

## SUPPLEMENTARY MATERIALS

# Choosing Variant Interpretation Tools for Clinical Applications: Context Matters

Josu Aguirre <sup>1,†</sup>, Natàlia Padilla <sup>1,†</sup>, Selen Özkan <sup>1</sup>, Casandra Riera <sup>1</sup>, Lúdia Feliubadaló <sup>2,3</sup> and Xavier de la Cruz <sup>1,4,\*</sup>

<sup>1</sup> Research Unit in Clinical and Translational Bioinformatics, Vall d'Hebron Institute of Research (VHIR), Universitat Autònoma de Barcelona, P/Vall d'Hebron, 119-129, 08035 Barcelona, Spain; aguirre.gomez.josu@gmail.com (J.A.); natalia.padilla@vhir.org (N.P.); selen.ozkan@vhir.org (S.Ö.); mcasandraria@gmail.com (C.R.)

<sup>2</sup> Hereditary Cancer Program, Program in Molecular Mechanisms and Experimental Therapy in Oncology (Oncobell), IDIBELL, Catalan Institute of Oncology, 08908 L'Hospitalet de Llobregat, Spain; lfeliubadalo@iconcologia.net

<sup>3</sup> Centro de Investigación Biomédica en Red de Cáncer (CIBERONC), 28929 Madrid, Spain

<sup>4</sup> Institució Catalana de Recerca i Estudis Avançats (ICREA), 08010 Barcelona, Spain

\* Correspondence: xavier.delacruz@vhir.org; Tel.: +34-934-893-000 (ext. 2687)

† These authors contributed equally to this work.

## **Appendix S1. Description of the MISC+REJ cost model and proof of the main results underlying the clinical space search algorithm under this model**

In this work, we present a framework to simultaneously compare the performance of multiple pathogenicity predictors in multiple deployment scenarios. The goal is to find which predictor is more appropriate for each scenario, or if a single predictor is enough for all scenarios. We treat pathogenicity predictors as classifiers with reject option because they have an incomplete coverage [1]. Our framework consists of two components. The first component is a theoretical body that includes a cost model and a set of propositions forming the foundation for the second component, the computational part. This computational part presents a procedure that divides the clinical space into regions where a single predictor prevails over the others in terms of cost. Here, we present the cost model, and in Sections S1-S3, we explain the computational procedure and the theoretical outcomes that support it.

**NOTE.** In the main body of the article, for comparison purposes, we present a framework based on the split misclassification errors (MISC) only; the reject term is not included. All the results about MISC are provided in the article.

**TERMINOLOGY NOTE.** We will use the terms ‘classifier’ and ‘predictor’ interchangeably. We will also refer to these tools as ‘methods’, particularly in the propositions, to underline the generalizability of the results.

The cost model we use as a starting point is the one used to evaluate classifiers with a reject option [2,3], which corresponds to the case of pathogenicity predictors. However, this model must be adapted to healthcare applications because of two reasons. First, it treats equally misclassification errors due to false positives and false negatives, and this is problematic because the medical consequences of these errors may vary substantially [4]. And second, it should take into account the frequency of the predicted classes [5–7], i.e., the frequency of

pathogenic and benign variants, in our case. In the following, we present a version of the cost model for reject classifiers (to which we will refer as MISC+REJ) based on these considerations (see the underlying probabilistic tree diagram in Supplementary Figure S2). In particular, the misclassification error is split into two terms, corresponding to false positive and false negative errors, and a parameter for the class probability. The cost model is then written as:

$$c = \alpha\rho(1-s_e)c_0 + \alpha(1-\rho)(1-s_p)c_1 + (1-\alpha)c_2 \quad (\text{SA1})$$

where the parameters  $c_0$  and  $c_1$  are the misclassification costs:  $c_0$  is the cost associated to annotating pathogenic variants as benign, and  $c_1$  is the cost associated to annotating benign variants as pathogenic.  $c_2$  is the cost associated to prediction rejection. In general cost models,  $\rho$  and  $1-\rho$  are the probabilities of the two predicted classes. In our case,  $\rho$  will be the frequency of pathogenic variants expected in the deployment context. The value of  $\rho$ , comprised between 0 and 1, varies with the genome region sequenced (gene panel, whole exome, etc.) and with the population of individuals tested (e.g., the individuals attending a specific hospital unit, ethnic group, etc.).  $s_e$  and  $s_p$  are the sensitivity and specificity of the pathogenicity predictor;  $\alpha$  is its coverage ( $1-\alpha$  is the rejection rate). These parameters are estimated testing the predictor in a set of  $N_{tot}$  variants as follows:

$$s_e = \frac{TP}{N_p} \quad (\text{SA1.1})$$

$$s_p = \frac{TN}{N_b} \quad (\text{SA1.2})$$

$$\alpha = \frac{N}{N_{tot}} \quad (\text{SA1.3})$$

where  $N(N \leq N_{tot})$  is the total number of predicted variants;  $N_p$ , and  $N_b$  are the numbers of predicted pathogenic and benign variants. We have that:  $N = N_p + N_b$ . And  $N_{tot} - N$  is the number of rejected predictions.

Following Hernández-Orallo et al. [7], instead of  $c$  we will use  $rc$ , the normalized average cost, which is obtained after dividing both sides of equation (1) by  $c_T (=c_0+c_1+c_2)$ :

$$rc = \frac{c}{c_T} = \alpha\rho(1-s_e)rc_0 + \alpha(1-\rho)(1-s_p)rc_1 + (1-\alpha)rc_2 \quad (SA2)$$

where  $rc_i = c_i/c_T$  ( $i=0,2$ ) are comprised between 0 and 1, and  $rc_0+rc_1+rc_2=1$ .

We can reduce the number of parameters in  $rc$  replacing  $rc_2$  by  $1-rc_0-rc_1$  in equation (2). We obtain, after some reordering:

$$rc = [\alpha\rho(1-s_e)+\alpha-1]rc_0 + [\alpha(1-\rho)(1-s_p)+\alpha-1]rc_1 + 1-\alpha \quad (SA3)$$

$rc$  is defined over a triangular region  $T$  in the  $rc_0$ - $rc_1$  plane, bounded by the axes  $rc_0$ ,  $rc_1$  and the line  $rc_0+rc_1=1$  (Figure 3a).  $T$  is conceptually equivalent to the interval  $I=(0,1)$  in the MISC case (see main body of the paper): each point in  $T$  corresponds to a clinical scenario. We will refer to  $T$  as ‘clinical space’ also.  $I$  and  $T$  differ in that the second is two-dimensional; i.e. clinical scenarios are represented by  $(rc_0, rc_1)$  pairs, not by a single value.

## Section S1. Generalizing predictor comparison to all clinical scenarios

For  $N$  predictors, in a clinical scenario defined by a pair  $(rc_0, rc_1)$ , the comparison is straightforward: we only have to compute and sort their respective  $rc$  values. The tool of choice will be the one with the lowest  $rc$ ; it will be the predictor with the lowest MISC+REJ cost. Here, the goal is to solve the more complex problem of comparing predictors across the clinical

space. That is, we want to find a division of this space in regions within which a single predictor prevails in terms of cost. The problem is more difficult for MISC+REJ than for MISC because the space we want to partition is two-dimensional instead of one-dimensional. To illustrate the procedure, we first describe the case of two predictors ( $N=2$ ) and then we extend the idea to an arbitrary number  $N$  of predictors.

Let  $M_i$  and  $M_j$  be two pathogenicity predictors, and  $rc(M_i)$  and  $rc(M_j)$  be their respective  $rc$ 's. We seek a division of  $\mathbf{T}$  into two regions:  $r_i$ , where  $M_i$  is preferable to  $M_j$  ( $rc(M_i) < rc(M_j)$ ), and  $r_j$ , where the opposite is the case ( $rc(M_i) > rc(M_j)$ ). The boundary between  $r_i$  and  $r_j$  is defined by the condition  $rc(M_i)=rc(M_j)$ , which, using equation (3) for  $rc(M_i)$  and  $rc(M_j)$ , gives:

$$\{\rho[\alpha_i(1-s_{e,i})-\alpha_j(1-s_{e,j})]+\alpha_i-\alpha_j\}rc_0+\{(1-\rho)[\alpha_i(1-s_{p,i})-\alpha_j(1-s_{p,j})]+\alpha_i-\alpha_j\}rc_1+\alpha_j-\alpha_i=0 \quad (\text{SA4})$$

where  $s_{e,k}$ ,  $s_{p,k}$  and  $\alpha_k$  are the sensitivity, specificity and coverage of predictor  $M_k$  ( $k=i,j$ ). Equation (4) shows that the boundary sought is a line (Figure 3a), which we will call  $l_{ij}$ , in the  $rc_0$ - $rc_1$  plane.

When  $l_{ij}$  crosses  $\mathbf{T}$ , it divides it into two convex polygons (Figure 3a), corresponding to the  $r_i$  and  $r_j$  regions. If  $l_{ij}$  does not cross  $\mathbf{T}$ , then only one of the two methods will have the lowest  $rc$  in all  $\mathbf{T}$  points.

From Equation (4), we see that  $l_{ij}$  depends on  $\rho$ ; consequently, different values of this parameter may change  $r_i$  and  $r_j$  (Supplementary Figure S3). In this appendix we concentrate on the problem of dividing  $\mathbf{T}$  when more than two predictors are available, keeping the value of  $\rho$  fixed. In the main body of the paper, we describe how different values of  $\rho$  affect the resulting division.

To generalize the comparison to more than two predictors ( $M_i$ ,  $i=1,N$ ;  $N>2$ ), we will develop a procedure that divides  $\mathbf{T}$  into  $m$  regions,  $\{r_k, k=1,m\}$ , such that only one method per region has the lowest  $rc$ . (Note that  $m \leq N$ , since there may be predictors that are never better

than the others). The next two sections are devoted to provide some fundamental results underlying this procedure and then to describe it, first in geometric (Section S2) and then in computational (Section S3) terms.

## **Section S2. Dividing $T$ into a set of regions $\{r_k, k=1,m\}$ in which only one predictor per region has the lowest $rc$**

Note. In our proofs we use several results about convex polygons that can be easily found in the books of Lee [8] and of Yaglom and Botyanskii [9]. The most important ones are explicitly cited.

Here, we address, in geometric terms, the problem of finding the  $\{r_k, k=1,m\}$  regions in which only one predictor per region has the lowest  $rc$ . The starting point is the set of  $N$  predictors that we want to compare. Previously, we have seen that there is a line associated to each pairwise comparison between predictors (Equation (4)). Therefore, after doing all the possible  $M=N.(N-1)/2$  pairwise comparisons, we obtain a set  $L_N=\{l_{ij}, i=1,N-1 \text{ and } j=i+1,N\}$  of lines. A first important result (Proposition S1) is that these lines cut  $T$ , producing a division of this triangle into a set  $P_N$  of convex polygons that will be key to finding the  $r_k$  regions.

**Proposition S1.** Let  $N \in \mathbb{N}$  be an arbitrary number of methods and let  $L_N=\{l_{ij}, i=1,N-1 \text{ and } j=i+1,N\}$  be the set of lines resulting from all the pairwise comparisons between these methods, using  $rc$ . These lines cut  $T$  into a set  $P_N=\{p_l\}$  of convex polygons.

*Proof. By induction.*

Base case. For  $N=2$ , there is only one line in  $L_N$ , since there is only one comparison between two methods. When the line contains no interior point of  $T$ , either because it does not intersect with  $T$ , or because it is a supporting line of it,  $P_2$  will have a single element,  $T$ , which is convex

because it is a triangle. If the line contains at least a point interior to  $T$ , then it will cut  $T$  into exactly two points [9]. The line segment uniting these two points is a chord of the polygon [8] and, by the ‘Polygon Splitting Theorem’ [8], divides  $T$  into two convex polygons.

Induction step. Here we show that if the proposition is true for  $N$ , then it is true for  $N+1$ . That is, we want to show that if  $L_N$  divides  $T$  into a set of convex polygons  $P_N$ , then  $L_{N+1}$  divides  $T$  into a new set of convex polygons that we will call  $P_{N+1}$ .

We know that the set of lines resulting from the comparison of  $N+1$  methods,  $L_{N+1}$ , will contain the lines corresponding to the comparisons between the first  $N$  methods,  $L_N$ , and between these  $N$  methods and an additional  $(N+1)$ th method,  $\{l_{i,N+1}\}_{i=1,N}$ , that is:

$$L_{N+1} = L_N \cup \{l_{i,N+1}, i=1, N\} \quad (\text{SA5})$$

Cutting  $T$  with the lines in  $L_{N+1}$  is equivalent to cutting it with the lines in  $L_N$  and then with those in  $\{l_{i,N+1}, i=1, N\}$  since order is irrelevant to the final result. Therefore,  $P_{N+1}$  will be the result of cutting the polygons in  $P_N$  with the lines in  $\{l_{i,N+1}, i=1, N\}$ . When we cut  $P_N$  with the first line,  $l_{1,N+1}$ , we create a new division of  $T$  in which each of the polygons split by  $l_{1,N+1}$  will be replaced by two children polygons (i.e., no polygon traversed by a line remains in the new division of  $T$ ). Next, we will repeat this process for the remaining lines in  $\{l_{i,N+1}, i=1, N\}$  until we obtain  $P_{N+1}$ . At the end of each step, the division of  $T$  will be constituted by the set of  $P_N$  polygons unaffected by the  $l_{i,N+1}$  line (these polygons are convex because the proposition is true for  $N$ ), and by the children of the affected polygons. Given that the affected polygons are convex (again because the proposition is true for  $N$ ), the children will also be convex, by the ‘Polygon Splitting Theorem’ [8]. Therefore, at the end of each step, the resulting division of  $T$  will be constituted by a set of convex polygons and, consequently,  $P_{N+1}$ , which is obtained at the end of the final step, will be constituted by convex polygons only. QED.

The polygons in  $P_N$  have several characteristics that are relevant for the computational algorithm used to list them (described in Section S3.3). First, their edges are noncollinear line segments that belong either to the  $L_N$  lines or to the three segments defining  $T$ . Second, their vertices can be: the  $T$  vertices, the intersection points between the  $L_N$  lines, and the intersection points between these lines and the triangle edges. Third, for each  $L_N$  line, the segment delimited by the intersection points of the line with the triangle is formed by a concatenation of edges from  $P_N$  polygons (Supplementary Figure S4). And fourth, the same happens for the three segments defining  $T$ , which are formed by a concatenation of edges from  $P_N$  polygons.

The  $P_N$  polygons also satisfy the following lemma.

**Lemma S1.** Let  $p$  be a polygon from  $P_N$ . Then none of its interior points belong to another polygon  $q \in P_N$ .

*Proof. By contradiction.* Let us assume that there exists a polygon  $p \in P_N$  such that one of its interior points belongs to  $q \in P_N$ . This point will belong to one of the edges of  $q$ . The line from  $L_N$  containing this edge will cut  $p$  at two points [9], that is, it will traverse  $p$ . This is in contradiction with the procedure utilized to generate  $P_N$ , in which any polygon traversed by a line from  $L_N$  is removed from the polygon list and replaced by the two children polygons. Therefore,  $p$  does not exist. QED.

Finally, we will show a key property of the  $P_N$  polygons, used to build the regions  $r_k$ .



**Proposition S2.** Let  $P_N$  be the set of convex polygons obtained after dividing  $T$  using  $L_N$ , the set of lines associated to the pair comparisons between  $N$  methods. For each polygon  $p \in P_N$ , the lowest  $rc$  value at all its interior points always corresponds to the same method.

*Proof. By contradiction.* Let us assume that the proposition is not true. That is, that there exists a polygon  $p \in P_N$  with an interior point  $m$  such that  $M_i$ , the method with the lowest  $rc$  value at  $m$ , is different from  $M_j$ , the method with the lowest  $rc$  value at the remaining interior points of  $p$ . Let us consider  $n$ , one of these remaining interior points. Then, according to our assumption, the  $rc$  values of  $M_i$  and  $M_j$  at  $m$  ( $rc_m(M_i)$  and  $rc_m(M_j)$ , respectively) and at  $n$  ( $rc_n(M_i)$  and  $rc_n(M_j)$ , respectively) satisfy the following inequalities:  $rc_m(M_i) < rc_m(M_j)$  and  $rc_n(M_i) > rc_n(M_j)$ , respectively. This is in contradiction with the fact that the boundary line between  $M_i$  and  $M_j$  does not pass between  $m$  and  $n$ , because during the construction of  $P_N$  (Proposition S1) any polygon traversed by a line from  $L_N$  is removed from  $P_N$  and replaced by the resulting children polygons. QED.

We will say that a method  $M_k$  is associated to a polygon  $p$ , when  $M_k$  has the lowest  $rc$  value for the interior of  $p$ . Note, that  $M_k$  will also have the lowest  $rc$  value at the line segments defining  $p$ .

The polygons in  $P_N$  do not necessarily coincide with the  $r_k$  regions but can be used to obtain them using the following procedure.

**Step 1. Find the method associated to each polygon.** For each polygon in  $P_N$  we apply the next three steps:

**Step 1.1.** Compute the average of its vertices, which is a point belonging to the interior of the polygon because the polygon is convex.

**Step 1.2.** Compute the  $rc$  value, at this average point, for each of the  $N$  methods.

**Step 1.3.** Sort the  $N$  methods according to their  $rc$ 's at the average point and choose the method with the lowest  $rc$  value. By Proposition S2, this method prevails (has the lowest  $rc$ ) at all the points interior to the polygon considered. This will be the method associated to the polygon.

**Step 2. Obtaining the  $\{r_k, k=1,m\}$  regions.** Each region  $r_k$  (Supplementary Figure S5) is obtained as the union of all polygons associated to the same method,  $M_k$ :

$$r_k = \bigcup_{i \in \Omega_k} p_i \quad (\text{SA6})$$

where  $\Omega_k$  are the indexes of the  $p_i$  polygons in  $P_N$  for which  $M_k$  has the lowest  $rc$  value. We will say that  $M_k$  is the method associated to the region  $r_k$ .

By application of Proposition S2 to the polygons associated to  $M_k$ , we know that there is only one method associated to each  $r_k$ . By the same reason, we know that no other region contains a point for which  $M_k$  is the method with the lowest  $rc$ . That is, by construction, there are no two regions with the same associated method; therefore,  $\{r_k, k=1,m\}$  is the set we are looking for.

Note. For clarity purposes, in this work we do not explicitly treat the fact that, at the boundary between adjacent  $r_k$  regions, two methods have the same  $rc$  value. This fact has no impact on the results presented, neither in the finding of the  $r_k$  regions, nor in computing their surfaces, etc.

**Section S3. Computational obtention of the  $P_N$  polygons using an adapted Breadth First Search (BFS)**

As we have seen in the previous section, once we know the polygons in  $P_N$  it is trivial to obtain the  $r_k$  regions, using (6). Here, we describe how we can obtain these polygons computationally. In particular, we show that the problem of finding a  $P_N$  polygon, in terms of its constituting vertices, can be modeled as a graph problem, and solved with an adapted version of the BFS algorithm that we will call aBFS. A python implementation of this procedure, CSP-rej (Clinical Space Partition, rejection), is available at:

[https://github.com/ClinicalTranslationalBioinformatics/clinical\\_space\\_partition](https://github.com/ClinicalTranslationalBioinformatics/clinical_space_partition)

This code reproduces the results presented in this work and allows users to explore other combinations of predictors. It must be noted that, when planning comparisons of large numbers of predictors, it is preferable to partition the original set of predictors into smaller sets, and run the program separately for each set. Then, aggregate the surviving methods from these separate runs and execute the program again. This approach will reduce the risk of numerical exceptions that appear in geometric computations [10], particularly when working with low  $\rho$  values. It can be repeated as many times as desired.

#### *Preliminary results: computing the set of edges and vertices of the $P_N$ polygons*

The first step in our approach is to compute the set of lines  $L_N = \{l_{ij}, i=1, N-1 \text{ and } j=i+1, N\}$ , applying equation (4) to all possible comparisons between the  $N$  methods.

The next step is to build  $VP$  and  $EP$ , the sets of vertices and edges, respectively, of the  $P_N$  polygons. For  $VP$ , we first compute the intersection between the lines in  $L_N$ , keeping only the points falling inside  $T$ . These points are included in  $VP$ . Then, we compute the intersection between the lines in  $L_N$  and the boundaries of  $T$ . The resulting points are added to  $VP$ . Finally, we include in  $VP$  the three vertices of  $T$ . The resulting set of vertices is used to obtain  $EP$ , which is constituted by all the  $\overline{v_i v_j}$  ( $v_i, v_j \in VP$ ) line segments uniting two consecutive vertices

in an  $L_N$  line or in the lines defining  $T$ . Note that every element of  $EP$  corresponds to the edge of a  $P_N$  polygon, as shown in the next lemma.

**Lemma S2.** Any segment in  $EP$  is the edge of at least one polygon in  $P_N$ .

*Proof.* By contradiction. We assume that there exists a segment  $\overline{v_i v_j} \in EP$  which is not the edge of any  $P_N$  polygon. In the following, we show that the possible options for  $\overline{v_i v_j}$  lead to a contradiction.

By construction of  $EP$ ,  $\overline{v_i v_j}$  belongs either to one of the lines in  $L_N$ , or to one of the three lines defining  $T$ . In all cases, within  $T$  these lines are formed by a concatenation of edges from polygons in  $P_N$  (Supplementary Figure S4). Therefore,  $\overline{v_i v_j}$  will overlap with some of these edges. Two situations are possible. It may happen that  $\overline{v_i v_j}$  spreads over two or more edges. In this case, some of the vertices of these edges will be comprised between  $v_i$  and  $v_j$ , in contradiction with the fact that  $v_i$  and  $v_j$  are consecutive. A second possibility is that  $\overline{v_i v_j} \subsetneq \overline{v_k v_l}$ , where  $\overline{v_k v_l}$  is one of the edges in the line. This is in contradiction with the fact that  $v_k$  and  $v_l$  are consecutive. Consequently,  $\overline{v_i v_j}$  must be an edge of a  $P_N$  polygon. QED.

Before describing how we obtain the  $P_N$  polygons, we need to prove a result about polygons sharing more than one edge that will be used to prove our computational procedure.

**Proposition S3.** Let  $p \in P_N$  be a polygon with two edges  $\overline{v_i v_j}, \overline{v_i v_k} \in EP$  forming a consecutive angle (Supplementary Figure S6). There exists no other convex polygon  $q$ , different from  $p$ , with  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$  among its edges and the remaining edges belonging to  $EP$ , and such that none of its interior points belongs to another polygon in  $P_N$ .

*Proof. By contradiction.* We will assume that  $q$  exists and explore the different possibilities that arise, showing that they all lead to contradiction. In particular, we will focus our reasoning on the relative position of the points in  $p$  and  $q$  outside  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$ . There are three possibilities, considering that both  $p$  and  $q$  are convex.

If in  $q$  these points are all interior to  $p$ , then the edges of  $q$  joining  $v_j$  and  $v_k$  will be interior to  $p$  (Supplementary Figure 7a). Because, these edges belong to  $EP$ , they then necessarily belong to polygons in  $P_N$  (by Lemma S2). This is in contradiction with the fact that none of the interior points of  $p$  belong to another  $P_N$  polygon (by Lemma S1).

If all the points in  $p$  outside  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$  are interior to  $q$  (Supplementary Figure S7b) then there are interior points of  $q$  will belonging to a polygon in  $P_N$ , because  $p \in P_N$ . This is in contradiction with the definition of  $q$  as having no interior points from  $P_N$  polygons.

Finally, we reach a similar contradiction when both  $p$  and  $q$  have interior points from each other (in this case, the points may come from full segments or fragments of segments) (Supplementary Figure S7c). QED.

### *Building $P_N$ with a graph-based approach*

The list of polygons in  $P_N$  can be obtained looping through all the vertices in  $VP$ , enumerating the polygons that meet at each vertex. Here, we show that we can model this polygon enumeration problem as a cycle enumeration problem in graph theory.

Our starting point is the unweighted, undirected graph  $G(V, E)$ , whose set of vertices,  $V$ , and edges,  $E$ , correspond to  $VP$  and  $EP$ , respectively. Because the list of vertices of a polygon is formally equivalent to that of a cycle, we can reformulate the original looping through  $VP$  elements as a looping through  $V$  elements. The lists of vertices of the  $P_N$  polygons meeting at a given vertex will now correspond to cycles in  $G(V, E)$ .

In this search, for each  $v_i \in V$ , we will use BFS as a shortest cycle generator, keeping only those cycles satisfying the following conditions:

- C1. A cycle cannot have more than one edge corresponding to a segment from the same line. This rule is applied to eliminate those sequences of edges produced by the BFS that do not correspond to convex polygons, according to the edge-line lemma [8]. Also, to avoid that edges from collinear segments are included.
- C2. A cycle cannot have repeated vertices, except the first and the last one, as in polygons all the vertices are different except the first and the last one [8].
- C3. Every edge in  $E$  has a counter that is decreased by one each time it is included in a cycle. Once the counter reaches zero, the edge is excluded from future searches. The starting value of the counter of each  $E$  edge depends on the location of its equivalent edge in  $EP$ . If the latter belongs to a side of the triangle, the counter will start at 1; otherwise (when it belongs to a line in  $L_N$ ) it will start at 2. This condition guarantees that, for each edge in  $EP$ , we enumerate all the polygons sharing it, thus ensuring, together with C4, that our polygon enumeration procedure is exhaustive for  $P_N$  elements.
- C4. Every vertex in  $V$  has a counter that is decreased by one each time it is included in a cycle. Once the counter reaches zero, the vertex is excluded from future searches. The starting value of the counter of each  $V$  vertex depends on the environment of its equivalent vertex in  $VP$ . More precisely, it depends on the number and origin of  $EP$  edges that include the  $VP$  vertex (see Supplementary Figure S8). This condition guarantees that, for each vertex, we enumerate all the polygons sharing it, thus ensuring, together with C3, that our polygon enumeration procedure is exhaustive for  $P_N$  elements.

- C5. Once a minimal cycle is found, it is excluded from future searches. This condition is introduced to avoid repetitions in the final list of cycles.
- C6. For every minimal cycle found, we check that the corresponding polygon has no interior points corresponding to vertices in  $VP$ . This condition, together with C7, prevents the inclusion in the final list of cycles of convex polygons not belonging to  $P_N$  (see Lemma S3 and Proposition S4 below).
- C7. For every minimal cycle found, we compute all the possible chords between the vertices of the corresponding polygon. If any of these chords corresponds to an edge in  $EP$ , the polygon is discarded. This condition excludes cycles with a list of vertices in which more than two vertices from the same line are included, thus limiting the chosen cycles to those corresponding to convex polygons (all pairs of non-adjacent edges of the chosen polygon are semiparallel and, by Theorem 8.7 in Lee, the polygon is convex). Condition C7, together with C6, also prevents the inclusion, in the final list of cycles, of convex polygons not belonging to  $P_N$  (see Lemma S3 and Proposition S4 below).

As mentioned above, we call aBFS this combination of BFS and C1-C7 conditions. We will now establish (Lemma S3 and Proposition S4) that a cycle found with aBFS corresponds to a polygon in  $P_N$ .

**Lemma S3.** Let  $c_{SC}$  be a cycle found by aBFS, and constituted by the sequence of edges  $\{v_i, v_{i+1}\} \in E$ , where  $i=1, N$  and  $v_I = v_{N+1}$ . Then,  $p_{sc}$ , the corresponding sequence of segments  $\overline{v_i v_{i+1}}$  from  $EP$ , is a convex polygon and has no interior points from any polygon in  $P_N$ .

*Proof.*  $c_{sc}$  is a cycle whose sequence of edges correspond to a sequence of segments, from  $EP$ , characteristic of a polygon [8]: it starts and ends at the same vertex, there are no collinear segments (condition C1), and it has not repeated vertices (condition C2) other than the first and last ones, which are equal.

$p_{sc}$  is convex because of conditions C1 and C7.

$p_{sc}$  has no interior points from any polygon in  $P_N$  because of conditions C6 and C7.

QED.

**Proposition S4.** Let  $c_{sc}$  be a Shortest Cycle identified by aBFS. Then,  $p_{sc}$ , the polygon corresponding to this cycle, belongs to  $P_N$ .

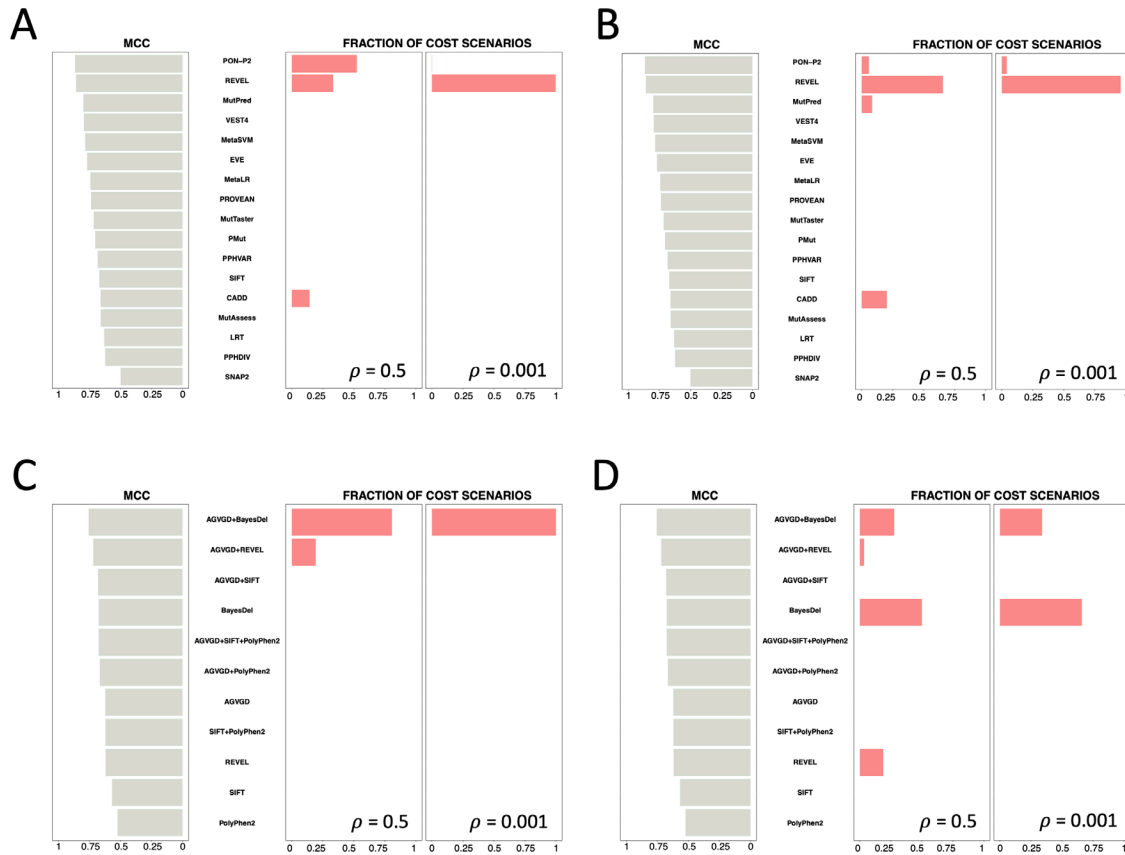
*Proof.* First, we know from Lemma S3 that  $p_{sc}$  is a convex polygon with no interior points from any polygon in  $P_N$ . Now, let us select an arbitrary pair of adjacent edges from  $p_{sc}$ ,  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$ . By construction of EP,  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$  are edges of polygons in  $P_N$  (Lemma S2). Because no line from  $L_N$  passes between them (by conditions C6 and C7),  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$  belong to the same polygon  $p \in P_N$ . From Proposition S3 we know that  $p$  is the only convex polygon with  $\overline{v_i v_j}$  and  $\overline{v_i v_k}$  among its edges and no interior points from any polygon in  $P_N$ , therefore,  $p_{sc} = p$  and  $p_{sc} \in P_N$ . QED.

For each vertex  $v_i$ , aBFS will find the shortest cycles corresponding to all the  $\overline{v_i v_j}, \overline{v_i v_k}$  pairs forming consecutive angles. This procedure will be repeated for all the vertices in  $G(V, E)$ , guaranteeing, through the use of counters (conditions C3 and C4 above), that the number of cycles found matches that of expected  $P_N$  polygons. By condition C5 and Proposition S4, we know that the shortest cycles found are unique and correspond to  $P_N$  polygons, respectively.



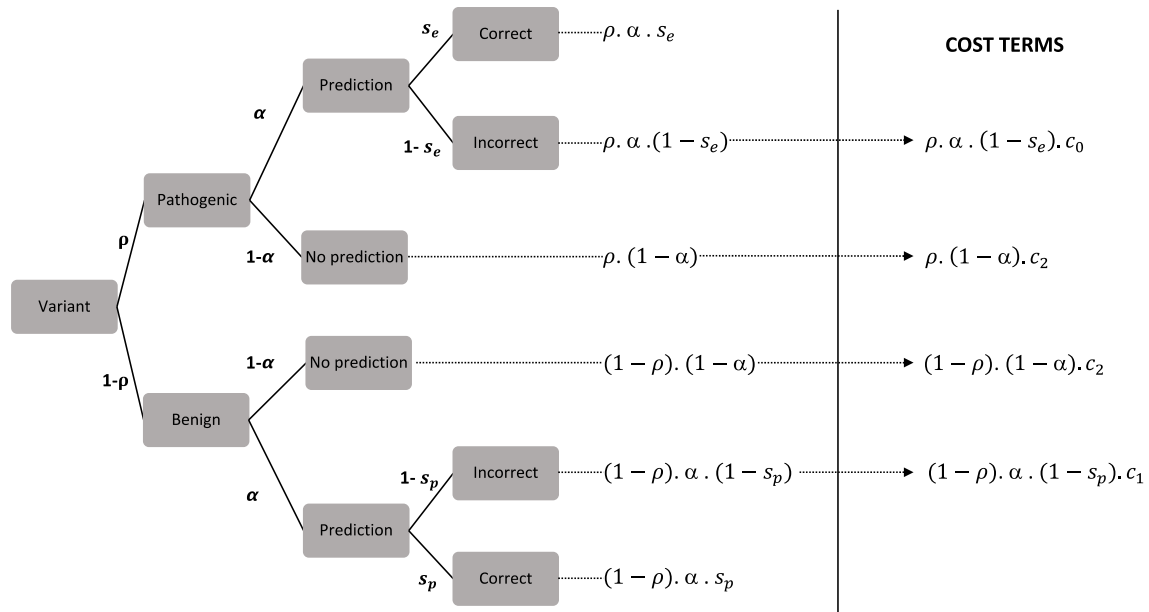
In summary, exhaustive application of aBFS to the vertices in  $E$  will produce the list of polygons in  $P_N$ . Each polygon will be defined by its list of vertices.

## Supplementary Figures

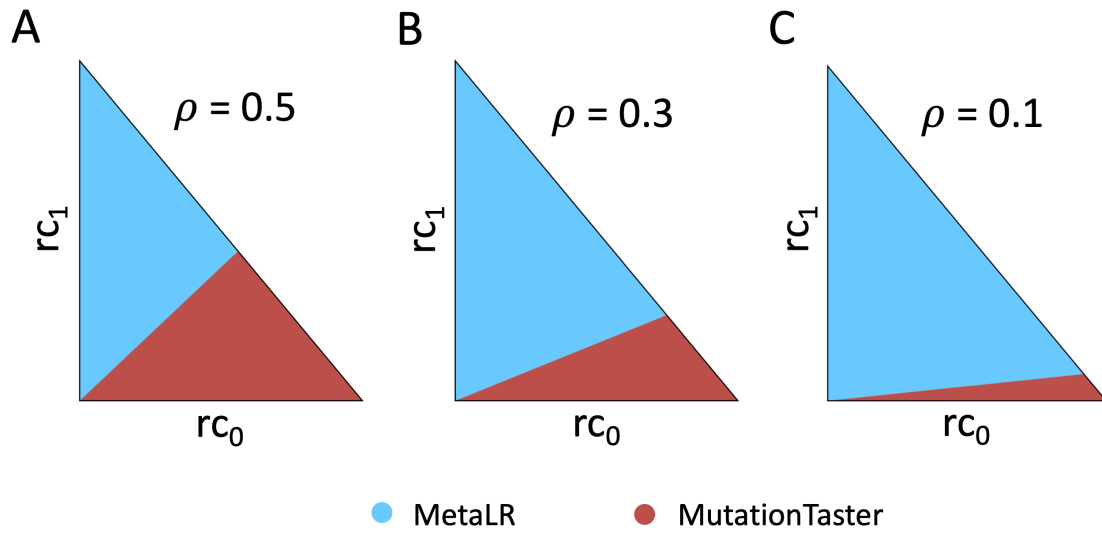


**Supplementary Figure SA1. Comparison between the MCC ranking of the seventeen pathogenicity predictors and their corresponding fractions of cost scenarios.** In all the parts of this figure (a-d), pathogenicity predictors are ranked according to their MCC (grey bars, left side) and fraction of cost scenarios for which each predictor is optimal are represented with pink bars, right side. The results are shown for  $\rho=0.5$  and  $\rho=0.001$ ). **A**, MISC analysis of the seventeen predictors. **B**, MISC+REJ analysis of the seventeen predictors. These figures are equivalent to those shown in Figs. 1b and 3d, respectively, for AUC. We have reproduced the MCC analysis using data for the TP53 gene and several predictors, retrieved from the work of Fortuno et al. [11]: **C**, MISC analysis, and **D**, MISC+REJ analysis. Here the cost space is

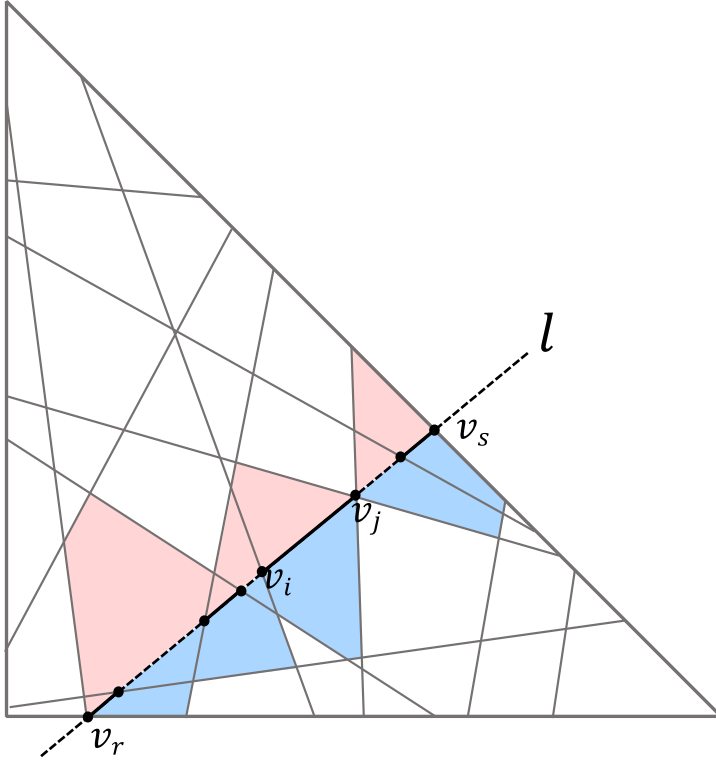
restricted to the Li-Fraumeni syndrome. The results confirm an incomplete correspondence between the MCC analysis and the fraction of cost scenarios in which each method predominates.



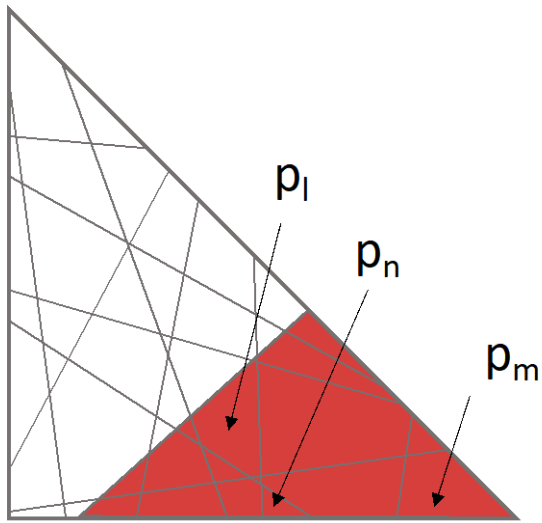
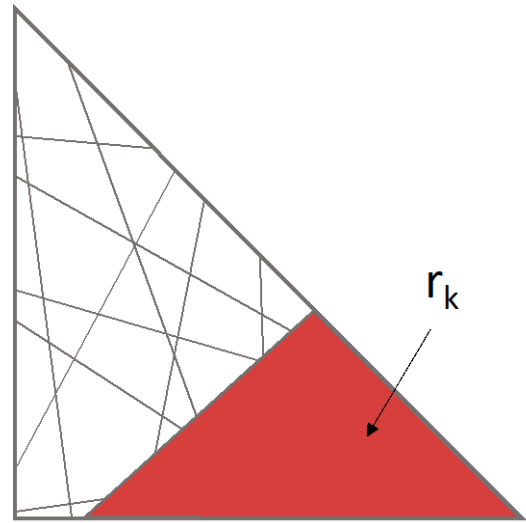
**Supplementary Figure S2. Probabilistic tree diagram underlying the cost-framework presented in this work.** Each branch of the tree corresponds to a different situation in the use of prediction methods, the situation's probability is written by its side. Multiplying probabilities along branches gives the probability of an event that is the combination of different situations. For example, a pathogenic variant can be incorrectly predicted as benign; the probability of this event is:  $\rho \cdot \alpha \cdot (1-s_e)$ . In the cost model MISC+REJ each of these events will contribute a term (shown to the right of the vertical line) that, after summation and reordering, will result in equation (1).



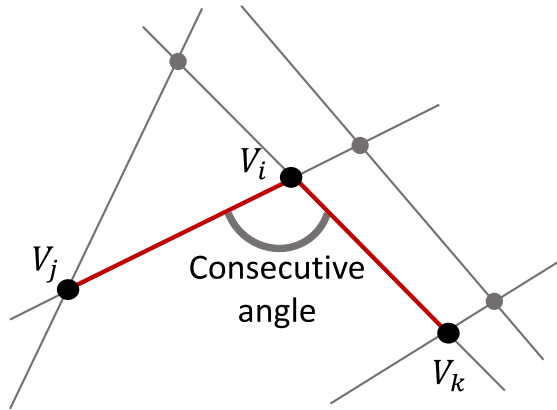
**Supplementary Figure S3. Effect of the fraction of pathogenic variants in the sample ( $\rho$ ) on the distribution of pathogenicity predictors over the cost domain.** This figure shows how different values of  $\rho$  (*a*,  $\rho = 0.5$ ; *b*,  $\rho = 0.3$ ; *c*,  $\rho = 0.1$ ) may substantially alter the size of the cost region assigned to each predictor. A version of this figure for the case of seventeen predictors is presented in Figure 4.



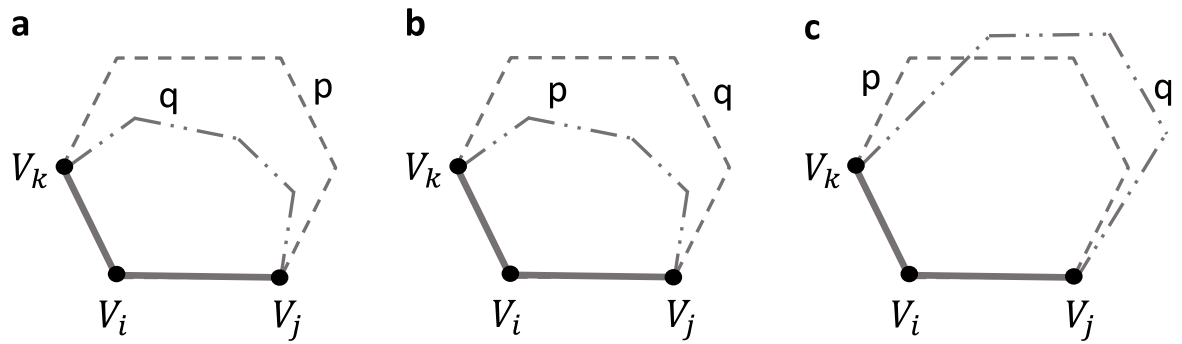
**Supplementary Figure S4. Lines from  $L_N$  are constituted by a concatenation of edges from  $P_N$  polygons, when crossing the triangle  $T$ .** The figure shows how a line  $l \in L_N$  is formed by a succession of segments that correspond to edges of  $P_N$  polygons (pink/blue), when traversing the triangle  $T$ , i.e., between the points  $v_r$  and  $v_s$ . As illustrated for  $\overline{v_i v_j}$ , each segment is shared by two polygons, one above (pink) and the other below (blue) the segment.

**a****b**

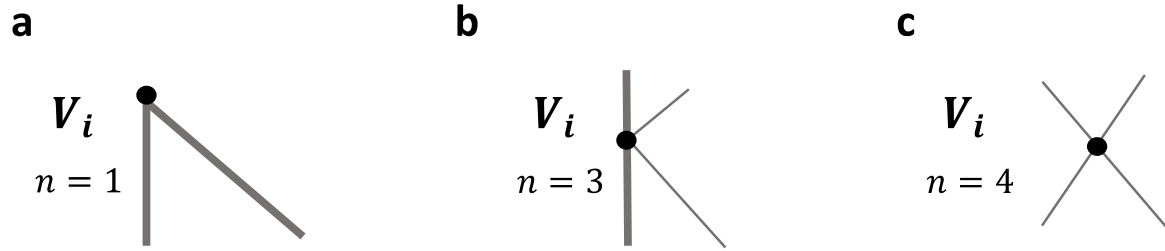
**Supplementary Figure S5. From polygons to regions.** Unification of the polygons in which the same method has the lowest average cost (**a**, shown in red) results in a more simplified view of the regions assigned to each predictor (**b**).



**Supplementary Figure S6. Illustration of the consecutive angle formed by vertices  $v_i$ ,  $v_j$ , and  $v_k$ .**



**Supplementary Figure S7. Different situations for the relative location of polygons  $p$  and  $q$ .** **a**, all the points in  $q$  are interior to  $p$ ; **b**, all the points in  $p$  are interior to  $q$ ; and **c**, both  $p$  and  $q$  have interior points from the other.



**Supplementary Figure S8. Initial values of the counter of each vertex.** In the three figures, thick grey lines correspond to the edges of the triangle  $T$  (the clinical space, see Appendix 1, Supplementary Materials), and thin grey lines correspond to the lines dividing  $T$  and associated with the different pair comparisons between methods. **a**, The vertex is one of the three triangle vertices. **b**, The vertex is the intersection between a triangle edge and the line associated with the comparison between two predictors. **c**, The vertex is the intersection between the lines associated to two different pair comparisons between predictors.



## Supplementary Tables

**Supplementary Table S1. The seventeen pathogenicity predictors used in this work.** We provide the performance parameters required for the cost computations: sensitivity, specificity and coverage/reject rate. The column Output homogenization shows the correspondence between our pathogenic/benign states and the output of each predictor. We also list the decision cutoff when it was not provided by dbNSFP.

Predictor	Sensitivity	Specificity	Coverage/Reject rate	Output homogenization & decision cutoff
<b>CADD</b>	1	0.68	1/0	P $\geq$ 15 B < 15
<b>EVE</b>	0.92	0.85	0.43/0.57	P = Pathogenic B = Benign
<b>LRT</b>	0.87	0.76	0.87/0.13	P = Deleterious B = Neutral
<b>MetaLR</b>	0.87	0.88	0.99/0.01	P = Deleterious B = Tolerated
<b>MetaSVM</b>	0.9	0.89	0.99/0.01	P = Deleterious B = Tolerated
<b>MutPred</b>	0.95	0.87	1/0	P $\geq$ 0.5 B < 0.5
<b>MutationAssessor</b>	0.89	0.77	0.86/0.14	P = High, Medium B = Low, Neutral
<b>MutationTaster</b>	0.98	0.74	0.99/0.01	P = Disease causing automatic, disease causing B = Polymorphism, polymorphism automatic
<b>PMut</b>	0.84	0.87	0.85/0.15	P = Disease B = Neutral
<b>PON-P2</b>	0.96	0.92	0.46/0.54	P = Pathogenic B = Neutral
<b>PROVEAN</b>	0.9	0.84	0.97/0.13	P = Deleterious B = Neutral
<b>Polyphen2_HDIV</b>	0.93	0.69	0.91/0.09	P = Probably damaging, possibly damaging B = Benign
<b>Polyphen2_HVAR</b>	0.9	0.78	0.91/0.09	P = Probably damaging, possibly damaging B = Benign
<b>REVEL</b>	0.92	0.94	1/0	P $\geq$ 0.5 B < 0.5

<b>SIFT</b>	0.93	0.75	0.97/0.03	P = Damaging B = Tolerated
<b>SNAP2</b>	0.85	0.65	0.82/0.18	P = Effect B = Neutral
<b>VEST4</b>	0.89	0.9	1/0	P >= 0.5 B < 0.5

**Supplementary Table S2. Predictive performance of the thirteen pathogenicity predictors**

**studied in Pejaver et al. [12].** In the first three columns we provide the performance parameters required for the cost computations (sensitivity, specificity and coverage/reject rate). In the last column, we give the score thresholds that define the pathogenic (P) and benign (B) classes for each predictor, according to Pejaver et al.'s (see Table 2 in [13]). The 'not classified' class pertains to variants with a prediction score falling between that of the pathogenic and benign classes. Note that for the pathogenic and benign classes we have unified the four levels (Supporting, Moderate, Strong, and Very Strong) provided by the authors.

<b>Predictor</b>	<b>Sensitivity</b>	<b>Specificity</b>	<b>Coverage/Reject rate</b>	<b>Output homogenization &amp; decision cutoff</b>
<b>BayesDel</b>	0.916	0.906	0.797/0.203	P >= 0.130 B <= -0.180
<b>CADD</b>	0.912	0.831	0.782/0.218	P >= 25.3 B <= 22.7
<b>EA</b>	0.893	0.882	0.556/0.444	P >= 0.685 B <= 0.262
<b>FATHMM</b>	0.981	0.526	0.170/0.830	P <= -4.140 B >= 3.320
<b>GERP++</b>	0.000	1.000	0.237/0.763	B <= 2.700
<b>MPC</b>	1.000	0.000	0.192/0.808	P >= 1.360
<b>MutPred2</b>	0.902	0.921	0.773/0.227	P >= 0.737 B <= 0.391
<b>PhyloP</b>	0.890	0.792	0.618/0.382	P >= 7.367 B <= 1.879
<b>PolyPhen2</b>	0.895	0.833	0.550/0.450	P >= 0.978 B <= 0.113

<b>PrimateAI</b>	0.914	0.844	0.572/0.428	P >= 0.790 B <= 0.483
<b>REVEL</b>	0.923	0.913	0.732/0.268	P >= 0.644 B <= 0.290
<b>SIFT</b>	0.903	0.824	0.559/0.441	P <= 0.001 B >= 0.080
<b>VEST4</b>	0.928	0.902	0.723/0.277	P >= 0.764 B <= 0.449

**Supplementary Table S3. Predictive performance of the rules for computational evidence in the two *ATM*-adapted versions of the ACMG/AMP guidelines.** In the first three columns we provide the performance parameters required for the cost computations (sensitivity, specificity and coverage/reject rate). In the last column, we give the score thresholds that define the pathogenic (P) and benign (B) classes for each predictor, according to Clingen's expert panel [14] and according to Feliubadalo et al. [15]. The 'not classified' class pertains to variants with a prediction score falling between that of the pathogenic and benign classes.

Predictor	Sensitivity	Specificity	Coverage/Reject rate	Output homogenization & decision cutoff
ClinGen	1.000	0.985	0.825/0.175	P $\geq$ 0.733 B $\leq$ 0.249
Feliubadaló	0.922	0.945	0.948/0.052	P = Pathogenic B = Benign

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