text{T}}(i) + |\Lambda^{\text{int}}|(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i+1)) & \geq \alpha^{\text{TC}}(k), \\
\alpha^{\text{TC}}(k) \geq \alpha^{\text{T}}(i) - |\Lambda^{\text{int}}|(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i+1)), & i \in [1, t_{\text{T}}], \\
\sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TC}}} m \cdot \delta_{\text{ac}}^{\text{TC}}(k, [\nu]^{\text{int}}) & = \beta^-(k), \\
\Delta_{\text{ac}}^{\text{TC}+}(k) + \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TC}}} [\mathbf{a}]^{\text{int}} \delta_{\text{ac}}^{\text{TC}}(k, [\nu]^{\text{int}}) & = \alpha^{\text{TC}}(k), \\
\Delta_{\text{ac}}^{\text{TC}-}(k) + \sum_{\nu=(\mathbf{a},\mathbf{b},m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TC}}} [\mathbf{b}]^{\text{int}} \delta_{\text{ac}}^{\text{TC}}(k, [\nu]^{\text{int}}) & = \alpha^{\text{C}}(\text{head}(k)), \\
\Delta_{\text{ac}}^{\text{TC}+}(k) + \Delta_{\text{ac}}^{\text{TC}-}(k) & \leq 2|\Lambda^{\text{int}}|(1 - \delta_{\chi}^{\text{T}}(k)), & k \in [1, k_{\text{C}}], \\
\sum_{k \in [1, k_{\text{C}}]} \delta_{\text{ac}}^{\text{TC}}(k, [\nu]^{\text{int}}) & = \text{ac}_{\text{TC}}([\nu]^{\text{int}}), & \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{TC}}, \tag{68}
\end{aligned}$$

$$\begin{aligned}
\alpha^F(i) + |\Lambda^{\text{int}}|(1 - \chi^F(i, c) + e^F(i)) &\geq \alpha^{\text{CF}}(c), \\
\alpha^{\text{CF}}(c) &\geq \alpha^F(i) - |\Lambda^{\text{int}}|(1 - \chi^F(i, c) + e^F(i)), & i \in [1, t_F], \\
\sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}} m \cdot \delta_{\text{ac}}^{\text{CF}}(c, [\nu]^{\text{int}}) &= \beta^{\text{in}}(c), \\
\Delta_{\text{ac}}^{\text{CF}+}(c) + \sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}} [\mathbf{a}]^{\text{int}} \delta_{\text{ac}}^{\text{CF}}(c, [\nu]^{\text{int}}) &= \alpha^{\text{C}}(\text{head}(c)), \\
\Delta_{\text{ac}}^{\text{CF}-}(c) + \sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}} [\mathbf{b}]^{\text{int}} \delta_{\text{ac}}^{\text{CF}}(c, [\nu]^{\text{int}}) &= \alpha^{\text{CF}}(c), \\
\Delta_{\text{ac}}^{\text{CF}+}(c) + \Delta_{\text{ac}}^{\text{CF}-}(c) &\leq 2 \max\{|\Lambda^{\text{int}}|, |\Lambda^{\text{int}}|\}(1 - \delta_X^F(c)), & c \in [1, \tilde{t}_C], \\
\sum_{c \in [1, \tilde{t}_C]} \delta_{\text{ac}}^{\text{CF}}(c, [\nu]^{\text{int}}) &= \text{ac}_{\text{CF}}([\nu]^{\text{int}}), & \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}, \quad (69)
\end{aligned}$$

$$\begin{aligned}
\alpha^F(j) + |\Lambda^{\text{int}}|(1 - \chi^F(j, i + \tilde{t}_C) + e^F(j)) &\geq \alpha^{\text{TF}}(i), \\
\alpha^{\text{TF}}(i) &\geq \alpha^F(j) - |\Lambda^{\text{int}}|(1 - \chi^F(j, i + \tilde{t}_C) + e^F(j)), & j \in [1, t_F], \\
\sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}} m \cdot \delta_{\text{ac}}^{\text{TF}}(i, [\nu]^{\text{int}}) &= \beta^{\text{in}}(i + \tilde{t}_C), \\
\Delta_{\text{ac}}^{\text{TF}+}(i) + \sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}} [\mathbf{a}]^{\text{int}} \delta_{\text{ac}}^{\text{TF}}(i, [\nu]^{\text{int}}) &= \alpha^{\text{T}}(i), \\
\Delta_{\text{ac}}^{\text{TF}-}(i) + \sum_{\nu=(\mathbf{a}, \mathbf{b}, m) \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}} [\mathbf{b}]^{\text{int}} \delta_{\text{ac}}^{\text{TF}}(i, [\nu]^{\text{int}}) &= \alpha^{\text{TF}}(i), \\
\Delta_{\text{ac}}^{\text{TF}+}(i) + \Delta_{\text{ac}}^{\text{TF}-}(i) &\leq 2 \max\{|\Lambda^{\text{int}}|, |\Lambda^{\text{int}}|\}(1 - \delta_X^F(i + \tilde{t}_C)), & i \in [1, t_T], \\
\sum_{i \in [1, t_T]} \delta_{\text{ac}}^{\text{TF}}(i, [\nu]^{\text{int}}) &= \text{ac}_{\text{TF}}([\nu]^{\text{int}}), & \nu \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}, \quad (70)
\end{aligned}$$

$$\begin{aligned}
\sum_{X \in \{\text{C}, \text{T}, \text{F}, \text{CT}, \text{TC}, \text{CF}, \text{TF}\}} (\text{ac}_X([\nu]^{\text{int}}) + \text{ac}_X([\bar{\nu}]^{\text{int}})) &= \text{ac}^{\text{int}}([\nu]^{\text{int}}), & \nu \in \Gamma_{\text{ac}, <}^{\text{int}}, \\
\sum_{X \in \{\text{C}, \text{T}, \text{F}, \text{CT}, \text{TC}, \text{CF}, \text{TF}\}} \text{ac}_X([\nu]^{\text{int}}) &= \text{ac}^{\text{int}}([\nu]^{\text{int}}), & \nu \in \Gamma_{\text{ac}, =}^{\text{int}}. \quad (71)
\end{aligned}$$

3.9 Descriptor for the Number of Chemical Symbols

We include constraints for computing the frequency of each chemical symbol in Λ_{dg} . Let $\text{cs}(v)$ denote the chemical symbol of a vertex v in a chemical graph G to be inferred; i.e., $\text{cs}(v) = \mu = \mathbf{ad} \in \Lambda_{\text{dg}}$ such that $\alpha(v) = \mathbf{a}$ and $\text{deg}_G(v) = d$.

constants:

- A set $\Lambda_{\text{dg}}^{\text{int}}$ of chemical symbols;
- Prepare a coding of each of the two sets $\Lambda_{\text{dg}}^{\text{int}}$ and let $[\mu]^{\text{int}}$ denote the coded integer of an element $\mu \in \Lambda_{\text{dg}}^{\text{int}}$;

- Choose subsets $\tilde{\Lambda}_{\text{dg}}^{\text{C}}, \tilde{\Lambda}_{\text{dg}}^{\text{T}}, \tilde{\Lambda}_{\text{dg}}^{\text{F}} \subseteq \Lambda_{\text{dg}}^{\text{int}}$: To compute the frequency of chemical symbols exactly, set $\tilde{\Lambda}_{\text{dg}}^{\text{C}} := \tilde{\Lambda}_{\text{dg}}^{\text{T}} := \tilde{\Lambda}_{\text{dg}}^{\text{F}} := \Lambda_{\text{dg}}^{\text{int}}$;

variables:

- $\text{ns}^{\text{int}}([\mu]^{\text{int}}) \in [0, n_{\text{UB}}^{\text{int}}]$, $\mu \in \Lambda_{\text{dg}}^{\text{int}}$: the number of interior-vertices v with $\text{cs}(v) = \mu$;
- $\delta_{\text{ns}}^{\text{X}}(i, [\mu]^{\text{int}}) \in [0, 1]$, $i \in [1, t_{\text{X}}]$, $\mu \in \Lambda_{\text{dg}}^{\text{int}}$, $\text{X} \in \{\text{C}, \text{T}, \text{F}\}$;

constraints:

$$\begin{aligned} \sum_{\mu \in \tilde{\Lambda}_{\text{dg}}^{\text{X}} \cup \{\epsilon\}} \delta_{\text{ns}}^{\text{X}}(i, [\mu]^{\text{int}}) &= 1, & \sum_{\mu = \text{ad} \in \tilde{\Lambda}_{\text{dg}}^{\text{X}}} [\mathbf{a}]^{\text{int}} \cdot \delta_{\text{ns}}^{\text{X}}(i, [\mu]^{\text{int}}) &= \alpha^{\text{X}}(i), \\ \sum_{\mu = \text{ad} \in \tilde{\Lambda}_{\text{dg}}^{\text{X}}} d \cdot \delta_{\text{ns}}^{\text{X}}(i, [\mu]^{\text{int}}) &= \text{deg}^{\text{X}}(i), \\ & & i \in [1, t_{\text{X}}], \text{X} \in \{\text{C}, \text{T}, \text{F}\}, \end{aligned} \quad (72)$$

$$\sum_{i \in [1, t_{\text{C}}]} \delta_{\text{ns}}^{\text{C}}(i, [\mu]^{\text{int}}) + \sum_{i \in [1, t_{\text{T}}]} \delta_{\text{ns}}^{\text{T}}(i, [\mu]^{\text{int}}) + \sum_{i \in [1, t_{\text{F}}]} \delta_{\text{ns}}^{\text{F}}(i, [\mu]^{\text{int}}) = \text{ns}^{\text{int}}([\mu]^{\text{int}}), \quad \mu \in \Lambda_{\text{dg}}^{\text{int}}. \quad (73)$$

3.10 Descriptor for the Number of Edge-configurations

We include constraints to compute the frequency of each edge-configuration in an inferred chemical graph G .

constants:

- A set Γ^{int} of edge-configurations $\gamma = (\mu, \xi, m)$ with $\mu \leq \xi$;
- Let $\Gamma_{<}^{\text{int}} = \{(\mu, \xi, m) \in \Gamma^{\text{int}} \mid \mu < \xi\}$, $\Gamma_{=}^{\text{int}} = \{(\mu, \xi, m) \in \Gamma^{\text{int}} \mid \mu = \xi\}$ and $\Gamma_{>}^{\text{int}} = \{(\xi, \mu, m) \mid (\mu, \xi, m) \in \Gamma_{<}^{\text{int}}\}$;
- Prepare a coding of the set $\Gamma^{\text{int}} \cup \Gamma_{>}^{\text{int}}$ and let $[\gamma]^{\text{int}}$ denote the coded integer of an element γ in $\Gamma^{\text{int}} \cup \Gamma_{>}^{\text{int}}$;
- Choose subsets $\tilde{\Gamma}_{\text{ec}}^{\text{C}}, \tilde{\Gamma}_{\text{ec}}^{\text{T}}, \tilde{\Gamma}_{\text{ec}}^{\text{CT}}, \tilde{\Gamma}_{\text{ec}}^{\text{TC}}, \tilde{\Gamma}_{\text{ec}}^{\text{F}}, \tilde{\Gamma}_{\text{ec}}^{\text{CF}}, \tilde{\Gamma}_{\text{ec}}^{\text{TF}} \subseteq \Gamma^{\text{int}} \cup \Gamma_{>}^{\text{int}}$; To compute the frequency of edge-configurations exactly, set $\tilde{\Gamma}_{\text{ec}}^{\text{C}} := \tilde{\Gamma}_{\text{ec}}^{\text{T}} := \tilde{\Gamma}_{\text{ec}}^{\text{CT}} := \tilde{\Gamma}_{\text{ec}}^{\text{TC}} := \tilde{\Gamma}_{\text{ec}}^{\text{F}} := \tilde{\Gamma}_{\text{ec}}^{\text{CF}} := \tilde{\Gamma}_{\text{ec}}^{\text{TF}} := \Gamma^{\text{int}} \cup \Gamma_{>}^{\text{int}}$;
- $\text{ec}_{\text{LB}}^{\text{int}}(\gamma), \text{ec}_{\text{UB}}^{\text{int}}(\gamma) \in [0, 2n_{\text{UB}}^{\text{int}}]$, $\gamma = (\mu, \xi, m) \in \Gamma^{\text{int}}$: lower and upper bounds on the number of interior-edges $e = uv$ with $\text{cs}(u) = \mu$, $\text{cs}(v) = \xi$ and $\beta(e) = m$;

variables:

- $\text{ec}^{\text{int}}([\gamma]^{\text{int}}) \in [\text{ec}_{\text{LB}}^{\text{int}}(\gamma), \text{ec}_{\text{UB}}^{\text{int}}(\gamma)]$, $\gamma \in \Gamma^{\text{int}}$: the number of interior-edges with edge-configuration γ ;
- $\text{ec}_{\text{C}}([\gamma]^{\text{int}}) \in [0, m_{\text{C}}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{C}}$, $\text{ec}_{\text{T}}([\gamma]^{\text{int}}) \in [0, t_{\text{T}}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{T}}$, $\text{ec}_{\text{F}}([\gamma]^{\text{int}}) \in [0, t_{\text{F}}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{F}}$: the number of edges $e^{\text{C}} \in E_{\text{C}}$ (resp., edges $e^{\text{T}} \in E_{\text{T}}$ and edges $e^{\text{F}} \in E_{\text{F}}$) with edge-configuration γ ;
- $\text{ec}_{\text{CT}}([\gamma]^{\text{int}}) \in [0, \min\{k_{\text{C}}, t_{\text{T}}\}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{CT}}$, $\text{ec}_{\text{TC}}([\gamma]^{\text{int}}) \in [0, \min\{k_{\text{C}}, t_{\text{T}}\}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{TC}}$, $\text{ec}_{\text{CF}}([\gamma]^{\text{int}}) \in [0, t_{\text{C}}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{CF}}$, $\text{ec}_{\text{TF}}([\gamma]^{\text{int}}) \in [0, t_{\text{T}}]$, $\gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{TF}}$: the number of edges $e^{\text{CT}} \in E_{\text{CT}}$ (resp., edges $e^{\text{TC}} \in E_{\text{TC}}$ and edges $e^{\text{CF}} \in E_{\text{CF}}$ and $e^{\text{TF}} \in E_{\text{TF}}$) with edge-configuration γ ;

- $\delta_{ec}^C(i, [\gamma]^{int}) \in [0, 1], i \in [\widetilde{k}_C + 1, m_C] = I_{(\geq 1)} \cup I_{(0/1)} \cup I_{(=1)}, \gamma \in \widetilde{\Gamma}_{ec}^C, \delta_{ec}^T(i, [\gamma]^{int}) \in [0, 1], i \in [2, t_T], \gamma \in \widetilde{\Gamma}_{ec}^T, \delta_{ec}^F(i, [\gamma]^{int}) \in [0, 1], i \in [2, t_F], \gamma \in \widetilde{\Gamma}_{ec}^F: \delta_{ec}^X(i, [\gamma]^t) = 1 \Leftrightarrow \text{edge } e^X_i \text{ has edge-configuration } \gamma;$
- $\delta_{ec,C}^{CT}(k, [\gamma]^{int}), \delta_{ec,C}^{TC}(k, [\gamma]^{int}) \in [0, 1], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}, \gamma \in \widetilde{\Gamma}_{ec}^{CT}: \delta_{ec,C}^{CT}(k, [\gamma]^{int}) = 1$ (resp., $\delta_{ec,C}^{TC}(k, [\gamma]^{int}) = 1$) $\Leftrightarrow \text{edge } e^{CT}_{tail(k),j}$ (resp., $e^{TC}_{head(k),j}$) for some $j \in [1, t_T]$ has edge-configuration $\gamma;$
- $\delta_{ec,C}^{CF}(c, [\gamma]^{int}) \in [0, 1], c \in [1, \widetilde{t}_C], \gamma \in \widetilde{\Gamma}_{ec}^{CF}: \delta_{ec,C}^{CF}(c, [\gamma]^{int}) = 1 \Leftrightarrow \text{edge } e^{CF}_{c,i}$ for some $i \in [1, t_F]$ has edge-configuration $\gamma;$
- $\delta_{ec,T}^{TF}(i, [\gamma]^{int}) \in [0, 1], i \in [1, t_T], \gamma \in \widetilde{\Gamma}_{ec}^{TF}: \delta_{ec,T}^{TF}(i, [\gamma]^{int}) = 1 \Leftrightarrow \text{edge } e^{TF}_{i,j}$ for some $j \in [1, t_F]$ has edge-configuration $\gamma;$
- $\text{deg}_T^{CT}(k), \text{deg}_T^{TC}(k) \in [0, 4], k \in [1, k_C]: \text{deg}_G(v)$ of an end-vertex $v \in V_T$ of the edge $(v^C_{tail(k)}, v) \in E_{CT}$ (resp., $(v, v^C_{head(k)}) \in E_{TC}$) if any;
- $\text{deg}_F^{CF}(c) \in [0, 4], c \in [1, \widetilde{t}_C]: \text{deg}_G(v)$ of an end-vertex $v \in V_F$ of the edge $(v^C_c, v) \in E_{CF}$ if any;
- $\text{deg}_F^{TF}(i) \in [0, 4], i \in [1, t_T]: \text{deg}_G(v)$ of an end-vertex $v \in V_F$ of the edge $(v^T_i, v) \in E_{TF}$ if any;
- $\Delta_{ec}^{C+}(i), \Delta_{ec}^{C-}(i) \in [0, 4], i \in [\widetilde{k}_C + 1, m_C], \Delta_{ec}^{T+}(i), \Delta_{ec}^{T-}(i) \in [0, 4], i \in [2, t_T], \Delta_{ec}^{F+}(i), \Delta_{ec}^{F-}(i) \in [0, 4], i \in [2, t_F]: \Delta_{ec}^{X+}(i) = \Delta_{ec}^{X-}(i) = 0$ (resp., $\Delta_{ec}^{X+}(i) = \text{deg}_G(u)$ and $\Delta_{ec}^{X-}(i) = \text{deg}_G(v)$) $\Leftrightarrow \text{edge } e^X_i = (u, v) \in E_X$ is used in G (resp., $e^X_i \notin E(G)$);
- $\Delta_{ec}^{CT+}(k), \Delta_{ec}^{CT-}(k) \in [0, 4], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}: \Delta_{ec}^{CT+}(k) = \Delta_{ec}^{CT-}(k) = 0$ (resp., $\Delta_{ec}^{CT+}(k) = \text{deg}_G(u)$ and $\Delta_{ec}^{CT-}(k) = \text{deg}_G(v)$) $\Leftrightarrow \text{edge } e^{CT}_{tail(k),j} = (u, v) \in E_{CT}$ for some $j \in [1, t_T]$ is used in G (resp., otherwise);
- $\Delta_{ec}^{TC+}(k), \Delta_{ec}^{TC-}(k) \in [0, 4], k \in [1, k_C] = I_{(\geq 2)} \cup I_{(\geq 1)}: \text{Analogous with } \Delta_{ec}^{CT+}(k) \text{ and } \Delta_{ec}^{CT-}(k);$
- $\Delta_{ec}^{CF+}(c), \Delta_{ec}^{CF-}(c) \in [0, 4], c \in [1, \widetilde{t}_C]: \Delta_{ec}^{CF+}(c) = \Delta_{ec}^{CF-}(c) = 0$ (resp., $\Delta_{ec}^{CF+}(c) = \text{deg}_G(u)$ and $\Delta_{ec}^{CF-}(c) = \text{deg}_G(v)$) $\Leftrightarrow \text{edge } e^{CF}_{c,j} = (u, v) \in E_{CF}$ for some $j \in [1, t_F]$ is used in G (resp., otherwise);
- $\Delta_{ec}^{TF+}(i), \Delta_{ec}^{TF-}(i) \in [0, 4], i \in [1, t_T]: \text{Analogous with } \Delta_{ec}^{CF+}(c) \text{ and } \Delta_{ec}^{CF-}(c);$

constraints:

$$\begin{aligned}
ec_C([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^C, \\
ec_T([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^T, \\
ec_F([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^F, \\
ec_{CT}([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^{CT}, \\
ec_{TC}([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^{TC}, \\
ec_{CF}([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^{CF}, \\
ec_{TF}([\gamma]^{int}) &= 0, & \gamma &\in \Gamma^{int} \setminus \widetilde{\Gamma}_{ec}^{TF},
\end{aligned}$$

(74)

$$\begin{aligned}
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_C([\gamma]^{\text{int}}) &= \sum_{i \in [\widetilde{k}_C + 1, m_C]} \delta_\beta^C(i, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_T([\gamma]^{\text{int}}) &= \sum_{i \in [2, t_T]} \delta_\beta^T(i, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_F([\gamma]^{\text{int}}) &= \sum_{i \in [2, t_F]} \delta_\beta^F(i, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_{CT}([\gamma]^{\text{int}}) &= \sum_{k \in [1, k_C]} \delta_\beta^+(k, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_{TC}([\gamma]^{\text{int}}) &= \sum_{k \in [1, k_C]} \delta_\beta^-(k, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_{CF}([\gamma]^{\text{int}}) &= \sum_{c \in [1, \widetilde{t}_C]} \delta_\beta^{\text{in}}(c, m), & m \in [1, 3], \\
\sum_{(\mu, \mu', m) = \gamma \in \Gamma^{\text{int}}} \text{ec}_{TF}([\gamma]^{\text{int}}) &= \sum_{c \in [\widetilde{t}_C + 1, c_F]} \delta_\beta^{\text{in}}(c, m), & m \in [1, 3],
\end{aligned} \tag{75}$$

$$\begin{aligned}
\sum_{\gamma = (\mathbf{ad}, \mathbf{bd}', m) \in \widetilde{\Gamma}_{\text{ec}}^C} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec}}^C(i, [\gamma]^{\text{int}}) &= \sum_{\nu \in \widetilde{\Gamma}_{\text{ac}}^C} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^C(i, [\nu]^{\text{int}}), \\
\Delta_{\text{ec}}^{C+}(i) + \sum_{\gamma = (\mathbf{ad}, \xi, m) \in \widetilde{\Gamma}_{\text{ec}}^C} d \cdot \delta_{\text{ec}}^C(i, [\gamma]^{\text{int}}) &= \text{deg}^C(\text{tail}(i)), \\
\Delta_{\text{ec}}^{C-}(i) + \sum_{\gamma = (\mu, \mathbf{bd}, m) \in \widetilde{\Gamma}_{\text{ec}}^C} d \cdot \delta_{\text{ec}}^C(i, [\gamma]^{\text{int}}) &= \text{deg}^C(\text{head}(i)), \\
\Delta_{\text{ec}}^{C+}(i) + \Delta_{\text{ec}}^{C-}(i) &\leq 8(1 - e^C(i)), & i \in [\widetilde{k}_C + 1, m_C], \\
\sum_{i \in [\widetilde{k}_C + 1, m_C]} \delta_{\text{ec}}^C(i, [\gamma]^{\text{int}}) &= \text{ec}_C([\gamma]^{\text{int}}), & \gamma \in \widetilde{\Gamma}_{\text{ec}}^C,
\end{aligned} \tag{76}$$

$$\begin{aligned}
\sum_{\gamma = (\mathbf{ad}, \mathbf{bd}', m) \in \widetilde{\Gamma}_{\text{ec}}^T} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec}}^T(i, [\gamma]^{\text{int}}) &= \sum_{\nu \in \widetilde{\Gamma}_{\text{ac}}^T} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^T(i, [\nu]^{\text{int}}), \\
\Delta_{\text{ec}}^{T+}(i) + \sum_{\gamma = (\mathbf{ad}, \xi, m) \in \widetilde{\Gamma}_{\text{ec}}^T} d \cdot \delta_{\text{ec}}^T(i, [\gamma]^{\text{int}}) &= \text{deg}^T(i - 1), \\
\Delta_{\text{ec}}^{T-}(i) + \sum_{\gamma = (\mu, \mathbf{bd}, m) \in \widetilde{\Gamma}_{\text{ec}}^T} d \cdot \delta_{\text{ec}}^T(i, [\gamma]^{\text{int}}) &= \text{deg}^T(i), \\
\Delta_{\text{ec}}^{T+}(i) + \Delta_{\text{ec}}^{T-}(i) &\leq 8(1 - e^T(i)), & i \in [2, t_T], \\
\sum_{i \in [2, t_T]} \delta_{\text{ec}}^T(i, [\gamma]^{\text{int}}) &= \text{ec}_T([\gamma]^{\text{int}}), & \gamma \in \widetilde{\Gamma}_{\text{ec}}^T,
\end{aligned} \tag{77}$$

$$\begin{aligned}
\sum_{\gamma=(\mathbf{ad},\mathbf{bd}',m)\in\tilde{\Gamma}_{\text{ec}}^{\text{F}}} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec}}^{\text{F}}(i, [\gamma]^{\text{int}}) &= \sum_{\nu\in\tilde{\Gamma}_{\text{ac}}^{\text{F}}} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^{\text{F}}(i, [\nu]^{\text{int}}), \\
\Delta_{\text{ec}}^{\text{F}+}(i) + \sum_{\gamma=(\mathbf{ad},\xi,m)\in\tilde{\Gamma}_{\text{ec}}^{\text{F}}} d \cdot \delta_{\text{ec}}^{\text{F}}(i, [\gamma]^{\text{int}}) &= \deg^{\text{F}}(i-1), \\
\Delta_{\text{ec}}^{\text{F}-}(i) + \sum_{\gamma=(\mu,\mathbf{bd},m)\in\tilde{\Gamma}_{\text{ec}}^{\text{F}}} d \cdot \delta_{\text{ec}}^{\text{F}}(i, [\gamma]^{\text{int}}) &= \deg^{\text{F}}(i), \\
\Delta_{\text{ec}}^{\text{F}+}(i) + \Delta_{\text{ec}}^{\text{F}-}(i) &\leq 8(1 - e^{\text{F}}(i)), & i \in [2, t_{\text{F}}], \\
\sum_{i\in[2,t_{\text{F}}]} \delta_{\text{ec}}^{\text{F}}(i, [\gamma]^{\text{int}}) &= \text{ec}_{\text{F}}([\gamma]^{\text{int}}), & \gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{F}}, \quad (78)
\end{aligned}$$

$$\begin{aligned}
\deg^{\text{T}}(i) + 4(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i)) &\geq \deg_{\text{T}}^{\text{CT}}(k), \\
\deg_{\text{T}}^{\text{CT}}(k) &\geq \deg^{\text{T}}(i) - 4(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i)), & i \in [1, t_{\text{T}}], \\
\sum_{\gamma=(\mathbf{ad},\mathbf{bd}',m)\in\tilde{\Gamma}_{\text{ec}}^{\text{CT}}} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec},\text{C}}^{\text{CT}}(k, [\gamma]^{\text{int}}) &= \sum_{\nu\in\tilde{\Gamma}_{\text{ac}}^{\text{CT}}} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^{\text{CT}}(k, [\nu]^{\text{int}}), \\
\Delta_{\text{ec}}^{\text{CT}+}(k) + \sum_{\gamma=(\mathbf{ad},\xi,m)\in\tilde{\Gamma}_{\text{ec}}^{\text{CT}}} d \cdot \delta_{\text{ec},\text{C}}^{\text{CT}}(k, [\gamma]^{\text{int}}) &= \deg^{\text{C}}(\text{tail}(k)), \\
\Delta_{\text{ec}}^{\text{CT}-}(k) + \sum_{\gamma=(\mu,\mathbf{bd},m)\in\tilde{\Gamma}_{\text{ec}}^{\text{CT}}} d \cdot \delta_{\text{ec},\text{C}}^{\text{CT}}(k, [\gamma]^{\text{int}}) &= \deg_{\text{T}}^{\text{CT}}(k), \\
\Delta_{\text{ec}}^{\text{CT}+}(k) + \Delta_{\text{ec}}^{\text{CT}-}(k) &\leq 8(1 - \delta_{\chi}^{\text{T}}(k)), & k \in [1, k_{\text{C}}], \\
\sum_{k\in[1,k_{\text{C}}]} \delta_{\text{ec},\text{C}}^{\text{CT}}(k, [\gamma]^{\text{int}}) &= \text{ec}_{\text{CT}}([\gamma]^{\text{int}}), & \gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{CT}}, \quad (79)
\end{aligned}$$

$$\begin{aligned}
\deg^{\text{T}}(i) + 4(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i+1)) &\geq \deg_{\text{T}}^{\text{TC}}(k), \\
\deg_{\text{T}}^{\text{TC}}(k) &\geq \deg^{\text{T}}(i) - 4(1 - \chi^{\text{T}}(i, k) + e^{\text{T}}(i+1)), & i \in [1, t_{\text{T}}], \\
\sum_{\gamma=(\mathbf{ad},\mathbf{bd}',m)\in\tilde{\Gamma}_{\text{ec}}^{\text{TC}}} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec},\text{C}}^{\text{TC}}(k, [\gamma]^{\text{int}}) &= \sum_{\nu\in\tilde{\Gamma}_{\text{ac}}^{\text{TC}}} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^{\text{TC}}(k, [\nu]^{\text{int}}), \\
\Delta_{\text{ec}}^{\text{TC}+}(k) + \sum_{\gamma=(\mathbf{ad},\xi,m)\in\tilde{\Gamma}_{\text{ec}}^{\text{TC}}} d \cdot \delta_{\text{ec},\text{C}}^{\text{TC}}(k, [\gamma]^{\text{int}}) &= \deg_{\text{T}}^{\text{TC}}(k), \\
\Delta_{\text{ec}}^{\text{TC}-}(k) + \sum_{\gamma=(\mu,\mathbf{bd},m)\in\tilde{\Gamma}_{\text{ec}}^{\text{TC}}} d \cdot \delta_{\text{ec},\text{C}}^{\text{TC}}(k, [\gamma]^{\text{int}}) &= \deg^{\text{C}}(\text{head}(k)), \\
\Delta_{\text{ec}}^{\text{TC}+}(k) + \Delta_{\text{ec}}^{\text{TC}-}(k) &\leq 8(1 - \delta_{\chi}^{\text{T}}(k)), & k \in [1, k_{\text{C}}], \\
\sum_{k\in[1,k_{\text{C}}]} \delta_{\text{ec},\text{C}}^{\text{TC}}(k, [\gamma]^{\text{int}}) &= \text{ec}_{\text{TC}}([\gamma]^{\text{int}}), & \gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{TC}}, \quad (80)
\end{aligned}$$

$$\begin{aligned}
& \deg^F(i) + 4(1 - \chi^F(i, c) + e^F(i)) \geq \deg_{\text{F}}^{\text{CF}}(c), \\
& \deg_{\text{F}}^{\text{CF}}(c) \geq \deg^F(i) - 4(1 - \chi^F(i, c) + e^F(i)), & i \in [1, t_{\text{F}}], \\
& \sum_{\gamma=(\mathbf{ad}, \mathbf{bd}', m) \in \tilde{\Gamma}_{\text{ec}}^{\text{CF}}} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec}, \text{C}}^{\text{CF}}(c, [\gamma]^{\text{int}}) = \sum_{\nu \in \tilde{\Gamma}_{\text{ac}}^{\text{CF}}} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^{\text{CF}}(c, [\nu]^{\text{int}}), \\
& \Delta_{\text{ec}}^{\text{CF}+}(c) + \sum_{\gamma=(\mathbf{ad}, \xi, m) \in \tilde{\Gamma}_{\text{ec}}^{\text{CF}}} d \cdot \delta_{\text{ec}, \text{C}}^{\text{CF}}(c, [\gamma]^{\text{int}}) = \deg^{\text{C}}(c), \\
& \Delta_{\text{ec}}^{\text{CF}-}(c) + \sum_{\gamma=(\mu, \mathbf{bd}, m) \in \tilde{\Gamma}_{\text{ec}}^{\text{CF}}} d \cdot \delta_{\text{ec}, \text{C}}^{\text{CF}}(c, [\gamma]^{\text{int}}) = \deg_{\text{F}}^{\text{CF}}(c), \\
& \Delta_{\text{ec}}^{\text{CF}+}(c) + \Delta_{\text{ec}}^{\text{CF}-}(c) \leq 8(1 - \delta_{\chi}^{\text{F}}(c)), & c \in [1, \tilde{t}_{\text{C}}], \\
& \sum_{c \in [1, \tilde{t}_{\text{C}}]} \delta_{\text{ec}, \text{C}}^{\text{CF}}(c, [\gamma]^{\text{int}}) = \text{ec}_{\text{CF}}([\gamma]^{\text{int}}), & \gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{CF}}, \quad (81)
\end{aligned}$$

$$\begin{aligned}
& \deg^F(j) + 4(1 - \chi^F(j, i + \tilde{t}_{\text{C}}) + e^F(j)) \geq \deg_{\text{F}}^{\text{TF}}(i), \\
& \deg_{\text{F}}^{\text{TF}}(i) \geq \deg^F(j) - 4(1 - \chi^F(j, i + \tilde{t}_{\text{C}}) + e^F(j)), & j \in [1, t_{\text{F}}], \\
& \sum_{\gamma=(\mathbf{ad}, \mathbf{bd}', m) \in \tilde{\Gamma}_{\text{ec}}^{\text{TF}}} [(\mathbf{a}, \mathbf{b}, m)]^{\text{int}} \cdot \delta_{\text{ec}, \text{T}}^{\text{TF}}(i, [\gamma]^{\text{int}}) = \sum_{\nu \in \tilde{\Gamma}_{\text{ac}}^{\text{TF}}} [\nu]^{\text{int}} \cdot \delta_{\text{ac}}^{\text{TF}}(i, [\nu]^{\text{int}}), \\
& \Delta_{\text{ec}}^{\text{TF}+}(i) + \sum_{\gamma=(\mathbf{ad}, \xi, m) \in \tilde{\Gamma}_{\text{ec}}^{\text{TF}}} d \cdot \delta_{\text{ec}, \text{T}}^{\text{TF}}(i, [\gamma]^{\text{int}}) = \deg^{\text{T}}(i), \\
& \Delta_{\text{ec}}^{\text{TF}-}(i) + \sum_{\gamma=(\mu, \mathbf{bd}, m) \in \tilde{\Gamma}_{\text{ec}}^{\text{TF}}} d \cdot \delta_{\text{ec}, \text{T}}^{\text{TF}}(i, [\gamma]^{\text{int}}) = \deg_{\text{F}}^{\text{TF}}(i), \\
& \Delta_{\text{ec}}^{\text{TF}+}(i) + \Delta_{\text{ec}}^{\text{TF}-}(i) \leq 8(1 - \delta_{\chi}^{\text{F}}(i + \tilde{t}_{\text{C}})), & i \in [1, t_{\text{T}}], \\
& \sum_{i \in [1, t_{\text{T}}]} \delta_{\text{ec}, \text{T}}^{\text{TF}}(i, [\gamma]^{\text{int}}) = \text{ec}_{\text{TF}}([\gamma]^{\text{int}}), & \gamma \in \tilde{\Gamma}_{\text{ec}}^{\text{TF}}, \quad (82)
\end{aligned}$$

$$\begin{aligned}
& \sum_{\text{X} \in \{\text{C}, \text{T}, \text{F}, \text{CT}, \text{TC}, \text{CF}, \text{TF}\}} (\text{ec}_{\text{X}}([\gamma]^{\text{int}}) + \text{ec}_{\text{X}}([\bar{\gamma}]^{\text{int}})) = \text{ec}^{\text{int}}([\gamma]^{\text{int}}), & \gamma \in \Gamma_{<}^{\text{int}}, \\
& \sum_{\text{X} \in \{\text{C}, \text{T}, \text{F}, \text{CT}, \text{TC}, \text{CF}, \text{TF}\}} \text{ec}_{\text{X}}([\gamma]^{\text{int}}) = \text{ec}^{\text{int}}([\gamma]^{\text{int}}), & \gamma \in \Gamma_{=}^{\text{int}}. \quad (83)
\end{aligned}$$

3.11 Descriptor for the Number of of Fringe-configurations

We include constraints to compute the frequency of each fringe-configuration in an inferred chemical graph G .

variables:

$\text{fc}([\psi]) \in [0, t_{\text{C}} + t_{\text{T}} + t_{\text{F}}]$, $\psi \in \mathcal{F}^*$: the frequency of a chemical rooted tree ψ in the set of ρ -fringe-trees in G ;

constraints:

$$\sum_{i \in [1, t_X], X \in \{C, T, F\}} \delta_{fr}^X(i, [\psi]) = fc([\psi]), \quad \psi \in \mathcal{F}^*. \quad (84)$$

3.12 Constraints for Normalization of Feature Vectors

By introducing a tolerance $\varepsilon > 0$ in the conversion between integers and reals, we include the following constraints for normalizing of a feature vector $f(G) = (x_1, x_2, \dots, x_K)$:

$$\frac{(1 - \varepsilon)(x_i - \min(\text{dcp}_i; D_\pi))}{\max(\text{dcp}_i; D_\pi) - \min(\text{dcp}_i; D_\pi)} \leq \hat{x}_i \leq \frac{(1 + \varepsilon)(x_i - \min(\text{dcp}_i; D_\pi))}{\max(\text{dcp}_i; D_\pi) - \min(\text{dcp}_i; D_\pi)}, \quad i \in [1, K]. \quad (85)$$

An example of a tolerance is $\varepsilon = 0.01$.

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